Constraint Programming

Search
(version of 23 November 2021)
Reminder: Branching = 2-Step Choice

- Variable selection.
- Value selection & partition selection.

The branching decisions can have a strong impact on the size of the search tree.
Black-Box Search: The Vision

- Constraint Program = Model + Search

The vision is that we should not have to worry too much about the search. A CP solver should be smart enough to learn or select a good search for the model.

- Since ~2004, many research efforts have been devoted to making this vision true.
The leftmost feasible solution is very important for pruning during branch-and-bound search.

Solution quality is mostly impacted by the value selection heuristic for optimization problems.

The decisions down the leftmost branch should look like a greedy algorithm.
Example: Traveling Salesperson Problem

Greedily, we would adopt a nearest-neighbor value selection heuristic:
Example: Traveling Salesperson Problem
Example: Traveling Salesperson Problem
Example: Traveling Salesperson Problem
Example: Traveling Salesperson Problem

Can we automate this for *any* problem in order to find a good first feasible solution?
Branching = 2 steps:
1. Variable selection.
2. Value selection & partition selection: bound-impact value selector (BIVS) selects a value with the smallest objective lower bound after propagation of the fixing of the selected variable to this value (when minimizing).

Jean-Guillaume Fages and Charles Prud'homme.
Making the first solution good!
*International Conference on Tools with Artificial Intelligence (ICTAI),* 2017.
Making the First Solution Good!

- Minimize $O = f(x_1, x_2, \ldots, x_{10})$ subject to some constraints.
- The variable selection heuristic selects $x_3$.
- The current domain is $D(x_3) = \{v_1, v_2, v_3, v_4\}$.

```py
best = null
bestObj = +inf
for v in {v_1, v_2, v_3, v_4}:
    sm.saveState()
    cp.post(x_3 = v)
    if min(O) < bestObj:
        best = v
        bestObj = min(O)
    sm.restoreState()
branch {cp.post(x_3 = best)}
{cp.post(x_3 != best)}
```
Making the First Solution Good: Advice

- This procedure (for making the first solution good) is quite costly: fixpoint computation for every value + save & restore the state.
- The advice is to do this only for discovering the first feasible solution and then to use a more naïve (but faster) value selection heuristic.
Value Selection: Phase Saving

- If a value v was successfully used for fixing a selected variable x – that is: if branching on x=v did not lead to failure of the fixpoint algorithm then store the value v as the **last success value** of x.

- Each time the variable x is selected by the variable selection heuristic, first try its last success value.

```java
SearchStatistics stats = makeDfs(cp,
   selectMin(q,
      qi -> qi.getSize() > 1,
      qi -> qi.getSize(),
      qi -> {
         int v = qi.getMin();
         return branch(() -> equal(qi,v),
                        () -> notEqual(qi,v));
      }));
```

One can easily implement a default search doing phase saving.
Branching

Reminder:
• Variable selection.
• Value selection & partition selection.

For the rest of this lecture, we focus on variable selection.
“To succeed, try first where you are most likely to fail.”

Robert M. Haralick and Gordon L. Elliott. 
Increasing tree search efficiency for constraint satisfaction problems. 
*International Joint Conference on Artificial Intelligence (IJCAI), 1979.*
First-Fail Principle

**First-fail** for variable selection:

Since all variables must eventually be fixed, if there are no solutions under a node (failure), then we prefer to detect this as soon as possible, so that not too much time is spent exploring the subtree under the node.

“To succeed, try first where you are most likely to fail.”

Robert M. Haralick and Gordon L. Elliott.
Increasing tree search efficiency for constraint satisfaction problems. *International Joint Conference on Artificial Intelligence (IJCAI), 1979.*
First-Fail Principle for Variable Selection

- Can be implemented in various ways:
  - Min-Dom
  - Dom+Deg
  - Dom/Deg
  - Dom/Wdeg
  - Impact-based search
  - Activity-based search
  - Last-conflict search
  - Conflict-ordering search
  - ...

Min-Dom Heuristic

- **Min-Dom**: Select an unfixed variable with the smallest domain size.
- This heuristic was shown experimentally to minimize search-tree depth.
- It is quite intuitive that for $D(x_1) = \{1,2\}$ and $D(x_2) = \{1,\ldots,100\}$ branching on $x_1$ first is likely to trigger more propagation.
Min-Dom Heuristic

```java
int n = 8;
Solver cp = Factory.makeSolver(false);
IntVar[] q = Factory.makeIntVarArray(cp, n, n);
// constraints ...

DFSearch search = Factory.makeDfs(cp, () -> {
    IntVar qs = selectMin(q,
        qi -> qi.size() > 1,
        qi -> qi.size());
    if (qs == null) return EMPTY;
    else {
        int v = qs.min();
        return branch(() -> Factory.equal(qs, v),
                    () -> Factory.notEqual(qs, v));
    }
});

search.onSolution(() -> println("solution:" + Arrays.toString(q)));
SearchStatistics stats = search.solve();
```

Select an unfixed variable with the smallest domain size.
Degree-Based Heuristics

If a variable is involved in many constraints, then it is likely that any filtering of the domain of this variable will trigger some filtering for other variables of those constraints.
Degree of a variable $x =$ the number of constraints on $x$.

- **Dom+Deg**: Min-Dom, and break ties with the degree.
- **Dom/Deg**: Select an unfixed variable with the smallest ratio between domain size and degree.
Weighted Degree (Wdeg) Heuristic

- Idea: Introduce learning to find out what the “difficult” variables are.
- Each time a constraint fails, its weight is increased (by +1).
  - The idea is that if a constraint has failed a lot in the past, then it will probably continue this trend in the future.
- Weighted degree of a variable $x$ = the sum of the weights of the constraints on $x$ and at least one other unfixed variable.
  - Dom/Wdeg: Select the first unfixed variable with the smallest ratio between domain size and weighted degree.
  - Weakness: A constraint that fails may not be the only guilty one for the failure (cascade of propagations in the fixpoint algorithm).
Dom/Wdeg Implementation

- Not shipped with MiniCP, but quite easy to implement.
- Each constraint should be aware of the variables in its scope.
- Add a method called `scope` to the class `Constraint`:
  it returns the list of variables in the scope of the constraint.
Impact-Based Search (IBS)

- Idea: Take a branching decision that has the lowest estimated impact in terms of filtering.
- Preliminaries:

\[ P = \langle X, D, C \rangle \]

\[ S(P) = \prod_{x \in X} |D(x)| \]

**CSP**

**estimate of the size of the search tree**

Philippe Refalo.
Impact-based search strategies for constraint programming.
Impact-Based Search (IBS)

\( P_0 = \langle X, D, C \rangle \)

\[ c_1 \]

\( P_1 = \langle X, D, C \cup \{c_1\} \rangle \)

\[ C_{k-1} = C \cup \{c_1, \ldots, c_{k-1}\} \]

\( P_{k-1} = \langle X, D_{k-1}, C_{k-1} \rangle \)

\[ c_k : x = a \]

\( P_k = \langle X, D_k, C_k \rangle \)

Contraction of search space with respect to parent node:

\[ I^k(x = a) = 1 - \frac{S(P_k)}{S(P_{k-1})} \]

= 0 (no contraction)
= 1 (failure, full contraction)
Impact-Based Search (IBS)

- We can estimate over a set $K$ of search nodes the average impact of $x=a$:

$$\bar{I}(x = a) = \frac{1}{|K|} \cdot \sum_{k \in K} I^k(x = a)$$

- This estimate can also be updated at every node, with a forget factor, instead of being averaged:

$$\bar{I}_k(x = a) = (1 - \alpha) \cdot \bar{I}_{k-1}(x = a) + \alpha \cdot I^k(x = a)$$

Exponential moving average $\alpha \in [0,1]$: the higher $\alpha$ is, the faster it will discount older observations.
Impact-Based Search (IBS)

- Estimation of the size of the search tree when trying $x=a$:

$$S(P) \cdot (1 - \bar{I}(x = a))$$

- Estimation of the size of the search tree if labeling the variable $x$:

$$\sum_{a \in D(x)} S(P) \cdot (1 - \bar{I}(x = a))$$

- Impact of a variable:

$$I(x) = \sum_{a \in D(x)} (1 - \bar{I}(x = a))$$

**Variable heuristic:** Select $x$ with minimum $I(x)$

**Value heuristic:** Select $a$ with minimum $\bar{I}(x = a)$

- Constant whatever the variable we branch on

$= 0$ (no contraction) 
$= 1$ (failure, full contraction)
Activity-Based Search (ABS)

- Idea: Track how often a variable domain is contracted during search, as this allows some learning.
- We expect that a variable whose domain is often contracted is a good candidate to branch on early in order to reduce the depth of the search tree.
Activity-Based Search (ABS)

Each time a search choice is executed, increment the counter $A(x)$, denoting the activity of $x$, of every variable $x$ having a domain contraction:

$$P = \langle X, D, C \rangle$$

$$x = a$$

$X' = \text{variables with domain changed}$

$X' \subseteq X$

$$\forall x \in X \text{ where } |D(x)| > 1 : A(x) := A(x) \cdot \gamma$$

$$\forall x \in X' : A(x) := A(x) + 1$$

Select $x$ maximizing

$$\frac{A(x)}{|D(x)|}$$
Initialization of Learning Heuristics

Activity-based search and impact-based search are initialized with some random dives before starting a complete search:
Can be used in combination with another heuristic; let us call it `fallbackHeuristic`.

`lastConflictVariable` stores a reference, initialized to `null`, to a decision variable: this variable is the last one we branched on that led to failure of the fixpoint algorithm.

```java
nextVarToBranchOn(X) {
    if (lastConflictVariable == null) {
        return fallbackHeuristic(X);
    } else {
        return lastConflictVariable;
    }
}
```

When branching on a variable leads to a failure, always branch on this variable first until it is successfully fixed.

Impossible to fix x4 to some value. Why? It may be in conflict with the choices on x3, x2, or x1. Since we do not know, we should always try to branch on x4 first, also after backtracking, as long as this conflict persists.
Last-Conflict Search: Example

1. `lastConflictVariable = null`

2. `y = c`
   - `x = a`
   - `x ≠ a`

3. `x = a`

4. `x ≠ b`

5. `lastConflictVariable = x`

6. `x = b`

7. `y ≠ c`
   - `x = a`

8. `lastConflictVariable = null: the next decision need not be on x.

Backtrack and select `x` during variable selection.

We branch on `x` rather than `y` or another variable.
Conflict-Ordering Search (COS)

- Generalization and extension of last-conflict search:
  - Each branching decision is timestamped with a shared counter, which is incremented at each failed visited node of the search tree.
  - Each variable has a timestamp: it is the timestamp of the most recent node where a failure occurred when branching on this variable, and null if it caused no failure so far.
  - COS: Select the unfixed variable with the largest timestamp. If no unfixed variable has a failure timestamp yet, then use a fallback heuristic.

Conflict-Ordering Search (COS)

- A **timestamp** is associated with each variable and is increased (to the node timestamp) at each failure it causes.
- **COS**: Select the variable with the largest timestamp.

### Example

<table>
<thead>
<tr>
<th>Variable</th>
<th>x</th>
<th>y</th>
<th>z</th>
</tr>
</thead>
<tbody>
<tr>
<td>Time-stamp</td>
<td>null</td>
<td>1</td>
<td>3</td>
</tr>
</tbody>
</table>

- **x = a**
- **x ≠ a**
- **y = b**
- **y ≠ b**
- **z = c**
- **z ≠ c**

Assume **y ≠ b** leads to fixing **y**.
Algorithm 1: COS\((P = (X,C) : \text{CSP})\)

<table>
<thead>
<tr>
<th>Output: true iff (P) is satisfiable</th>
</tr>
</thead>
<tbody>
<tr>
<td>(P \leftarrow \phi(P))</td>
</tr>
<tr>
<td>(\text{if } P = \bot \text{ then})</td>
</tr>
<tr>
<td>(\quad\text{if lastVar }\neq \text{null then})</td>
</tr>
<tr>
<td>(\quad\quad\text{nConflicts }\leftarrow \text{nConflicts }+ 1)</td>
</tr>
<tr>
<td>(\quad\quad\text{stamp[lastVar] }\leftarrow \text{nConflicts})</td>
</tr>
<tr>
<td>(\quad\text{return false})</td>
</tr>
<tr>
<td>(\text{if } \forall x \in X,</td>
</tr>
<tr>
<td>(\quad\text{return true})</td>
</tr>
<tr>
<td>(\text{failed }\leftarrow {x \in X : \text{stamp}[x] &gt; 0 \land</td>
</tr>
<tr>
<td>(\text{if failed }= \emptyset \text{ then})</td>
</tr>
<tr>
<td>(\quad\text{lastVar }\leftarrow \text{varHeuristic.select()})</td>
</tr>
<tr>
<td>(\text{else})</td>
</tr>
<tr>
<td>(\quad\text{lastVar }\leftarrow \text{argmax}_{x \in \text{failed}}{\text{stamp}[x]})</td>
</tr>
<tr>
<td>(\quad v \leftarrow \text{valHeuristic[lastVar].select()})</td>
</tr>
<tr>
<td>(\quad\text{return } \text{COS}(P</td>
</tr>
</tbody>
</table>

*Failure: update the timestamp.*

*Failure: select the variable with the largest timestamp.*
Example

Consider the CSP below, whose infeasibility is not detected by the fixpoint algorithm because of the AllDifferent decomposition. Assume we branch in the order $x_1$, $\ldots$, $x_8$.

\[
\begin{align*}
  x_1 & \in \{0,1\} \\
  x_2 & \in \{0,1\} \\
  x_3 & \in \{0,1\} \\
  x_4 & \in \{0,1\} \\
  x_5 & \in \{0,1,2\} \\
  x_6 & \in \{0,1,2\} \\
  x_7 & \in \{0,1,2\} \\
  x_8 & \in \{0,1,2\}
\end{align*}
\]
The problem is that there is a subset of variables, \(\{x_5, x_6, x_7, x_8\}\), causing the conflict. Ideally we should detect this earlier in the search tree.
Example: Search Trees with 2 Static Orders

Search tree obtained with the static order \([x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8]\). Ideally, a heuristic learns this order.

Search tree obtained with the static order \([x_5, x_6, x_7, x_8, x_1, x_2, x_3, x_4]\): ideally, a heuristic learns this order.
Example: Search Tree with COS
Discrepancy-Based Search Combinator
Question: Assume I have a good heuristic

Then this solution should be really good.

But is this one supposed to be better than...

...this one?
Question: Assume I have a good heuristic.

Then this solution should be really good.

But is this one supposed to be better than...

...this one?

Classify solutions according to the mistakes of the heuristic!
Discrepancy-Based Search

- **Discrepancy** = the number of right decisions.

Given that:
- I trust my heuristic
- Wrong decisions generally occur at early stages

We should visit solutions by increasing discrepancy.
Objective obj = cp.minimize(totCost);

for (int dL = 0; dL < x.length; dL++) {
    DFSelf dfs = makeDfs(cp, limitedDiscrepancy(firstFail(x), dL));
    dfs.optimize(obj);
}
public class LimitedDiscrepancyBranching implements Supplier<Procedure[]> {

    private int curD;
    private final int maxD;
    private final Supplier<Procedure[]> bs;

    public LimitedDiscrepancyBranching(Supplier<Procedure[]> branching, int maxDiscrepancy) {
        if (maxDiscrepancy < 0) throw new IllegalArgumentException("max discrepancy should be >= 0");
        this.bs = branching;
        this.maxD = maxDiscrepancy;
    }

    @Override
    public Procedure[] get() {
        // TODO
        throw new NotImplementedException();
    }
}

- Eliminate alternatives that would exceed maxD
- Wrap each alternative (closure) such that the call method of the wrapped alternatives:
  - Augments curD depending on its position
  - +0 for alts[0], ..., +i for alts[i]
Search: Summary

- Variable selection using the first-fail heuristic:
  - Degree and Weighted Degree
  - Impact-Based Search (IBS):
    - Choose the variable and value that have a strong impact in terms of filtering.
    - The impacts are estimated and learned during the search.
  - Activity-Based Search (ABS):
    - Choose the variable whose domain is often filtered: activity.
    - The activity is learned during the search.
  - Last-Conflict Search (LC):
    - Choose the variable that led to the latest conflict, if any.
  - Conflict-Ordering Search (COS):
    - Choose the variable with the largest timestamp.
- Value selection heuristic:
  - Bound-Impact Value Selector (BIVS)
  - Phase saving
  - Discrepancy-Based Search Combinator