



Instructions: This is a multiple-choice exam, to save you the time of tidying up the presentation of your answers. There is exactly *one* correct answer per question. You can keep the question sheets and should *hand in only the answer sheet*: you are *not* expected to explain your answers. Unfortunately, the teacher cannot attend this exam. *Also read the instructions on the answer sheet before starting.*

0 Warm-Up Questions: General Culture

Question –2: How is Edsger W. Dijkstra’s surname pronounced in his native language?

- A dəkstra B daikstra C dikstra D dɜɪkstra E deikstra

Question –1: In which year was the Alan M. Turing Award *not* won by a scientist whose work was discussed in this course?

- A 1972 B 1974 C 1976 D 1978 E 1982

1 String Matching

Question 1: On which of the following length- m patterns P does the naïve string matching algorithm reach its *worst*-case runtime when looking for *all* occurrences of P in the text $T = 0^n$ (that is, a string of n occurrences of the character ‘0’), with $n \geq m \geq 3$?

- A 1^m B $1^{m-1}0$ C $1(0^{m-1})$ D $0(1^{m-1})$ E $0^{m-1}1$

Question 2: For the Rabin-Karp string matching algorithm, let p denote the fingerprint of the length- m pattern P , and let t_s denote the fingerprint of the length- m substring T_s for shift s in text T (of length at least m). On which assumption does the algorithm rely?

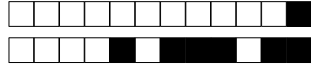
- A $p = t_s \Rightarrow P = T_s$ D $p \neq t_s \Leftarrow \exists k \in 1 \dots m : P[k] \neq T_s[k]$
 B $p = t_s \Leftrightarrow P = T_s$
 C $p = t_s \Leftarrow P = T_s$ E $p \neq t_s \Rightarrow \forall k \in 1 \dots m : P[k] \neq T_s[k]$

Question 3: How many spurious hits does the Rabin-Karp string matching algorithm encounter in the text $T = \text{“3141512653849792”}$ when looking for *all* occurrences of the pattern $P = \text{“26”}$, working modulo $q = 11$ and over the alphabet $\Sigma = \{0, 1, 2, \dots, 9\}$?

- A 0 B 1 C 2 D 3 E 4

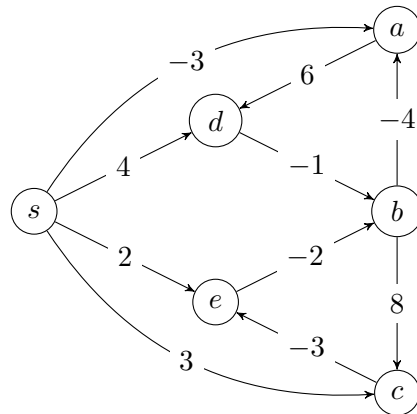
Question 4: On which of the following patterns P does the Rabin-Karp string matching algorithm reach its *best*-case runtime when looking for *all* occurrences of P in the text $T = 0^n$ (that is, a string of n occurrences of the character ‘0’), with $n \geq 3$, working modulo $q = 3$ and over the alphabet $\Sigma = \{0, 1, 2, \dots, 9\}$?

- A “660” B “300” C “099” D “007” E “000”



2 Greedy Algorithms

Consider a weighted directed graph G with a set V of n vertices, a set E of m edges, and an edge-weight function $w: E \rightarrow \mathbb{R}$, such as in the figure to the right. In this example, some edges have **negative** weights: Dijkstra's single-source shortest paths algorithm is **not** applicable for finding shortest paths (that is, minimum-weight paths) from vertex s to all vertices; for instance, a shortest path from s to d is s - c - e - b - a - d and has weight $3 + (-3) + (-2) + (-4) + 6 = 0$.



Question 5: Let G have some negative weights: which condition below is **sufficient** for the existence of a shortest path from a source vertex s to a vertex reachable from s ?

- A** if G is acyclic
- B** if G is connected
- C** if real weights
- D** if rational weights
- E** if integer weights

Question 6: What weight does Dijkstra's algorithm find for a shortest path from s to d ?

- A** it loops forever
- B** 0
- C** 3
- D** 4
- E** it crashes

Question 7: If G has some negative weights, violating the precondition of Dijkstra's algorithm, then let us first modify G by adding $-\ell$ to all weights, where $\ell < 0$ is the smallest weight, giving us a new graph G' . When does Dijkstra's algorithm find shortest paths in G' that correspond (but with exaggerated weights) 1-to-1 to shortest paths in G ?

- A** always
- B** sometimes
- C** never
- D** un-decidable
- E** we do not know

Question 8: Dijkstra's greedy algorithm for the single-source shortest paths problem performs n iterations: does it stay correct if we drop the inner loop of relaxation checks at the n th iteration of the outer loop (extraction of a minimum from a min-priority queue)?

- A** always
- B** sometimes
- C** never
- D** un-decidable
- E** we do not know

Question 9: Let G possibly have some negative weights: what is the maximum number of edges of a shortest path that is simple (that is, it does not re-visit vertices)?

- A** $-\infty$
- B** $n - 1$
- C** n
- D** $n + 1$
- E** $+\infty$



3 Dynamic Programming

Consider a weighted directed graph G with a set V of n vertices, a set E of m edges, an edge-weight function $w: E \rightarrow \mathbb{R}$ with no negative-weight cycles, and a source vertex s . Consider the following recurrence, parameterised by $\langle \alpha, \beta_1, \dots, \beta_4, \gamma \rangle$, for a quantity $T[i, v]$:

$$T[i, v] = \begin{cases} 0 & \text{if } \alpha \wedge (v = s) \\ \beta_1 & \text{if } \alpha \wedge (v \neq s) \\ \gamma \left(T[i-1, \beta_2], \min_{(x,v) \in E} (T[i-1, \beta_3] + \beta_4) \right) & \text{if } \neg \alpha \end{cases}$$

Question 10: If $T[n-1, v]$ is returned by a correct algorithm for computing the weight of a shortest path in G from s to v , then $T[i, v]$, with $0 \leq i < n$ and $v \in V$, denotes the weight of a shortest path from s to v with how many edges?

- A $< i$ B $\leq i$ C $= i$ D $\geq i$ E $> i$

Question 11: For the graph in the figure of Section 2, what is the sum $T[3, a] + T[3, b]$?

- A -8 B -6 C -5 D -4 E -3

Question 12: What is the logical condition α ?

- A $i = 0$ B $i = n - 1$ C $i = m - 1$ D $i = m$ E $i \cdot v = 0$

Question 13: What is the numeric expression β_1 ?

- A $-\infty$ B -1 C 0 D $+1$ E $+\infty$

Question 14: What is the index expression β_2 ?

- A s B $v - 1$ C v D $v + 1$ E x

Question 15: What is the index expression β_3 ?

- A s B $v - 1$ C v D $v + 1$ E x

Question 16: What is the numeric expression β_4 ?

- A $w(s, v)$ B $w(s, x)$ C $w(v, x)$ D $w(x, v)$ E 1

Question 17: What is the two-argument operator γ (written in prefix form above)?

- A $+$ B Σ C Π D w E \min

Question 18: What is an ordering of the indices i and v under which the elements $T[i, v]$ of the table T can be filled without performing any redundant computations?

- A by lexicographically increasing $\langle v, i \rangle$ D by lexicographically decreasing $\langle v, -i \rangle$
 B any order
 C by lexicographically increasing $\langle -i, v \rangle$ E by lexicographically decreasing $\langle -i, v \rangle$

We say that $\langle a, b \rangle$ is *lexicographically smaller* than $\langle c, d \rangle$ if either $a < c$ or $a = c \wedge b < d$.



4 Complexity

Consider a weighted directed graph G with a set V of n vertices, a set E of m edges, and an edge-weight function $w: E \rightarrow \mathbb{R}$. Some of the following questions refer to the recurrence on $T[i, v]$ of Section 3, but they can be answered without knowing the correct answers to the questions of that section.

Question 19: What is the *space* complexity of a dynamic program obtained directly (that is, without improvements) from the recurrence on $T[i, v]$?

- A $\Theta(n \cdot m)$ B $\Theta(n \cdot m^2)$ C $\Theta(n^2)$ D $\Theta(n^2 \cdot m)$ E $\Theta(n^3)$

Question 20: What is the *tightest time* complexity of a dynamic program for the recurrence on $T[i, v]$?

- A $\mathcal{O}(n \cdot m)$ B $\mathcal{O}(n \cdot m^2)$ C $\mathcal{O}(n^2)$ D $\mathcal{O}(n^2 \cdot m)$ E $\mathcal{O}(n^3)$

Question 21: A *useful* decision version of the shortest path problem, which is an optimisation problem, asks whether from a given source vertex to a given target vertex there exists a path whose weight, say ω , is in what relationship to a given weight W ?

- A $\omega = n - 1$ B $\omega = n$ C $\omega \leq W$ D $\omega = W$ E $\omega \geq W$

Question 22: What is the *tightest time* complexity class, as *currently* known, of a decision version of the single-source shortest paths problem, considering the existence of a dynamic program for the recurrence on $T[i, v]$?

- A P B pseudo-polynomial C NP-complete D NP-hard E we do not know

Question 23: What is the *tightest time* complexity class, as *currently* known, of a decision version of the optimisation problem of finding *longest* paths from a given source vertex to all other vertices of a weighted directed graph?

- A P B pseudo-polynomial C NP-complete D NP-hard E we do not know

