

# Chapter 7: Process Synchronization

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- Semaphores
- Classical Problems of Synchronization
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# Background

- Concurrent access to shared data may result in data inconsistency.
- Maintaining data consistency requires mechanisms to ensure the orderly execution of cooperating processes.
- Shared-memory solution to bounded-buffer problem (Chapter 4) allows at most  $n - 1$  items in buffer at the same time. A solution, where all  $N$  buffers are used is not simple.
  - ◆ Suppose that we modify the producer-consumer code by adding a variable counter, initialized to 0 and incremented each time a new item is added to the buffer

# Bounded-Buffer

- Shared data

```
#define BUFFER_SIZE 10
typedef struct {
    . . .
} item;
item buffer[BUFFER_SIZE];
int in = 0;
int out = 0;
int counter = 0;
```

# Bounded-Buffer

- Producer process

```
item nextProduced;
```

```
while (1) {  
  while (counter == BUFFER_SIZE)  
    ; /* do nothing */  
  buffer[in] = nextProduced;  
  in = (in + 1) % BUFFER_SIZE;  
  counter++;  
}
```

# Bounded-Buffer

- Consumer process

```
item nextConsumed;
```

```
while (1) {  
  while (counter == 0)  
    ; /* do nothing */  
  nextConsumed = buffer[out];  
  out = (out + 1) % BUFFER_SIZE;  
  counter--;  
}
```

# Bounded Buffer

- The statements

```
counter++;  
counter--;
```

must be performed atomically.

- Atomic operation means an operation that completes in its entirety without interruption.

# Bounded Buffer

- The statement “**count++**” may be implemented in machine language as:

```
register1 = counter  
register1 = register1 + 1  
counter = register1
```

- The statement “**count—**” may be implemented as:

```
register2 = counter  
register2 = register2 – 1  
counter = register2
```

# Bounded Buffer

- If both the producer and consumer attempt to update the buffer concurrently, the assembly language statements may get interleaved.
- Interleaving depends upon how the producer and consumer processes are scheduled.



# Bounded Buffer

- Assume **counter** is initially 5. One interleaving of statements is:

producer: **register1 = counter** (register1 = 5)

producer: **register1 = register1 + 1** (register1 = 6)

consumer: **register2 = counter** (register2 = 5)

consumer: **register2 = register2 - 1** (register2 = 4)

producer: **counter = register1** (counter = 6)

consumer: **counter = register2** (counter = 4)

- The value of **count** may be either 4 or 6, where the correct result should be 5.

# Race Condition

- **Race condition:** The situation where several processes access – and manipulate shared data concurrently. The final value of the shared data depends upon which process finishes last.
- To prevent race conditions, concurrent processes must be **synchronized**.

# The Critical-Section Problem

- $n$  processes all competing to use some shared data
- Each process has a code segment, called critical section, in which the shared data is accessed.
- Problem – ensure that when one process is executing in its critical section, no other process is allowed to execute in its critical section.

# Solution to Critical-Section Problem

1. **Mutual Exclusion.** If process  $P_i$  is executing in its critical section, then no other processes can be executing in their critical sections.
2. **Progress.** If no process is executing in its critical section and there exist some processes that wish to enter their critical section, then the selection of the processes that will enter the critical section next cannot be postponed indefinitely.
3. **Bounded Waiting.** A bound must exist on the number of times that other processes are allowed to enter their critical sections after a process has made a request to enter its critical section and before that request is granted.
  - Assume that each process executes at a nonzero speed
  - No assumption concerning relative speed of the  $n$  processes.

# Initial Attempts to Solve Problem

- Only 2 processes,  $P_0$  and  $P_1$
- General structure of process  $P_i$  (other process  $P_j$ )

do {

entry section



critical section

exit section



remainder section

} while (1);

- Processes may share some common variables to synchronize their actions.

# Algorithm 1

- Shared variables:

- ◆ **int turn;**  
initially **turn = 0**
- ◆ **turn - i**  $\wedge$   $P_i$  can enter its critical section

- Process  $P_i$

```
do {  
while (turn != i) ;  
critical section  
turn = j;  
reminder section  
} while (1);
```

- Satisfies mutual exclusion, but not progress

# Algorithm 2

- Shared variables

- ◆ **boolean flag[2];**  
initially **flag [0] = flag [1] = false.**
- ◆ **flag [i] = true**  $\wedge$   $P_i$  ready to enter its critical section

- Process  $P_i$

do {

**flag[i] := true;**

**hile (flag[j]) ;**

section

**critical**

**flag [i] = false;**

**remainder section**

**} while (1);**

- Satisfies mutual exclusion, but not progress requirement.

# Algorithm 3

- Combined shared variables of algorithms 1 and 2.

- Process  $P_i$

do {

**flag [i] := true;**

**urn = j;**

**hile (flag [j] and turn = j) ;**

  critical section

**flag [i] = false;**

  remainder section

} **while (1);**

- Meets all three requirements; solves the critical-section problem for two processes.



# Bakery Algorithm

## Critical section for n processes

- Before entering its critical section, process receives a number. Holder of the smallest number enters the critical section.
- If processes  $P_i$  and  $P_j$  receive the same number, if  $i < j$ , then  $P_i$  is served first; else  $P_j$  is served first.
- The numbering scheme always generates numbers in increasing order of enumeration; i.e., 1,2,3,3,3,3,4,5...

# Bakery Algorithm

- Notation  $<^0$  lexicographical order (ticket #, process id #)
  - ◆  $(a,b) < c,d$  if  $a < c$  or if  $a = c$  and  $b < d$
  - ◆  $\max(a_0, \dots, a_{n-1})$  is a number,  $k$ , such that  $k \geq a_i$  for  $i = 0, \dots, n - 1$

- Shared data

**boolean choosing[n];**

**int number[n];**

Data structures are initialized to **false** and **0** respectively

# Bakery Algorithm

```
do {  
  choosing[i] = true;  
  number[i] = max(number[0], number[1], ...❖ number [n - 1])+1;  
  choosing[i] = false;  
  for (j = 0; j < n; j++) {  
    while (choosing[j]) ;  
    while ((number[j] != 0) && (number[j,j] < number[i,i])) ;  
  }  
  critical section  
  number[i] = 0;  
  remainder section  
} while (1);
```

# Synchronization Hardware

- Test and modify the content of a word atomically

```
boolean TestAndSet(boolean &target) {  
    boolean rv = target;  
    target = true;  
  
    return rv;  
}
```

# Mutual Exclusion with Test-and-Set

- Shared data:

```
boolean lock = false;
```

- Process  $P_i$

```
do {  
while (TestAndSet(lock)) ;  
critical section  
lock = false;  
remainder section  
}
```

# Synchronization Hardware

- Atomically swap two variables.

```
void Swap(boolean &a, boolean &b) {  
    boolean temp = a;  
    a = b;  
    b = temp;  
}
```

# Mutual Exclusion with Swap

- Shared data (initialized to **false**):  
**boolean lock;**  
**boolean waiting[n];**
- Process  $P_i$   
do {  
**key = true;**  
**while (key == true)**  
**Swap(lock,key);**  
**critical section**  
**lock = false;**  
**remainder section**  
}

# Semaphores

- Synchronization tool that does not require busy waiting.
- Semaphore S – integer variable
- can only be accessed via two indivisible (atomic) operations

wait (S):

```
while S =< 0 do no-op;  
  S--;
```

signal (S):

```
S++;
```



# Critical Section of n Processes

- Shared data:

- **semaphore mutex;** //initially mutex = 1

- Process  $P_i$ :

```
do {  
    wait(mutex);  
    critical section  
    signal(mutex);  
    remainder section  
} while (1);
```

# Semaphore Implementation

- Define a semaphore as a record

```
typedef struct {  
    int value;  
    struct process *L;  
} semaphore;
```

- Assume two simple operations:
  - ◆ **block** suspends the process that invokes it.
  - ◆ **wakeup(P)** resumes the execution of a blocked process **P**.

# Implementation

- Semaphore operations now defined as

wait(S):

**S.value--;**

**if (S.value < 0) {**

**add this process to S.L;**

**lock;**

**}**

signal(S):

**S.value++;**

**if (S.value <= 0) {**

**remove a process P from S.L;**

**akeup(P);**

**}**

# Semaphore as a General Synchronization Tool

- Execute B in  $P_j$  only after A executed in  $P_i$
- Use semaphore flag initialized to 0
- Code:

$P_i$	$P_j$
M	M
A	wait(flag)
signal(flag)	B

# Deadlock and Starvation

- **Deadlock** – two or more processes are waiting indefinitely for an event that can be caused by only one of the waiting processes.

- Let S and Q be two semaphores initialized to 1

P0	P1
wait(S);	wait(Q);
wait(Q);	wait(S);
M	M
signal(S);	signal(Q);
signal(Q)	signal(S);

- **Starvation** – indefinite blocking. A process may never be removed from the semaphore queue in which it is suspended.

# Two Types of Semaphores

- Counting semaphore – integer value can range over an unrestricted domain.
- Binary semaphore – integer value can range only between 0 and 1; can be simpler to implement.
- Can implement a counting semaphore  $S$  as a binary semaphore.

# Implementing S as a Binary Semaphore

- Data structures:  
**binary-semaphore S1, S2;**  
**int C;**

- Initialization:  
**S1 = 1**  
**S2 = 0**  
**C = initial value of semaphore S**

# Implementing S

- wait operation

```
wait(S1);
```

```
C--;
```

```
if (C < 0) {
```

```
  signal(S1);
```

```
  wait(S2);
```

```
}
```

```
signal(S1);
```

- signal operation

```
wait(S1);
```

```
C ++;
```

```
if (C <= 0)
```

```
  signal(S2);
```

```
else
```

```
  signal(S1);
```



# Classical Problems of Synchronization

- Bounded-Buffer Problem
- Readers and Writers Problem
- Dining-Philosophers Problem

# Bounded-Buffer Problem

- Shared data

**semaphore full, empty, mutex;**

Initially:

**full = 0, empty = n, mutex = 1**

# Bounded-Buffer Problem Producer Process

```
do {  
  ...  
  produce an item in nextp  
  ...  
  wait(empty);  
  wait(mutex);  
  ...  
  add nextp to buffer  
  ...  
  signal(mutex);  
  signal(full);  
} while (1);
```

# Bounded-Buffer Problem Consumer Process

```
do {  
  wait(full)  
  wait(mutex);  
  ...  
  remove an item from buffer to nextc  
  ...  
  signal(mutex);  
  signal(empty);  
  ...  
  consume the item in nextc  
  ...  
} while (1);
```

# Readers-Writers Problem

- Shared data

**semaphore mutex, wrt;**

Initially

**mutex = 1, wrt = 1, readcount = 0**

# Readers-Writers Problem Writer Process

**wait(wrt);**

**...**

**writing is performed**

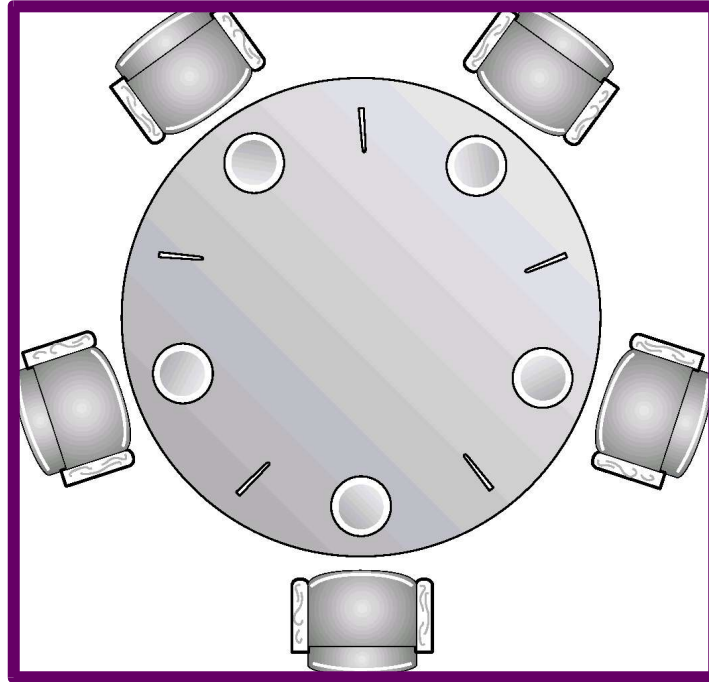
**...**

**signal(wrt);**

# Readers-Writers Problem Reader Process

```
wait(mutex);  
readcount++;  
if (readcount == 1)  
wait(rt);  
signal(mutex);  
...  
reading is performed  
...  
wait(mutex);  
readcount--;  
if (readcount == 0)  
signal(wrt);  
signal(mutex);
```

# Dining-Philosophers Problem



- Shared data  
**semaphore chopstick[5];**  
Initially all values are 1



# Dining-Philosophers Problem

```
■ Philosopher i:  
do {  
wait(chopstick[i])  
wait(chopstick[(i+1) % 5])  
...  
eat  
...  
signal(chopstick[i]);  
signal(chopstick[(i+1) % 5]);  
...  
think  
...  
} while (1);
```

# Critical Regions

- High-level synchronization construct
- A shared variable **v** of type **T**, is declared as:  
**v: shared T**
- Variable **v** accessed only inside statement  
**region v when B do S**

ere **B** is a boolean expression.

- While statement **S** is being executed, no other process can access variable **v**.

# Critical Regions

- Regions referring to the same shared variable exclude each other in time.
- When a process tries to execute the region statement, the Boolean expression B is evaluated. If B is true, statement S is executed. If it is false, the process is delayed until B becomes true and no other process is in the region associated with v.

# Example – Bounded Buffer

- Shared data:

```
struct buffer {  
    int pool[n];  
    int count, in, out;  
}
```

# Bounded Buffer Producer Process

- Producer process inserts **nextp** into the shared buffer

```
region buffer when( count < n) {  
    pool[in] = nextp;  
    in:= (in+1) % n;  
    count++;  
}
```

# Bounded Buffer Consumer Process

- Consumer process removes an item from the shared buffer and puts it in **nextc**

```
region buffer when (count > 0) {  
    nextc = pool[out];  
    out = (out+1) % n;  
    count--;
```

# Implementation region x when B do S

- Associate with the shared variable x, the following variables:

```
semaphore mutex, first-delay, second-delay;  
int first-count, second-count;
```

- Mutually exclusive access to the critical section is provided by **mutex**.
- If a process cannot enter the critical section because the Boolean expression **B** is false, it initially waits on the **first-delay** semaphore; moved to the **second-delay** semaphore before it is allowed to reevaluate B.

# Implementation

- Keep track of the number of processes waiting on **first-delay** and **second-delay**, with **first-count** and **second-count** respectively.
- The algorithm assumes a FIFO ordering in the queuing of processes for a semaphore.
- For an arbitrary queuing discipline, a more complicated implementation is required.



# Monitors

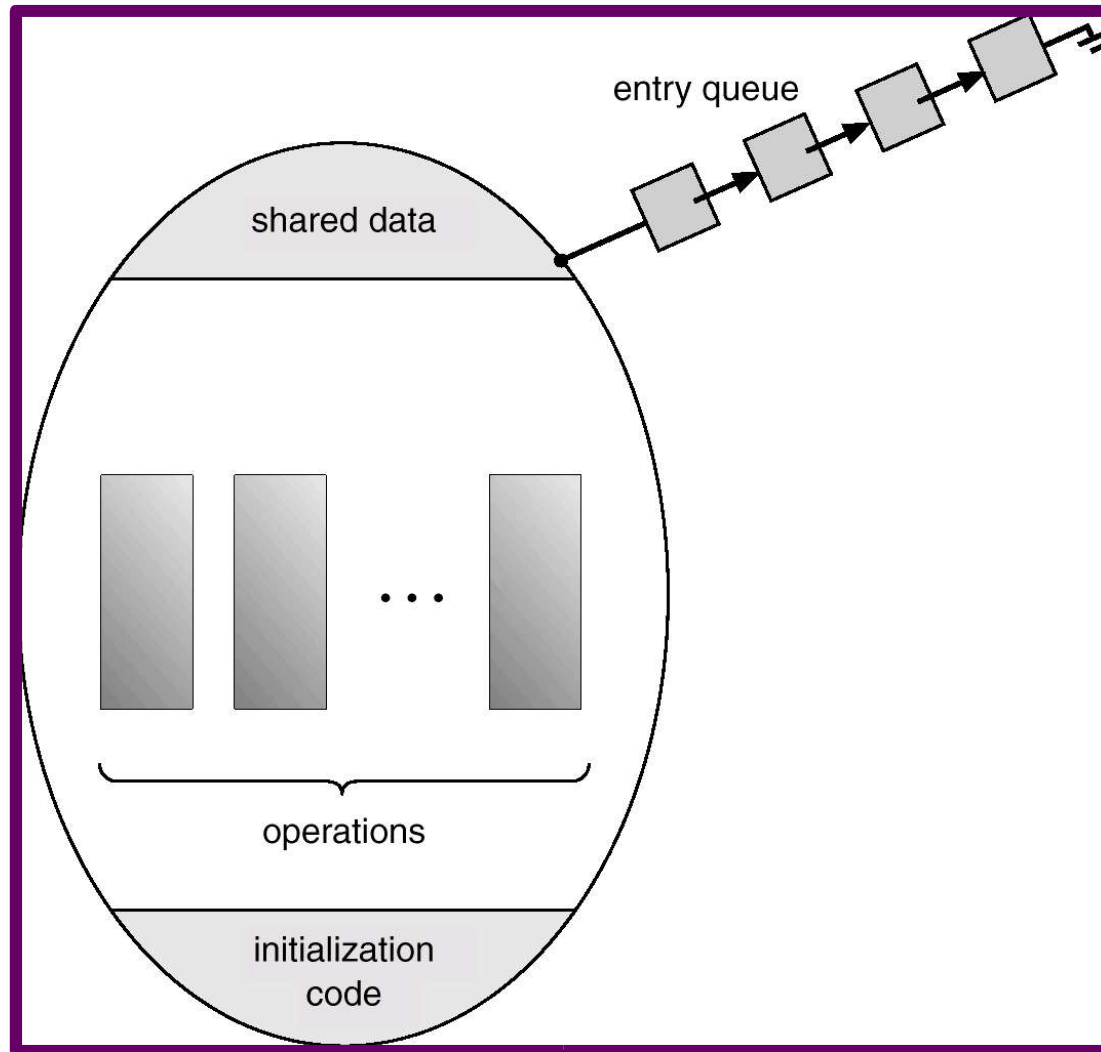
- High-level synchronization construct that allows the safe sharing of an abstract data type among concurrent processes.

```
monitor monitor-name
{
shared variable declarations
procedure body P1 (...) {
...
}
procedure body P2 (...) {
...
}
procedure body Pn (...) {
...
}
}
{
initialization code
}
}
```

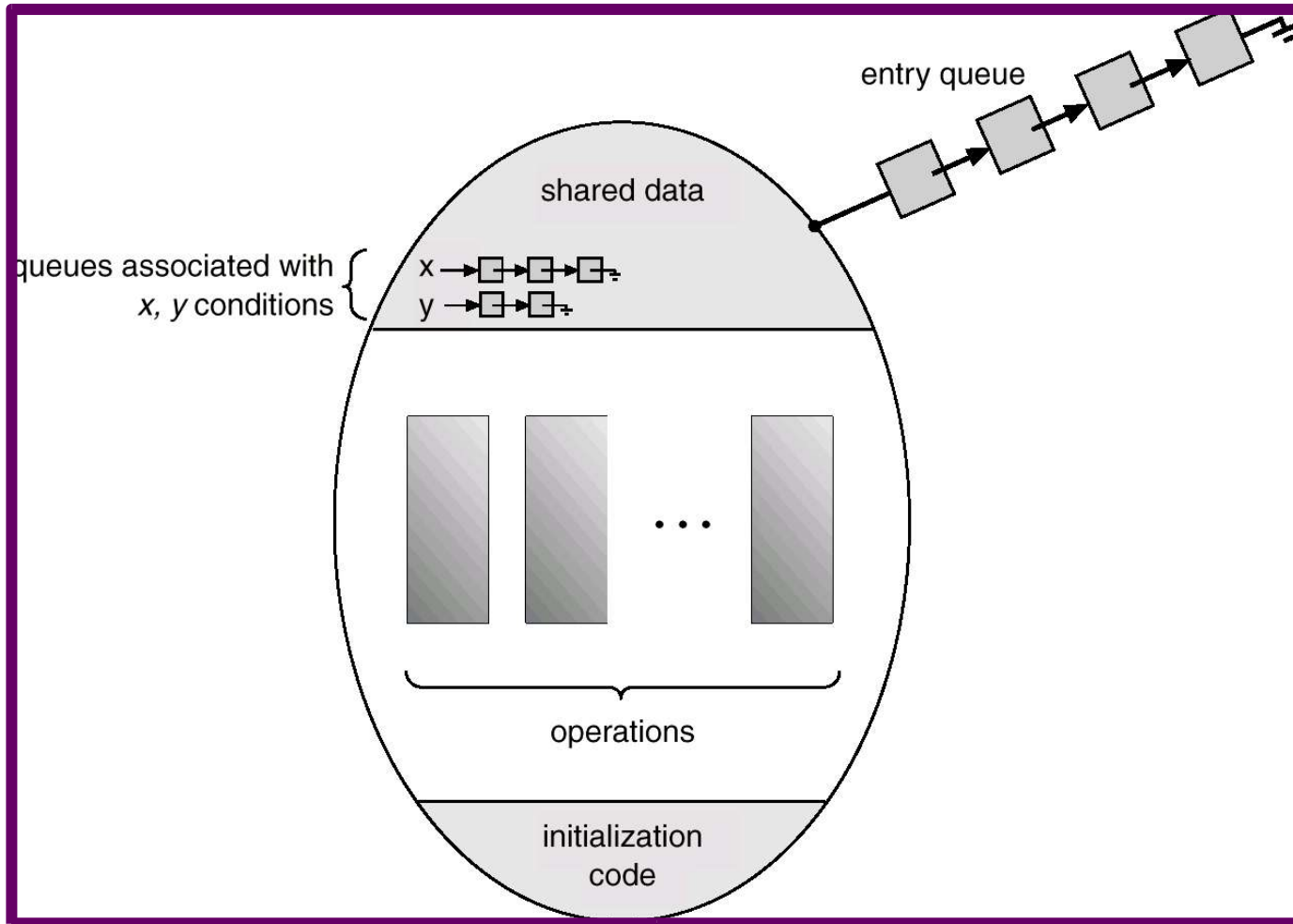
# Monitors

- To allow a process to wait within the monitor, a **condition** variable must be declared, as  
**condition x, y;**
- Condition variable can only be used with the operations **wait** and **signal**.
  - ◆ The operation **x.wait(); m** means that the process invoking this operation is suspended until another process invokes **x.signal();**
  - ◆ The **x.signal** operation resumes exactly one suspended process. If no process is suspended, then the **signal** operation has no effect.

# Schematic View of a Monitor



# Monitor With Condition Variables



# Dining Philosophers Example

```
monitor dp
{
enum {thinking, hungry, eating} state[5];
condition self[5];
void pickup(int i)           // following slides
void putdown(int i)         // following slides
void test(int i)            // following slides
void init() {
for (int i = 0; i < 5; i++)
state[i] = thinking;
}
}
```

# Dining Philosophers

```
void pickup(int i) {  
    state[i] = hungry;  
    test[i];  
    if (state[i] != eating)  
        self[i].wait();  
}
```

```
void putdown(int i) {  
    state[i] = thinking;  
    // test left and right neighbors  
    test((i+4) % 5);  
    test((i+1) % 5);  
}
```

# Dining Philosophers

```
void test(int i) {  
    if ( (state[(i + 4) % 5] != eating) &&  
        (state[i] == hungry) &&  
        (state[(i + 1) % 5] != eating)) {  
        state[i] = eating;  
        self[i].signal();  
    }  
}
```

# Monitor Implementation Using Semaphores

- Variables

```
semaphore mutex; // (initially = 1)
semaphore next;  // (initially = 0)
int next-count = 0;
```

- Each external procedure **F** will be replaced by **wait(mutex);**

```
    ...
    body of F;
```

```
    ...
if (next-count > 0)
    signal(next)
else
    signal(mutex);
```

- Mutual exclusion within a monitor is ensured.



# Monitor Implementation

- For each condition variable **x**, we have:

```
semaphore x-sem; // (initially = 0)  
int x-count = 0;
```

- The operation **x.wait** can be implemented as:

```
x-count++;  
if (next-count > 0)  
  signal(next);  
else  
  signal(mutex);  
  wait(x-sem);  
x-count--;
```

# Monitor Implementation

- The operation **x.signal** can be implemented as:

```
if (x-count > 0) {  
  next-count++;  
  signal(x-sem);  
  wait(next);  
  next-count--;  
}
```

# Monitor Implementation

- Conditional-wait construct: **x.wait(c);**
  - ◆ **c** – integer expression evaluated when the **wait** operation is executed.
  - ◆ value of **c** (a priority number) stored with the name of the process that is suspended.
  - ◆ when **x.signal** is executed, process with smallest associated priority number is resumed next.
- Check two conditions to establish correctness of system:
  - ◆ User processes must always make their calls on the monitor in a correct sequence.
  - ◆ Must ensure that an uncooperative process does not ignore the mutual-exclusion gateway provided by the monitor, and try to access the shared resource directly, without using the access protocols.

# Solaris 2 Synchronization

- Implements a variety of locks to support multitasking, multithreading (including real-time threads), and multiprocessing.
- Uses adaptive mutexes for efficiency when protecting data from short code segments.
- Uses condition variables and readers-writers locks when longer sections of code need access to data.
- Uses turnstiles to order the list of threads waiting to acquire either an adaptive mutex or reader-writer lock.

# Windows 2000 Synchronization

- Uses interrupt masks to protect access to global resources on uniprocessor systems.
- Uses spinlocks on multiprocessor systems.
- Also provides dispatcher objects which may act as wither mutexes and semaphores.
- Dispatcher objects may also provide events. An event acts much like a condition variable.

