Instructions:

- **Language of the Machine**
- **More primitive than higher level languages**
  e.g., no sophisticated control flow
- **Very restrictive**
  e.g., MIPS Arithmetic Instructions

- **We'll be working with the MIPS instruction set architecture**
  - **similar to other architectures developed since the 1980' s**
  - **used by NEC, Nintendo, Silicon Graphics, Sony**

Design goals: maximize performance and minimize cost, reduce design time

MIPS arithmetic

- All instructions have 3 operands
- Operand order is fixed (destination first)

Example:

C code: `A = B + C`

MIPS code: `add $s0, $s1, $s2`

(associated with variables by compiler)

©1998 Morgan Kaufmann Publishers

MIPS arithmetic

- Design Principle: simplicity favors regularity. Why?
- Of course this complicates some things...

C code:
```
A = B + C + D;
E = F - A;
```

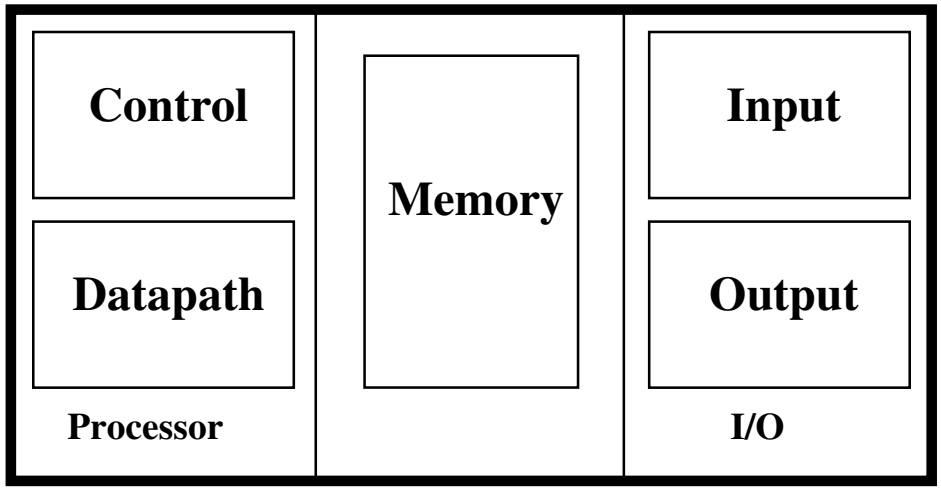
MIPS code:
```
add $t0, $s1, $s2
add $s0, $t0, $s3
sub $s4, $s5, $s0
```

- Operands must be registers, only 32 registers provided
- Design Principle: smaller is faster. Why?

©1998 Morgan Kaufmann Publishers

Registers vs. Memory

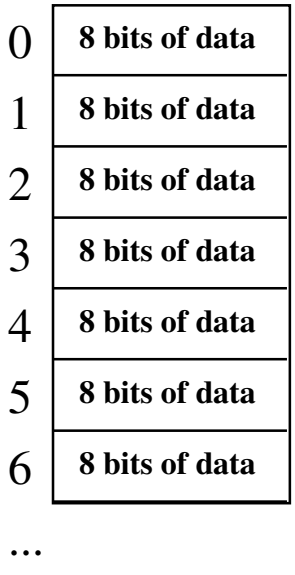
- Arithmetic instructions operands must be registers,
  — only 32 registers provided
- Compiler associates variables with registers
- What about programs with lots of variables



©1998 Morgan Kaufmann Publishers

Memory Organization

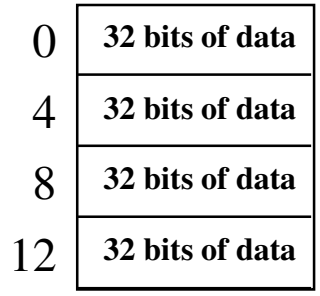
- **Viewed as a large, single-dimension array, with an address.**
- **A memory address is an index into the array**
- **"Byte addressing" means that the index points to a byte of memory.**



©1998 Morgan Kaufmann Publishers

Memory Organization

- Bytes are nice, but most data items use larger "words"
- For MIPS, a word is 32 bits or 4 bytes.



Registers hold 32 bits of data

- 2^{32} bytes with byte addresses from 0 to $2^{32}-1$
- 2^{30} words with byte addresses 0, 4, 8, ... $2^{32}-4$
- Words are aligned

  i.e., what are the least 2 significant bits of a word address?

©1998 Morgan Kaufmann Publishers

Instructions

- Load and store instructions
- Example:

C code:

```
A[8] = h + A[8];
```

MIPS code:

```
lw $t0, 32($s3)
add $t0, $s2, $t0
sw $t0, 32($s3)
```

- Store word has destination last
- Remember arithmetic operands are registers, not memory!

©1998 Morgan Kaufmann Publishers

Our First Example

- **Can we figure out the code?**

```
swap(int v[], int k);
{ int temp;
     temp = v[k]
     v[k] = v[k+1];
     v[k+1] = temp;
}
```

©1998 Morgan Kaufmann Publishers

So far we've learned:

- **MIPS**
  - **— loading words but addressing bytes**
  - **— arithmetic on registers only**

- **Instruction** **Meaning**

Instruction	Meaning
`add $s1, $s2, $s3`	`$s1 = $s2 + $s3`
`sub $s1, $s2, $s3`	`$s1 = $s2 - $s3`
`lw $s1, 100($s2)`	`$s1 = Memory[$s2+100]`
`sw $s1, 100($s2)`	`Memory[$s2+100] = $s1`

©1998 Morgan Kaufmann Publishers

Machine Language

- Instructions, like registers and words of data, are also 32 bits long
  - Example: `add $t0, $s1, $s2`
  - registers have numbers, `$t0=9, $s1=17, $s2=18`

- Instruction Format:

000000	10001	10010	01000	00000	100000

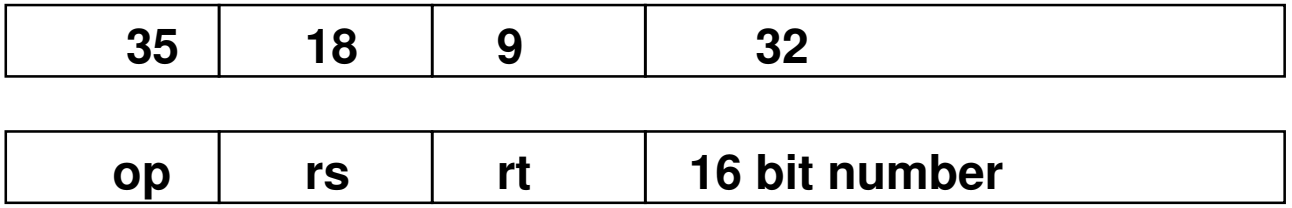
op	rs	rt	rd	shamt	funct

- *Can you guess what the field names stand for?*

©1998 Morgan Kaufmann Publishers

Machine Language

- Consider the load-word and store-word instructions,
  - What would the regularity principle have us do?
  - New principle: Good design demands a compromise
- Introduce a new type of instruction format
  - I-type for data transfer instructions
  - other format was R-type for register
- Example: `lw $t0, 32($s2)`

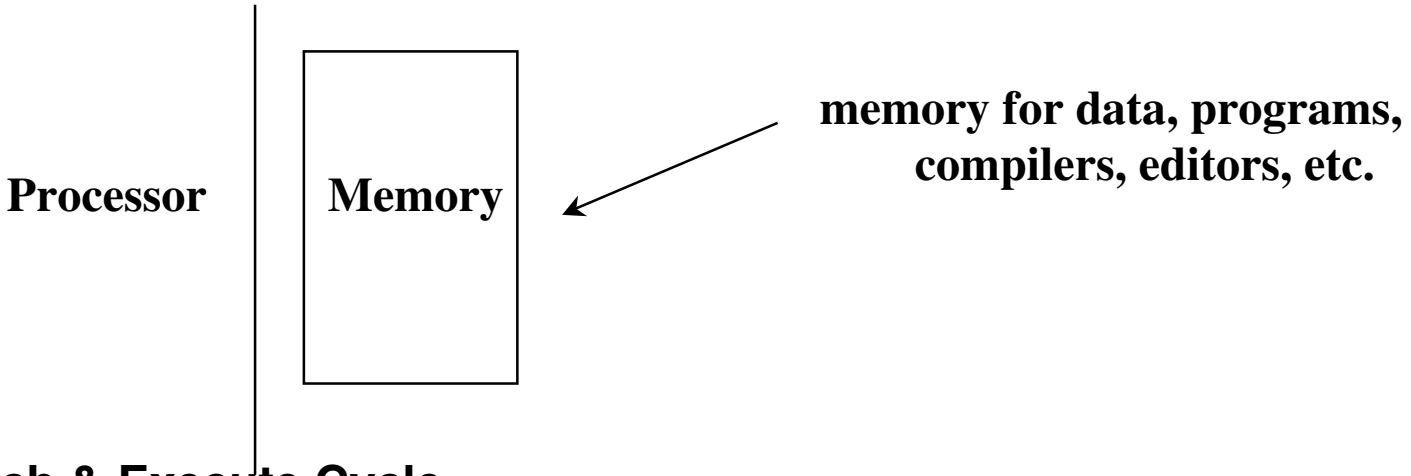


- Where's the compromise?

©1998 Morgan Kaufmann Publishers

Stored Program Concept

- Instructions are bits
- Programs are stored in memory
  — to be read or written just like data



- Fetch & Execute Cycle
  - Instructions are fetched and put into a special register
  - Bits in the register "control" the subsequent actions
  - Fetch the “next” instruction and continue

©1998 Morgan Kaufmann Publishers

Control

- Decision making instructions
  - alter the control flow,
  - i.e., change the "next" instruction to be executed

- MIPS conditional branch instructions:

```
bne $t0, $t1, Label
beq $t0, $t1, Label
```

- Example: if (i==j) h = i + j;

```
        bne $s0, $s1, Label
        add $s3, $s0, $s1
Label: ....
```

©1998 Morgan Kaufmann Publishers

Control

- **MIPS unconditional branch instructions:**

```
j  label
```

- **Example:**

```
if (i!=j)                beq $s4, $s5, Lab1
    h=i+j;               add $s3, $s4, $s5
else                     j Lab2
    h=i-j;               Lab1:  sub $s3, $s4, $s5
                         Lab2:  ...
```

- *Can you build a simple for loop?*

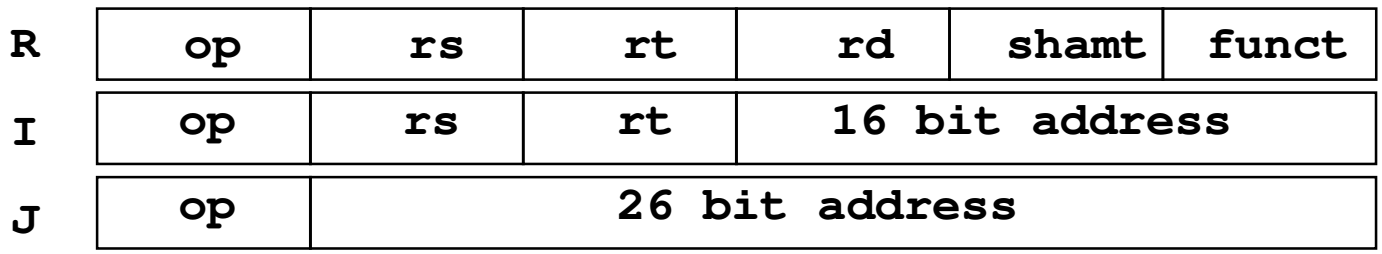
©1998 Morgan Kaufmann Publishers

So far:

- **Instruction** **Meaning**

```
add $s1,$s2,$s3   $s1 = $s2 + $s3
sub $s1,$s2,$s3   $s1 = $s2 – $s3
lw $s1,100($s2)   $s1 = Memory[$s2+100]
sw $s1,100($s2)   Memory[$s2+100] = $s1
bne $s4,$s5,L     Next instr. is at Label if $s4 ° $s5
beq $s4,$s5,L     Next instr. is at Label if $s4 = $s5
j Label           Next instr. is at Label
```

- **Formats:**



Control Flow

- We have: beq, bne, what about Branch-if-less-than?
- New instruction:

```
slt $t0, $s1, $s2
```

```
if  $s1 < $s2 then
    $t0 = 1
else
    $t0 = 0
```

- Can use this instruction to build "`blt $s1, $s2, Label`"
  — can now build general control structures
- Note that the assembler needs a register to do this,
  — there are policy of use conventions for registers

©1998 Morgan Kaufmann Publishers

Policy of Use Conventions

Name	Register number	Usage
`$zero`	0	the constant value 0
`$v0-$v1`	2-3	values for results and expression evaluation
`$a0-$a3`	4-7	arguments
`$t0-$t7`	8-15	temporaries
`$s0-$s7`	16-23	saved
`$t8-$t9`	24-25	more temporaries
`$gp`	28	global pointer
`$sp`	29	stack pointer
`$fp`	30	frame pointer
`$ra`	31	return address

©1998 Morgan Kaufmann Publishers

Constants

- Small constants are used quite frequently (50% of operands)

  e.g., A = A + 5;
  B = B + 1;
  C = C - 18;

- Solutions? Why not?
  - put 'typical constants' in memory and load them.
  - create hard-wired registers (like $zero) for constants like one.

- MIPS Instructions:

```
addi $29, $29, 4
slti $8, $18, 10
andi $29, $29, 6
ori $29, $29, 4
```

- How do we make this work?

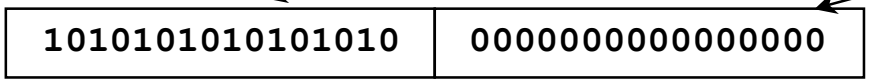
©1998 Morgan Kaufmann Publishers

How about larger constants?

- We'd like to be able to load a 32 bit constant into a register
- Must use two instructions, new "load upper immediate" instruction

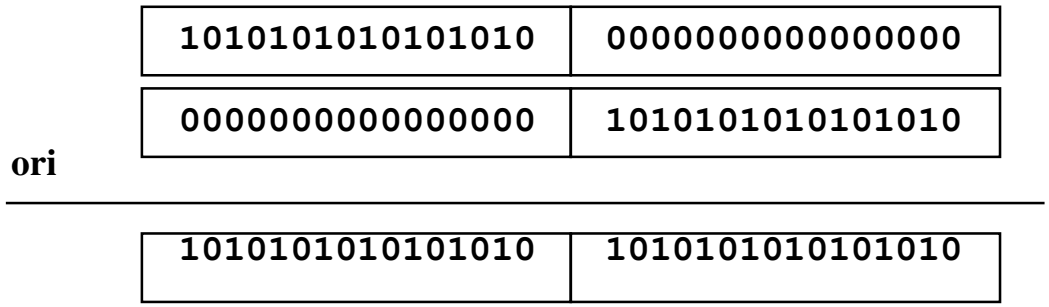
```
lui $t0, 1010101010101010
```

filled with zeros



- Then must get the lower order bits right, i.e.,

```
ori $t0, $t0, 1010101010101010
```



©1998 Morgan Kaufmann Publishers

Assembly Language vs. Machine Language

- Assembly provides convenient symbolic representation
  - much easier than writing down numbers
  - e.g., destination first
- Machine language is the underlying reality
  - e.g., destination is no longer first
- Assembly can provide 'pseudoinstructions'
  - e.g., “move $t0, $t1” exists only in Assembly
  - would be implemented using “add $t0,$t1,$zero”
- When considering performance you should count real instructions

©1998 Morgan Kaufmann Publishers

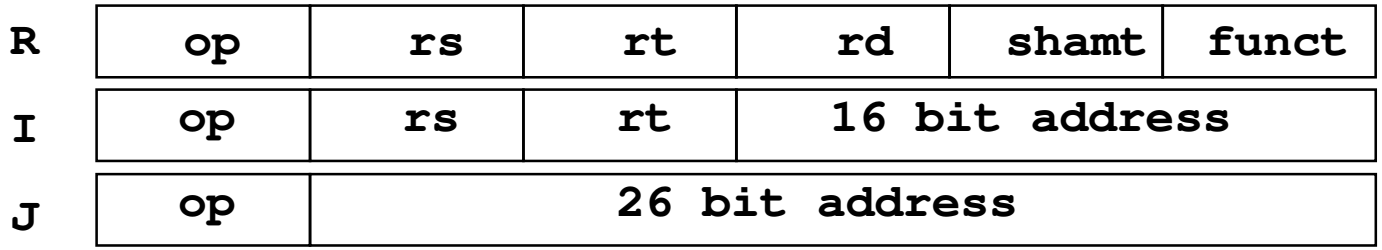
Other Issues

- Things we are not going to cover
  - support for procedures
  - linkers, loaders, memory layout
  - stacks, frames, recursion
  - manipulating strings and pointers
  - interrupts and exceptions
  - system calls and conventions
- Some of these we'll talk about later
- We've focused on architectural issues
  - basics of MIPS assembly language and machine code
  - we'll build a processor to execute these instructions.

©1998 Morgan Kaufmann Publishers

Overview of MIPS

- simple instructions all 32 bits wide
- very structured, no unnecessary baggage
- only three instruction formats



- rely on compiler to achieve performance
  — what are the compiler's goals?
- help compiler where we can

©1998 Morgan Kaufmann Publishers

Addresses in Branches and Jumps

- **Instructions:**

`bne $t4,$t5,Label`	Next instruction is at Label if `$t4 ° $t5`
`beq $t4,$t5,Label`	Next instruction is at Label if `$t4 = $t5`
`j Label`	Next instruction is at Label

- **Formats:**

- **Addresses are not 32 bits**
  - **— How do we handle this with load and store instructions?**

©1998 Morgan Kaufmann Publishers

Addresses in Branches

- **Instructions:**

`bne $t4,$t5,Label`	Next instruction is at Label if $t4°$t5
`beq $t4,$t5,Label`	Next instruction is at Label if $t4=$t5

- **Formats:**

- **Could specify a register (like lw and sw) and add it to address**
  - **use Instruction Address Register (PC = program counter)**
  - **most branches are local (principle of locality)**
- **Jump instructions just use high order bits of PC**
  - **address boundaries of 256 MB**

©1998 Morgan Kaufmann Publishers

To summarize:

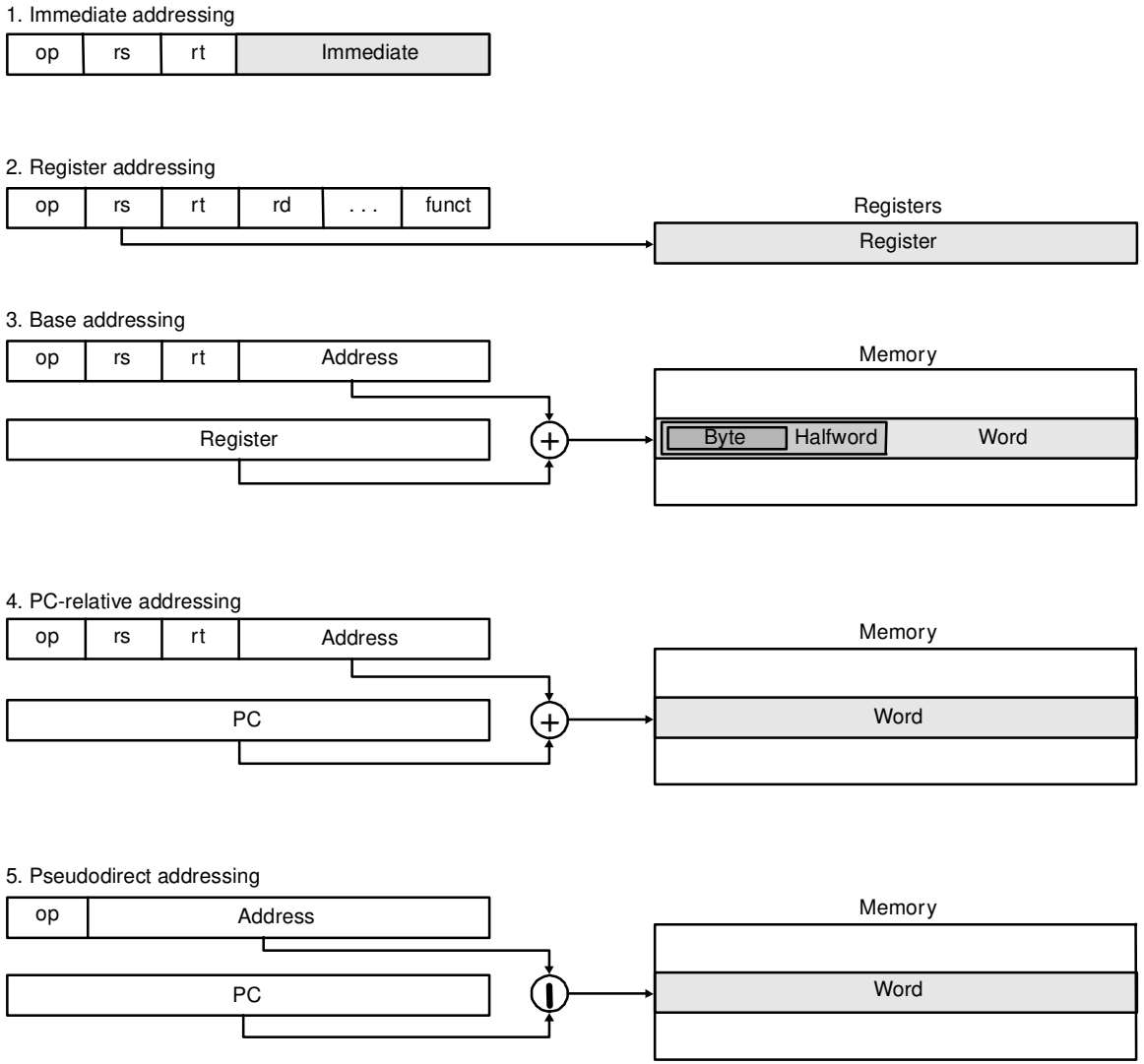
MIPS operands

Name	Example	Comments
32 registers	`$s0-$s7, $t0-$t9, $zero, $a0-$a3, $v0-$v1, $gp, $fp, $sp, $ra, $at`	Fast locations for data. In MIPS, data must be in registers to perform arithmetic. MIPS register $zero always equals 0. Register $at is reserved for the assembler to handle large constants.
2^{30} memory words	Memory[0], Memory[4], ..., Memory[4294967292]	Accessed only by data transfer instructions. MIPS uses byte addresses, so sequential words differ by 4. Memory holds data structures, such as arrays, and spilled registers, such as those saved on procedure calls.

MIPS assembly language

Category	Instruction	Example	Meaning	Comments
Arithmetic	add	`add $s1, $s2, $s3`	`$s1 = $s2 + $s3`	Three operands; data in registers
	subtract	`sub $s1, $s2, $s3`	`$s1 = $s2 - $s3`	Three operands; data in registers
	add immediate	`addi $s1, $s2, 100`	`$s1 = $s2 + 100`	Used to add constants
Data transfer	load word	`lw $s1, 100($s2)`	`$s1` = Memory[`$s2` + 100]	Word from memory to register
	store word	`sw $s1, 100($s2)`	Memory[`$s2` + 100] = $s1	Word from register to memory
	load byte	`lb $s1, 100($s2)`	`$s1` = Memory[`$s2` + 100]	Byte from memory to register
	store byte	`sb $s1, 100($s2)`	Memory[`$s2` + 100] = $s1	Byte from register to memory
	load upper immediate	`lui $s1, 100`	$s1 = 100 * 2^{16}	Loads constant in upper 16 bits
Conditional branch	branch on equal	`beq $s1, $s2, 25`	if (`$s1 == $s2`) go to PC + 4 + 100	Equal test; PC-relative branch
	branch on not equal	`bne $s1, $s2, 25`	if (`$s1 != $s2`) go to PC + 4 + 100	Not equal test; PC-relative
	set on less than	`slt $s1, $s2, $s3`	if (`$s2 < $s3`) `$s1` = 1; else `$s1` = 0	Compare less than; for beq, bne
	set less than immediate	`slti $s1, $s2, 100`	if (`$s2 < 100`) `$s1` = 1; else `$s1` = 0	Compare less than constant
Uncondi-tional jump	jump	`j 2500`	go to 10000	Jump to target address
	jump register	`jr $ra`	go to `$ra`	For switch, procedure return
	jump and link	`jal 2500`	`$ra` = PC + 4; go to 10000	For procedure call

©1998 Morgan Kaufmann Publishers



©1998 Morgan Kaufmann Publishers

Alternative Architectures

- **Design alternative:**
  - **provide more powerful operations**
  - **goal is to reduce number of instructions executed**
  - **danger is a slower cycle time and/or a higher CPI**
- **Sometimes referred to as "RISC vs. CISC"**
  - **virtually all new instruction sets since 1982 have been RISC**
  - **VAX: minimize code size, make assembly language easy**

    instructions from 1 to 54 bytes long!
- **We'll look at PowerPC and 80x86**

©1998 Morgan Kaufmann Publishers

PowerPC

- Indexed addressing
  - example: `lw $t1,$a0+$s3  #$t1=Memory[$a0+$s3]`
  - What do we have to do in MIPS?

- Update addressing
  - update a register as part of load (for marching through arrays)
  - example: `lwu $t0,4($s3) #$t0=Memory[$s3+4];$s3=$s3+4`
  - What do we have to do in MIPS?
- Others:
  - load multiple/store multiple
  - a special counter register "`bc Loop`"

    decrement counter, if not 0 goto loop

80x86

- 1978: The Intel 8086 is announced (16 bit architecture)
- 1980: The 8087 floating point coprocessor is added
- 1982: The 80286 increases address space to 24 bits, +instructions
- 1985: The 80386 extends to 32 bits, new addressing modes
- 1989-1995: The 80486, Pentium, Pentium Pro add a few instructions (mostly designed for higher performance)
- 1997: MMX is added

"This history illustrates the impact of the "golden handcuffs" of compatibility

"adding new features as someone might add clothing to a packed bag"

"an architecture that is difficult to explain and impossible to love"

A dominant architecture: 80x86

- See your textbook for a more detailed description
- Complexity:
  - Instructions from 1 to 17 bytes long
  - one operand must act as both a source and destination
  - one operand can come from memory
  - complex addressing modes
    e.g., “base or scaled index with 8 or 32 bit displacement”
- Saving grace:
  - the most frequently used instructions are not too difficult to build
  - compilers avoid the portions of the architecture that are slow

*“what the 80x86 lacks in style is made up in quantity,
making it beautiful from the right perspective”*

©1998 Morgan Kaufmann Publishers

Summary

- Instruction complexity is only one variable
  - lower instruction count vs. higher CPI / lower clock rate
- Design Principles:
  - simplicity favors regularity
  - smaller is faster
  - good design demands compromise
  - make the common case fast
- Instruction set architecture
  - a very important abstraction indeed!