

MPI

Message Passing Interface

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Previous ...

- Communicators
- Send and Receive: blocking/non-blocking
- Other Point-to-Point Functions
- Global Functions
- Datatypes
- Topology
- Timing



Outline

- One-sided communication
- Dynamic process creation
- Parallel I/O
- Example



One-sided communication

■ What is two-sided communication?

- The data movement has to be specified on both sides.
- P2P functions: Send/Recv
- Global functions: Bcast, Scatter, ...

■ One-sided?

- A process access another process address space without any explicit participation in that communication operation by the remote process.



One-sided communication

■ Advantage:

- Direct remote memory access (RMA)
- No hand-shaking
- Flexible and dynamic data distribution
- No extra buffer

■ How?

- Create accessing window
- Define the pattern
- Sync (fence)
- Reduction
- Lighter syncs



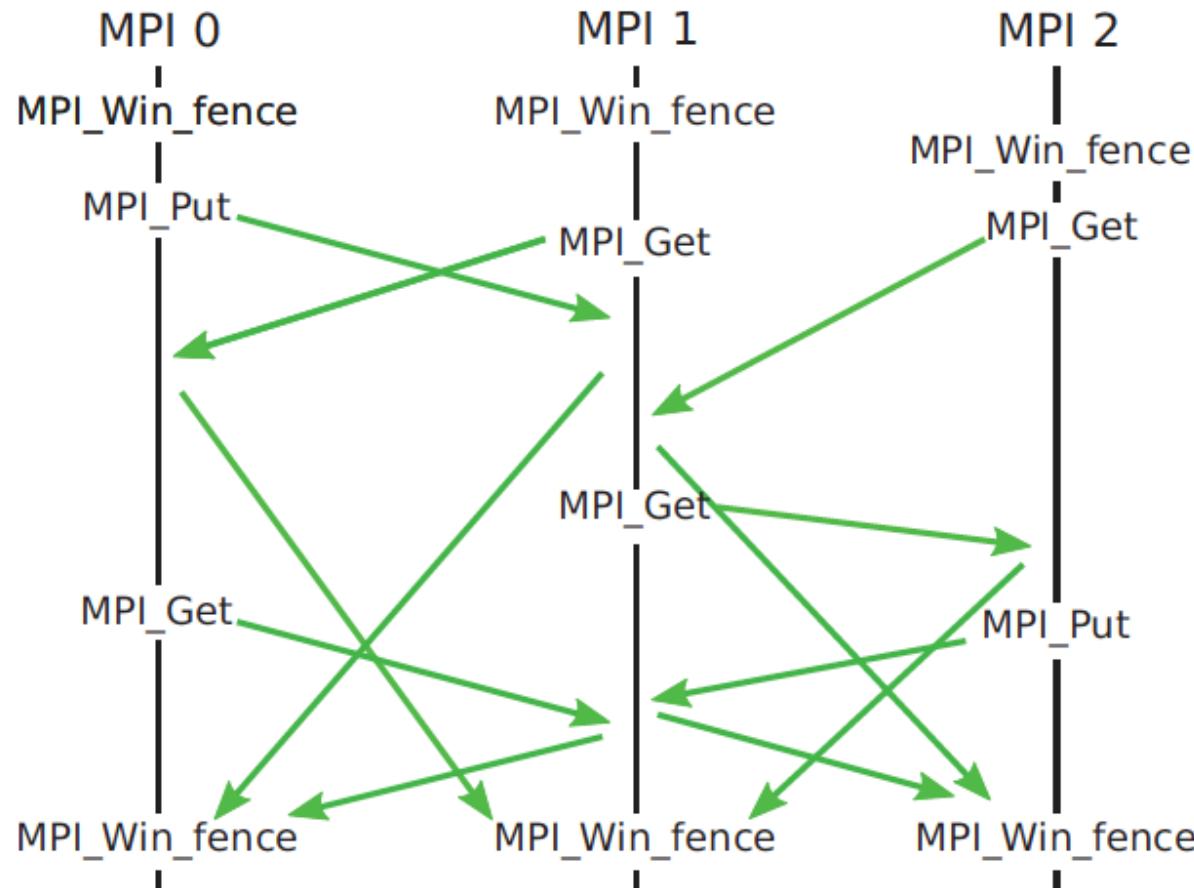
One-sided communication

■ Synchronization

- ✿ MPI_Win_fence
 - Similar to MPI_Barrier call
 - All communication calls should be inside a pair of MPI_Win_fence
- ✿ Post-Start-Complete-Wait
 - Synchronize within a group
- ✿ Lock-Unlock
 - like mutexes or other concurrent operations



MPI_Win_fence

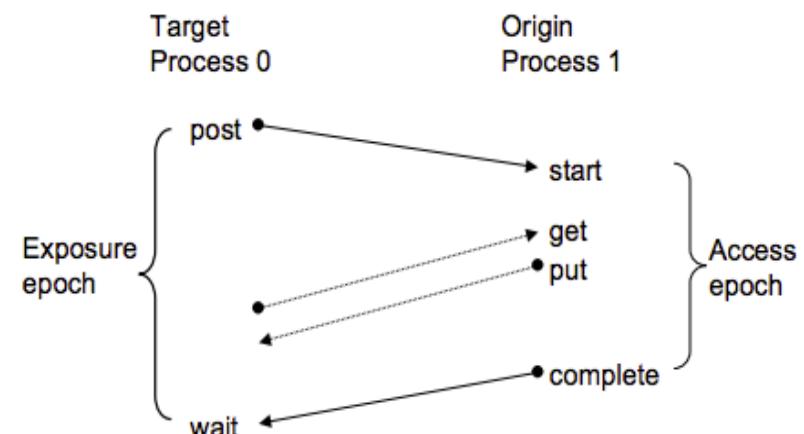


Demo:
Matrix
transpose
and
accumulate

Demo_fence

Post-Start-Complete-Wait

- Work within a group (local and lighter than MPI_Fence)
 - * Origin:
 - MPI_Win_start
 - MPI_Win_complete
 - * Target:
 - MPI_Win_post
 - MPI_Win_wait
 - Demo: [Demo_put](#)





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Dynamic Process Creation

- A MPI process can spawn new MPI processes at run time which starts running a new program.
 - ★ `MPI_Comm_spawn(...);`
 - ★ `MPI_Comm_spawn_multiple(...)`

Demo



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Parallel I/O

The I/O Challenge:

- Problems are increasingly computationally challenging
 - Large parallel machines needed to perform calculations
 - Critical to leverage parallelism in all phases
- Data access is a huge challenge
 - Using parallelism to obtain performance
 - Finding usable, efficient, portable interfaces
 - Understanding and tuning I/O
- Data stored in a single simulation for some projects: –
 $O(100)$ TB !!



Parallel I/O

- I/O approach
 - ✳ Gather all data to one process
 - ✳ Each process write to one file
 - ✳ All process write to one file
- MPI-I/O: the Basics
 - ✳ MPI-IO provides a low-level
 - ✳ Simply compile and link as you in any normal MPI program.



- MPI-IO can be done in 2 basic ways :
 - ✿ Independent – For independent I/O each MPI task is handling the I/O independently using non- collective calls like `MPI_File_write()` and `MPI_File_read()`.
 - ✿ Collective – When doing collective I/O all MPI tasks participating in I/O has to call the same routines. Basic routines are `MPI_File_write_all()` and `MPI_File_read_all()`

Demo



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Case study

■ The number of primes

- ★ A prime number is a natural number greater than 1 that has no positive divisors other than 1 and itself.
- ★ $n \% j \neq 0, j = [2,3,\dots,n-1]$
- ★ MPI_Bcast
- ★ MPI_Reduce

Demo