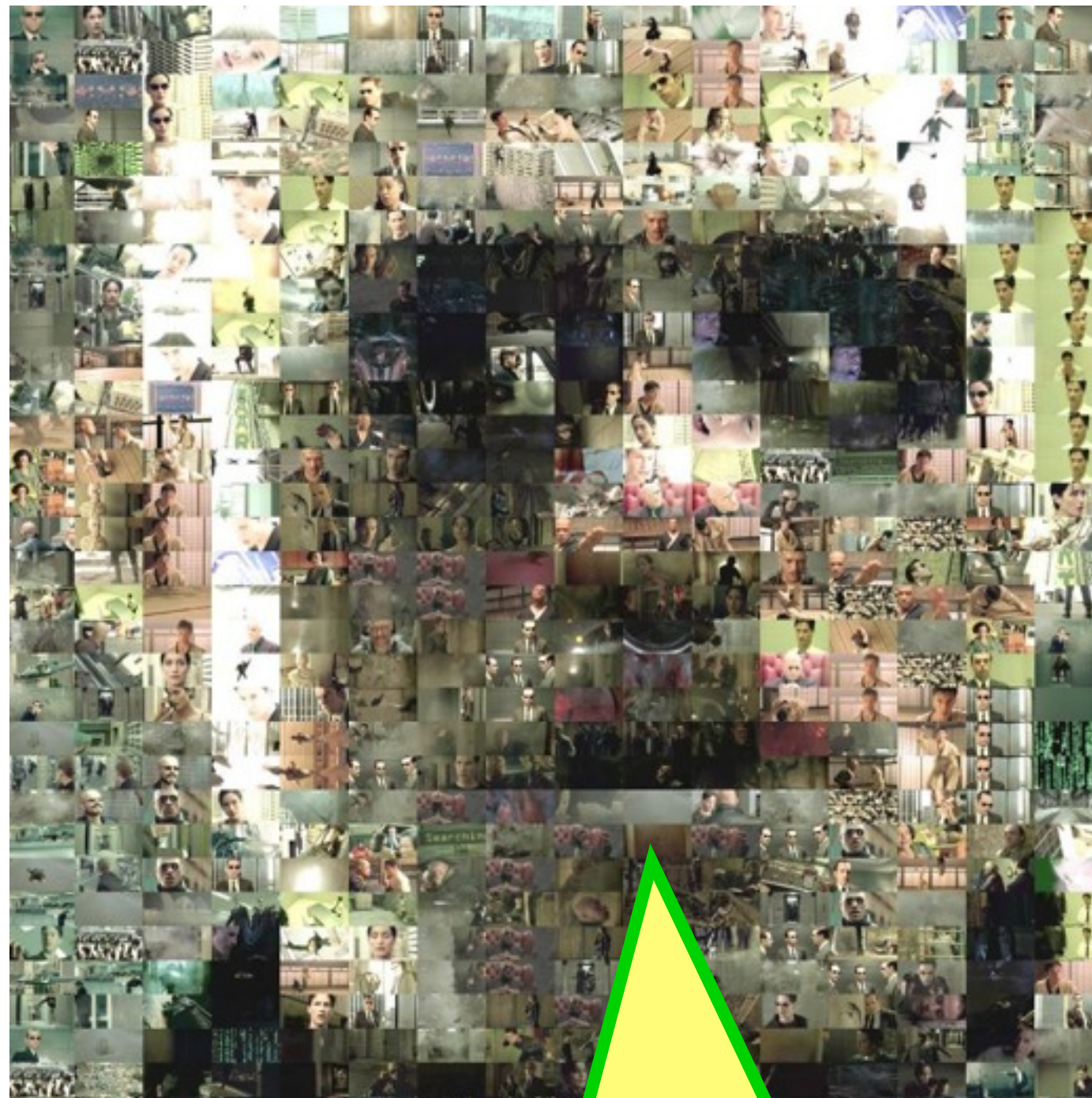


MIPS and SPIM tutorial

Part Three: Subroutines, jal, jr and the stack (push & pop)

November 2009

karl.marklund@it.uu.se



Get ready for part three of your MIPS assembly programming training.



A piece of advice...

...when you program, try to identify pieces that can be reused.

Ok, so what if we want to "invent" an addm instruction...

Calculate the sum $a + b$ and store the result to memory location m .

addm m, a, b

But no such instruction exist



Good.... Goood
thinking young man.

This (addm) is something I
would like to do over and
over again for different
values of a, b and m.

I see... You are now ready to learn
how to strucure your programs
using **subroutines**.

Calculate the sum $a + b$ and store the result to memory location m.

addm m, a, b

But no such such instruction exist

Start by describing
your **subroutine**:

```
#####  
#  
# DESCRIPTION: Calculate the sum  $S = A + B$  and store the result  
#               to memory location M.  
#  
# INPUT: $a0 - M (address)  
#        $a1 - A (integer)  
#        $a2 - B (integer)  
#  
# OUTPUT: $v0 - S (integer)  
#  
# SIDE EFFECTS: The the sum S is written to the memory location M.  
#  
#####
```

Input is allways in \$a-
registers

Output is allways in a
\$v0-register

Any side effects must
be stated.

Introduce *abstractions*

Don't use register names in your comments.

```
#####  
#  
# DESCRIPTION: Calculate the sum  $S = A + B$  and store the result  
# to memory location M.  
#  
# INPUT: $a0 - M (address)  
# $a1 - A (integer)  
# $a2 - B (integer)  
#  
# OUTPUT: $v0 - S (integer)  
#  
# SIDE EFFECTS: The the sum S is written to the memory location M.  
#  
#####
```

Refer to the previously introduced abstractions in your comments.

Any *side effects* must be stated.

```
#####
#
# DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#              to memory location M.
#
# INPUT: $a0 - M (address)
#        $a1 - A (integer)
#        $a2 - B (integer)
#
# OUTPUT: $v0 - S (integer)
#
# SIDE EFFECTS: The the sum S is written to the memory location M.
#####
```

addm:

```
add    $v0, $a1, $a2
sw     $v0, 0($a0)
```

```
jr     $ra
```

```
# S = A + B
# Store S to memory at address M
```

```
# Return to caller
```

Now you can write the code to implement the desired functionality.

Refer to the previously introduced abstractions in your comments.

Give the subroutine a name – a **label**

Always end a subroutine with jr \$ra (jump register \$ra)

```
#####
#
#   DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#               to memory location M.
#
#       INPUT:  $a0 - M (address)
#              $a1 - A (integer)
#              $a2 - B (integer)
#
#       OUTPUT: $v0 - S (integer)
#
#   SIDE EFFECTS: The the sum S is written to the memory location M.
#
#####
```



Writing *good comments* and *good looking code* does matter!

addm:

```
add    $v0, $a1, $a2
```

```
sw     $v0, 0($a0)
```

```
jr     $ra
```

```
# S = A + B
```

```
# Store S to memory at address M
```

```
# Return to caller
```

Use *tabs* to align your code.

Align your comments.


```

.data

X:    .word 0x11111111
Y:    .word 0x22222222

.text

        .globl main

main:   la      $a0, X           # Address of X
        li      $a1, 127        # a = 127
        li      $a2, 0xa        # b = 0xa = 10 (decimal)

        jal     addm            # store s = a + b in memory at
                                # address X

        la      $a0, Y           # Address of Y
        add     $a1, $v0,$zero   # a = s
        li      $a2, -3         # b = -3
        jal     addm            # store a + b in memory at
                                # address Y

        nop

        jr      $ra

```

```

.data
X:    .word 0x11111111
Y:    .word 0x22222222

.text
      .globl main

```

```

main:  la      $a0, X           # Address of X
      li      $a1, 127         # a = 127
      li      $a2, 0xa        # b = 0xa = 10 (decimal)

```

```
      addm
```

Load Immediate (li)

Set the content of register \$a2 to $10_{10} = a_{16} = 0xa$ using an immediate constant.

To test the subroutine we store these values in the data segment

Introduce and reuse abstractions in your comments.

Can use decimal or hexadecimal notation.

```

      # store a + b in memory at address Y
      # a = 127
      # b = -3
      store a + b in memory at address Y

```

```
      nop

```

```
      jr      $ra

```



**When you load
subroutines.s**

Registers

R0 (r0)	= 00000000	R13 (t5)	= 00000000	R20 (s4)	= 00000000	R27 (k1)	= 00000000
R1 (at)	= 00000000	R14 (t6)	= 00000000	R21 (s5)	= 00000000	R28 (gp)	= 10008000
R2 (v0)	= 00000000	R15 (t7)	= 00000000	R22 (s6)	= 00000000	R29 (sp)	= 7ffffeffc
R3 (v1)	= 00000000	R16 (t8)	= 00000000	R23 (s7)	= 00000000		
R4 (a0)	= 00000000	R17 (t9)	= 00000000	R24 (t8)	= 00000000		
R5 (a1)	= 00000000	R18 (t10)	= 00000000	R25 (t9)	= 00000000		
R6 (a2)	= 00000000	R19 (t11)	= 00000000	R26 (k0)	= 00000000		

Remember: X and Y are labels, labels are nothing but named memory locations (addresses).

At memory address X (0x10010000) the value 0x11111111 is stored.

At memory address Y (0x10010004) the value 0x22222222 is stored.

DATA

[0x10000000]...	[0x10010000]	0x00000000	0x00000000
[0x10010000]		0x11111111	0x22222222
[0x10010010]...	[0x10040000]	0x00000000	0x00000000

STACK

[0x7ffffeffc]	0x00000000
---------------	------------

**"X" and "Y" are stored in the
data segment.**

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All Rights Reserved.

DOS and Windows ports by David ...

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See the file README for a full copyright notice.

Loaded: C:\Program Files\PCSpim\exceptions.s

C:\Documents and Settings\Karl Marklund\My Documents\Teaching\Dark ht 2008\Tutorials By Karl Marklund\subr

.data

X: **.word 0x11111111**
Y: **.word 0x22222222**

.text
.globl main

main: **la** **\$a0, X** **# Address of X**
 li **\$a1, 127** **# a = 127**
 li **\$a2, 0xa** **# b = 0xa = 10 (decimal)**

jal **addm** **# store s = a + b in memory at address X**

\$ra → **la** **\$a0, Y** **# Address of Y**
 add **\$a1, \$v0,** **Jump and Link to**
 li **\$a2, -3** **the subroutine.**
 addm **memory address Y**

Jump and Link stores the address to the instruction following the jal in \$ra.

The execution will now continue in the subroutine...

Prepare the call to the subroutine by providing input in the \$a-registers.

```
#####
#
#   DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#               to memory location M.
#
#       INPUT: $a0 - M (address)
#              $a1 - A (integer)
#              $a2 - B (integer)
#
#       OUTPUT: $v0 - S (integer)
#
#   SIDE EFFECTS: The the sum S is written to the memory location M.
#
#####
addm:
PC → add    $v0, $a1, $a2      #  $S = A + B$ 
    sw      $v0, 0($a0)      # Store S to memory at address M

    jr      $ra              # Return to caller
```

```
#####
#
#   DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#               to memory location M.
#
#       INPUT: $a0 - M (address)
#              $a1 - A (integer)
#              $a2 - B (integer)
#
#       OUTPUT: $v0 - S (integer)
#
#   SIDE EFFECTS: The the sum S is written to the memory location M.
#
#####
addm:
    add    $v0, $a1, $a2        #  $S = A + B$ 
    sw     $v0, 0($a0)         # Store S to memory at address M
    jr     $ra                 # Return to caller
```

PC →


```
#####
#
#  DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#                to memory location M.
#
#      INPUT: $a0 - M (address)
#             $a1 - A (integer)
#             $a2 - B (integer)
#
#      OUTPUT: $v0 - S (integer)
#
#  SIDE EFFECTS: The the sum S is written to the memory location M.
#
#####
addm:
    add    $v0, $a1, $a2        #  $S = A + B$ 
    sw     $v0, 0($a0)          # Store S to memory at address M

PC → jr     $ra                # Return to caller
```

We can return to the caller using the stored return address in \$ra

The sum S is *returned* in the output register \$v0.

.data

X: .word 0x11111111

Y: .word 0x22222222

.text

.globl main

main:

la	\$a0, X	# Address of X
li	\$a1, 127	# a = 127
li	\$a2, 0xa	# b = 0xa = 10 (decimal)

jal	addm	# store s = a + b in memory at address X
------------	-------------	---

\$ra →

la	\$a0, Y	# Address of Y
add	\$a1, \$v0,\$zero	# a = S
li	\$a2, -3	# b = -3
jal	addm	# store a + b in memory at address Y

nop

jr \$ra

.data

X: .word 0x11111111

Y: .word 0x22222222

.text

.globl main

main:

la	\$a0, X	# Address of X
li	\$a1, 127	# a = 127
li	\$a2, 0xa	# b = 0xa = 10 (decimal)

jal	addm	# store s = a + b in memory at address X
------------	-------------	---

PC →

la	\$a0, Y	# Address of Y
add	\$a1, \$v0,\$zero	# a = S
li	\$a2, -3	# b = -3
jal	addm	# store a + b in memory at address Y

nop

jr \$ra

.data

X: .word 0x11111111

Y: .word 0x22222222

.text

.globl main

main: la \$a0, X # Address of X
li \$a1, 127 # a = 127
li \$a2, 0xa # b = 0xa = 10 (decimal)
jal addm # store s = a + b in memory

la \$a0, Y # Address of Y
add \$a1, \$v0,\$zero # a = S
li \$a2, -3 # b = -3

jal addm # store a + b in memory at address Y

\$ra → nop

jr \$ra

**Prepare a 2nd call
to the subroutine
by providing input
in the \$a-registers.**

**Jump and Link stores the
address to the instruction
following the jal in \$ra.**

```
#####
#
#   DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#               to memory location M.
#
#       INPUT: $a0 - M (address)
#              $a1 - A (integer)
#              $a2 - B (integer)
#
#       OUTPUT: $v0 - S (integer)
#
#   SIDE EFFECTS: The the sum S is written to the memory location M.
#
#####
addm:
PC → add    $v0, $a1, $a2      # S = A + B
      sw     $v0, 0($a0)      # Store S to memory at address M

      jr     $ra              # Return to caller
```

```
#####
#
#   DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#               to memory location M.
#
#       INPUT:  $a0 - M (address)
#              $a1 - A (integer)
#              $a2 - B (integer)
#
#       OUTPUT: $v0 - S (integer)
#
#   SIDE EFFECTS: The the sum S is written to the memory location M.
#
#####
addm:
    add    $v0, $a1, $a2        #  $S = A + B$ 
PC → sw    $v0, 0($a0)         # Store S to memory at address M

    jr     $ra                 # Return to caller
```



```
#####
#
#  DESCRIPTION: Calculate the sum  $S = A + B$  and store the result
#                to memory location M.
#
#      INPUT: $a0 - M (address)
#             $a1 - A (integer)
#             $a2 - B (integer)
#
#      OUTPUT: $v0 - S (integer)
#
#  SIDE EFFECTS: The the sum S is written to the memory location M.
#
#####
addm:
    add    $v0, $a1, $a2        #  $S = A + B$ 
    sw     $v0, 0($a0)          # Store S to memory at address M

PC → jr     $ra                # Return to caller
```

.data

X: .word 0x11111111
Y: .word 0x22222222

.text
.globl main

main: la \$a0, X # Address of X
 li \$a1, 127 # a = 127
 li \$a2, 0xa # b = 0xa = 10 (decimal)

 jal addm # store s = a + b in memory at address X

la \$a0, Y # Address of Y
add \$a1, \$v0,\$zero # a = s
li \$a2, -3 # b = -3
jal addm # store a + b in memory at address Y

\$ra → nop

jr \$ra

No Operation – nop

An instruction that does nothing.

.data

X: .word 0x11111111

Y: .word 0x22222222

.text

.globl main

main: la \$a0, X # Address of X

li \$a1, 127 # a = 127

li \$a2, 0xa # b = 0xa = 10 (decimal)

jal addm # store s = a + b in memory at address X

la \$a0, Y # Address of Y

add \$a1, \$v0,\$zero # a = s

li \$a2, -3 # b = -3

jal addm # store a + b in memory at address Y

\$PC → nop

jr \$ra

.data

X: .word 0x11111111

Y: .word 0x22222222

.text

.globl main

main: la \$a0, X # Address of X

li \$a1, 127 # a = 127

li \$a2, 0xa # b = 0xa = 10 (decimal)

jal addm # store s = a + b in memory at address X

la \$a0, Y # Address of Y

add \$a1, \$v0,\$zero # a = s

li \$a2, -3 # b = -3

jal addm # store a + b in memory at address Y

nop

\$PC → jr \$ra

.data

X: .word 0x11111111

Y: .word 0x22222222

.text

.globl main

main: la \$a0, X # Address of X

li \$a1, 127 # a = 127

li \$a2, 0xa # b = 0xa = 10 (decimal)

jal addm # store s = a + b in memory at address X

la \$a0, Y # Address of Y

add \$a1, \$v0,\$zero # a = s

li \$a2, -3 # b = -3

jal addm # store a + b in memory at address Y

\$PC → nop

jr \$ra

.data

X: .word 0x11111111

Y: .word 0x22222222

.text

.globl main

main: la \$a0, X # Address of X
li \$a1, 127 # a = 127
li \$a2, 0xa # b = 0xa = 10 (decimal)

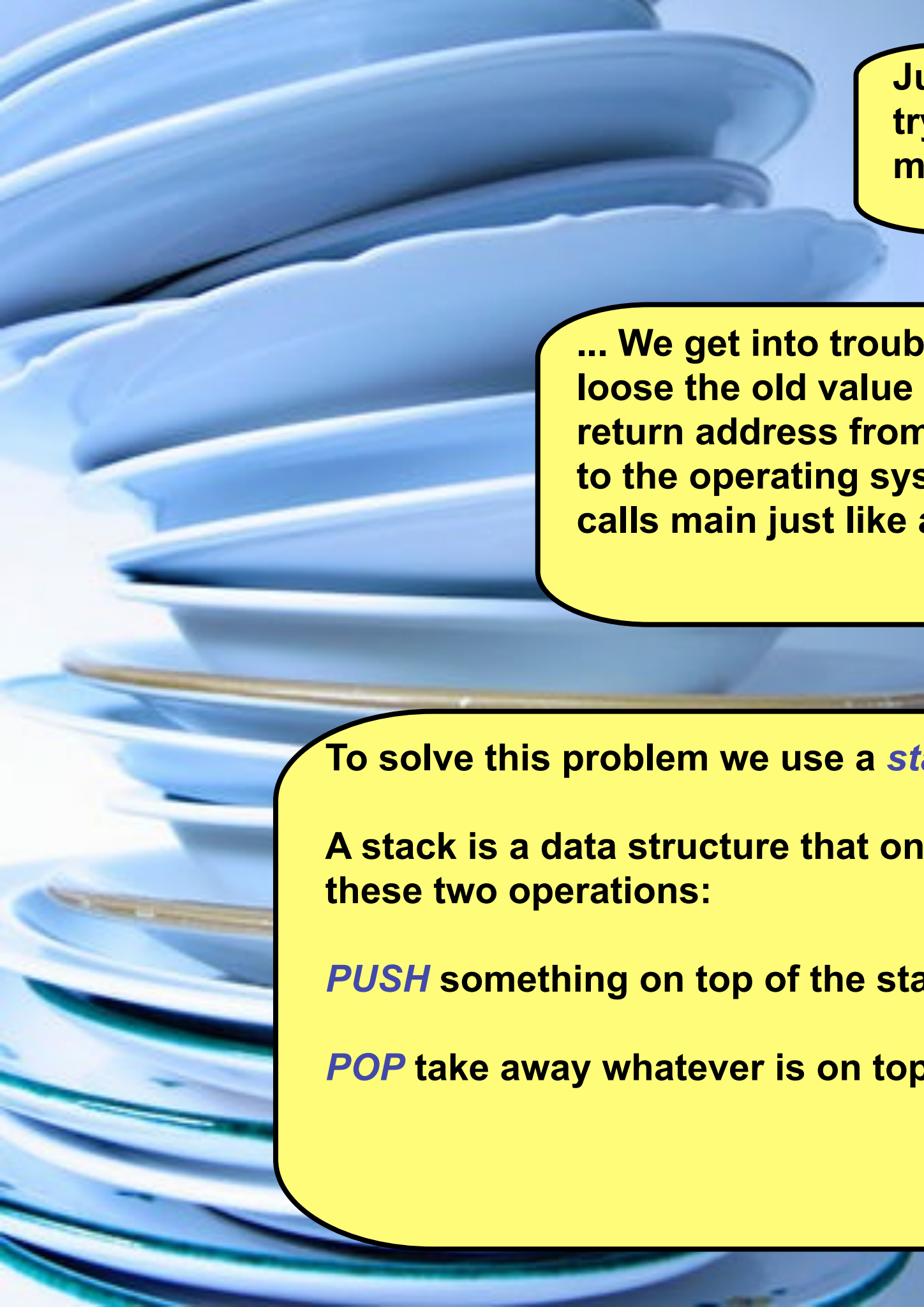
jal addm # store s = a + b in memory at address X

la \$a0, Y # Address of Y
add \$a1, \$v0,\$zero # a = s
li \$a2, -3 # b = -3
jal addm # store a + b in memory at address Y

\$ra → nop

\$PC → jr \$ra

We are stuck in an infinite loop since \$ra points here.



Just like it's a bad idea to try to pick a plate in the middle of the pile...

... We get into trouble when we lose the old value in \$ra, the return address from main back to the operating system (who calls main just like a subroutine).

To solve this problem we use a *stack*.

A stack is a data structure that only allows these two operations:

PUSH something on top of the stack.

POP take away whatever is on top of the stack.

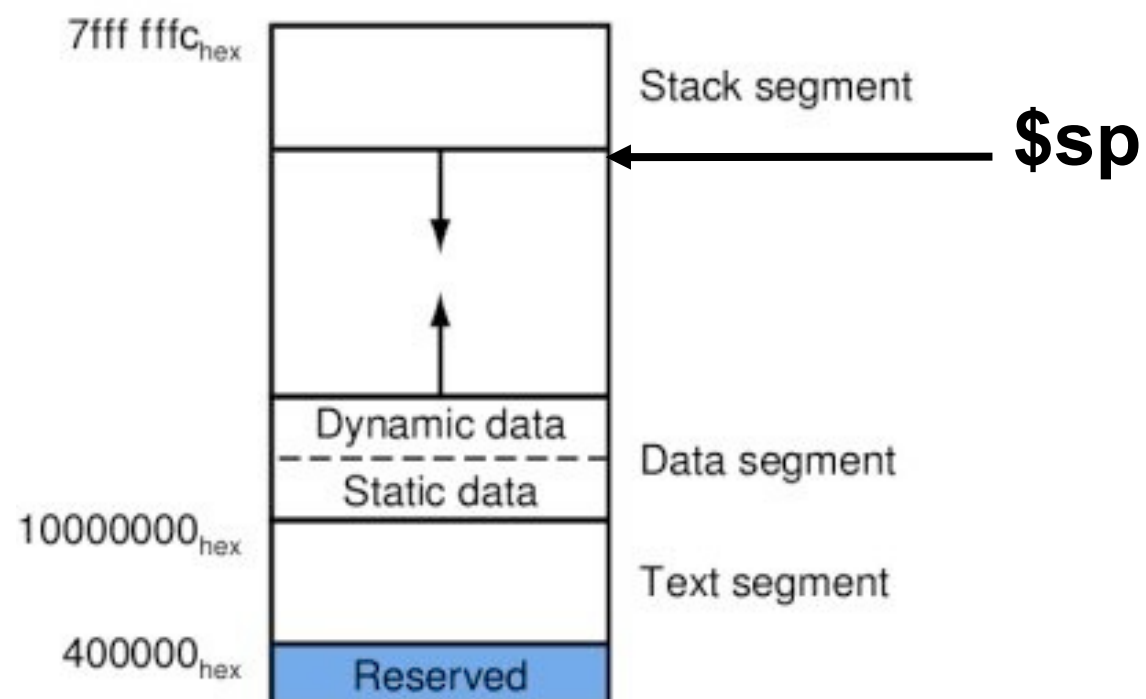




In MIPS the special $\$sp$ register is used to define the top of the stack.

The stack grows downward from high addresses to low addresses:

This is to not interfere with the data segment.



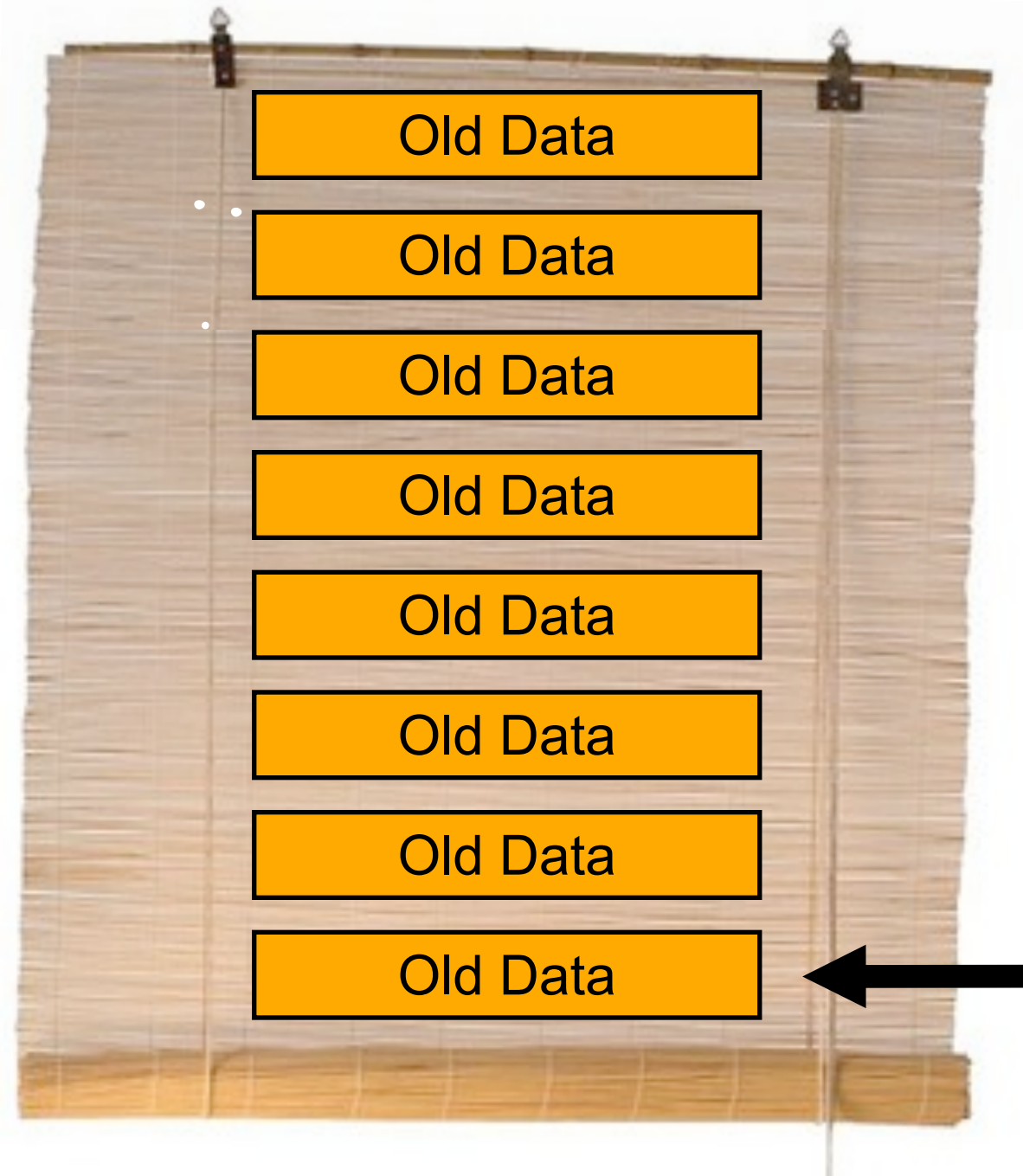
PUSH register \$t0

```
addi    $sp, $sp, -4
sw       $t0, 0($sp)
```

POP value on top of stack to \$t0

```
lw       $t0, 0($sp)
addi     $sp, $sp, 4
```


High Address



Low Address

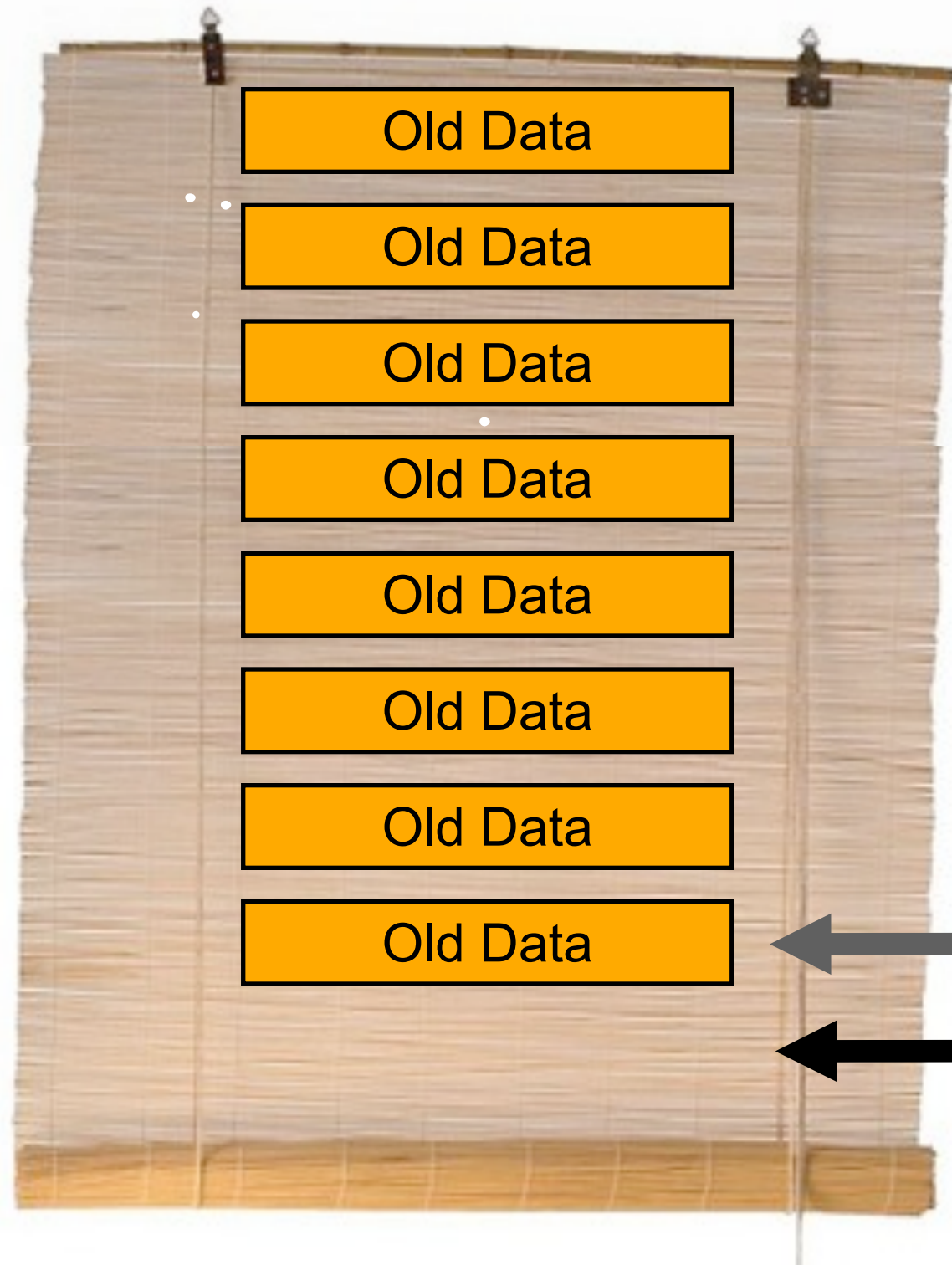
In MIPS, the stack is like a pull down curtain.

On the curtain we write data words (4 bytes).

\$sp



High Address



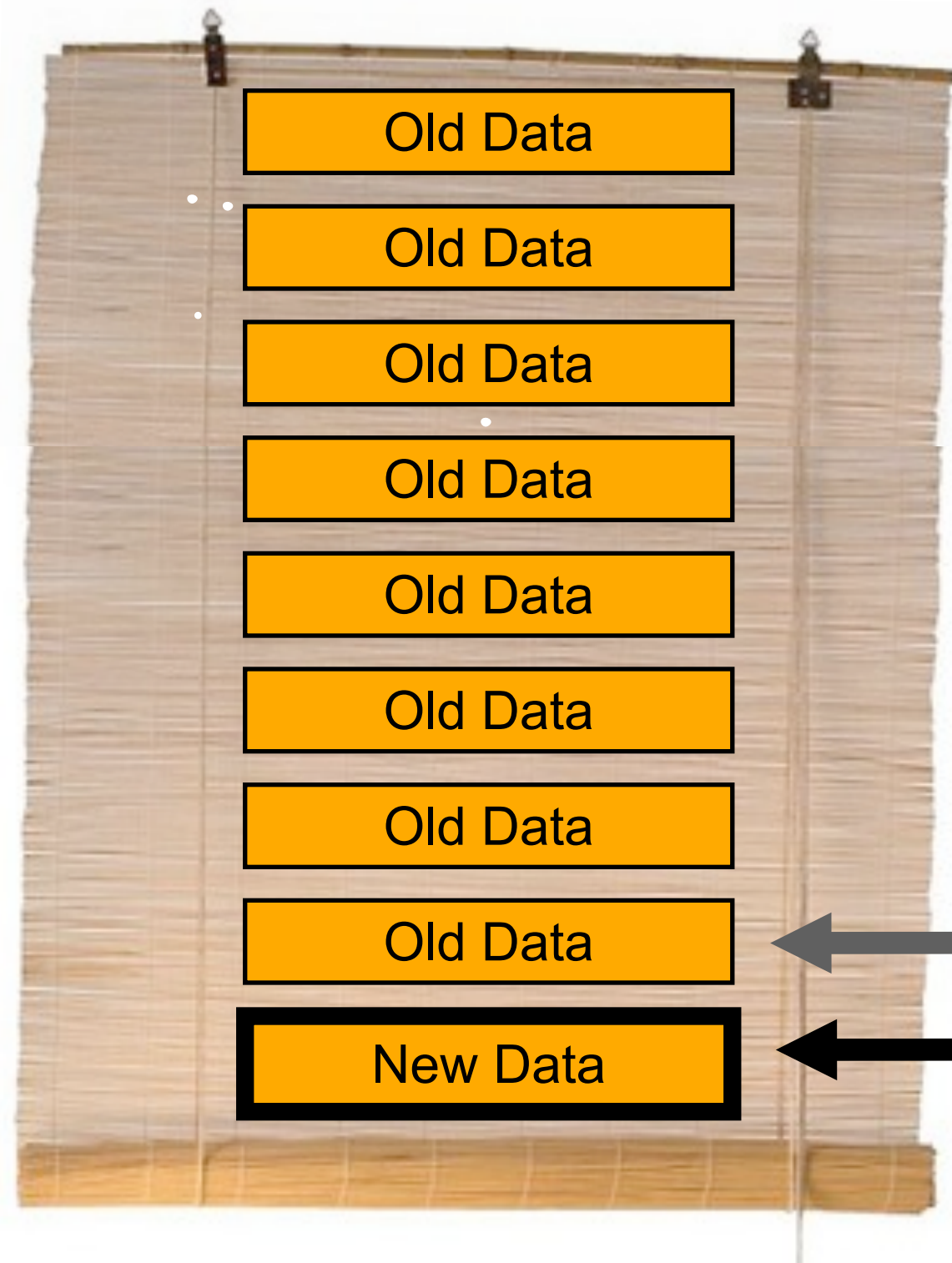
Low Address

To PUSH new data on the stack, pull the "curtain" down to make space for the new data.

```
addi $sp, $sp, -4
```



High Address



Low Address

Now you can store the new data (for example $\$t0$) by "writing" to the new space on the "curtain"

```
sw $t0, 0($sp)
```



```

.data
X:    .word 0x11111111
Y:    .word 0x22222222
.text
.globl main

```

```

main:  # Push return address
      addi    $sp, $sp, -4
      sw      $ra, 0($sp)

```

```

      la      $a0, X           # Address of X
      li      $a1, 127         # a = 127
      li      $a2, 0xa         # b = 0xa = 10 (decimal)
      jal     addm             # store s = a + b in memory at address X

```

```

      la      $a0, Y           # Address of Y
      add     $a1, $v0, $zero   # a = S
      li      $a2, -3          # b = -3
      jal     addm             # store a + b in memory at address Y

```

```

nop

```

```

      # Pop return address
      lw      $ra, 0($sp)
      addi    $sp, $sp, 4

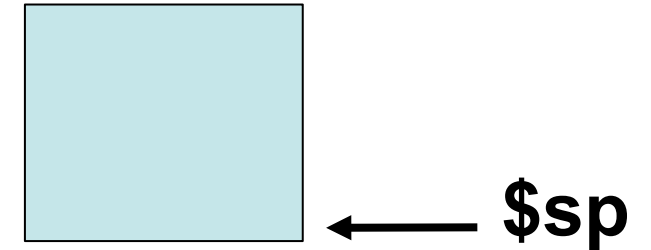
```

```

      jr      $ra

```

In the beginning of main we store the return address (back to the caller of main) on the stack (push).



Before the PUSH

```

.data
X:    .word 0x11111111
Y:    .word 0x22222222
.text
.globl main

```

```

main:  # Push return address
      addi    $sp, $sp, -4
      sw      $ra, 0($sp)

```

```

      la      $a0, X           # Address of X
      li      $a1, 127         # a = 127
      li      $a2, 0xa         # b = 0xa = 10 (decimal)
      jal     addm             # store s = a + b in memory at address X

```

```

      la      $a0, Y           # Address of Y
      add     $a1, $v0, $zero   # a = S
      li      $a2, -3          # b = -3
      jal     addm             # store a + b in memory at address Y

```

```

nop

```

```

      # Pop return address
      lw      $ra, 0($sp)
      addi    $sp, $sp, 4

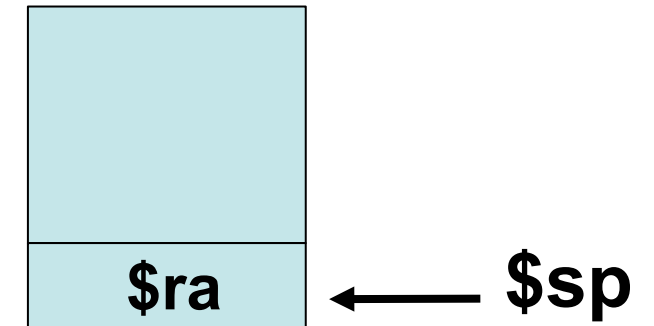
```

```

      jr      $ra

```

In the beginning of main we store the return address (back to the caller of main) on the stack (push).



After the PUSH

```

.data
X:    .word 0x11111111
Y:    .word 0x22222222
.text
.globl main

```

```

main:  # Push return address
      addi    $sp, $sp, -4
      sw      $ra, 0($sp)

      la      $a0, X           # Address of X
      li      $a1, 127         # a = 127
      li      $a2, 0xa         # b = 0xa = 10 (decimal)
      jal     addm             # store s = a + b in memory at address X

      la      $a0, Y           # Address of Y
      add     $a1, $v0, $zero   # a = S
      li      $a2, -3          # b = -3
      jal     addm             # store a + b in memory at address Y

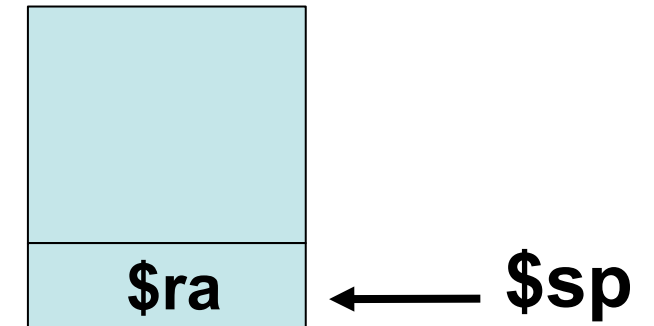
      nop

      # Pop return address
      lw      $ra, 0($sp)
      addi    $sp, $sp, 4

      jr      $ra

```

In the beginning of main we store the return address (back to the caller of main) on the stack (push).



Before the POP

At the end of main we restore the return address (back to the caller of main) by popping it from the stack.

```

X:      .data
        .word 0x11111111
Y:      .word 0x22222222
        .text
        .globl main

```

```

main:   # Push return address
        addi    $sp, $sp, -4
        sw      $ra, 0($sp)

```

```

        la      $a0, X           # Address of X
        li      $a1, 127         # a = 127
        li      $a2, 0xa         # b = 0xa = 10 (decimal)
        jal     addm             # store s = a + b in memory at address X

```

```

        la      $a0, Y           # Address of Y
        add     $a1, $v0, $zero   # a = S
        li      $a2, -3          # b = -3
        jal     addm             # store a + b in memory at address Y

```

```

nop

```

```

        # Pop return address
        lw      $ra, 0($sp)
        addi    $sp, $sp, 4

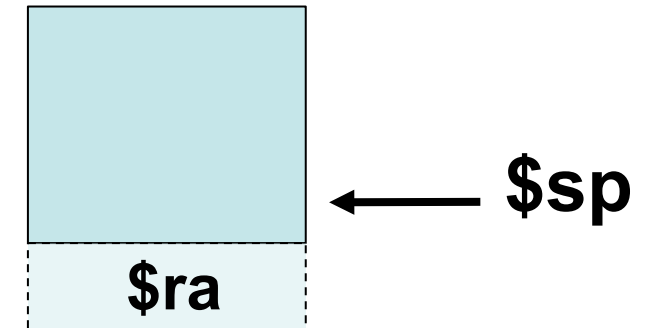
```

```

        jr      $ra

```

In the beginning of main we store the return address (back to the caller of main) on the stack (push).



After the POP

At the end of main we restore the return address (back to the caller of main) by popping it from the stack.



Ok, and this applies to main also since "someone" is calling main.

Whenever you write a subroutine that calls other subroutines, or calls itself (recursion), you must **push** the old value of \$ra on the stack.

Before you do jr \$ra at the end of the subroutine you must **pop** back the old value to \$ra from the stack..

For subroutines that does not call other subroutines (a.k.a **leaf procedures**) it is not neccessary to push and pop \$ra.



The MIPS Register Convention

The first four arguments to a subroutine are passed in the registers ***\$a0-\$a3***; subsequent arguments are passed on the stack.

Subroutine return values in ***\$v0, \$v1 registers***. In case of more than two return values, push return data on the stack.



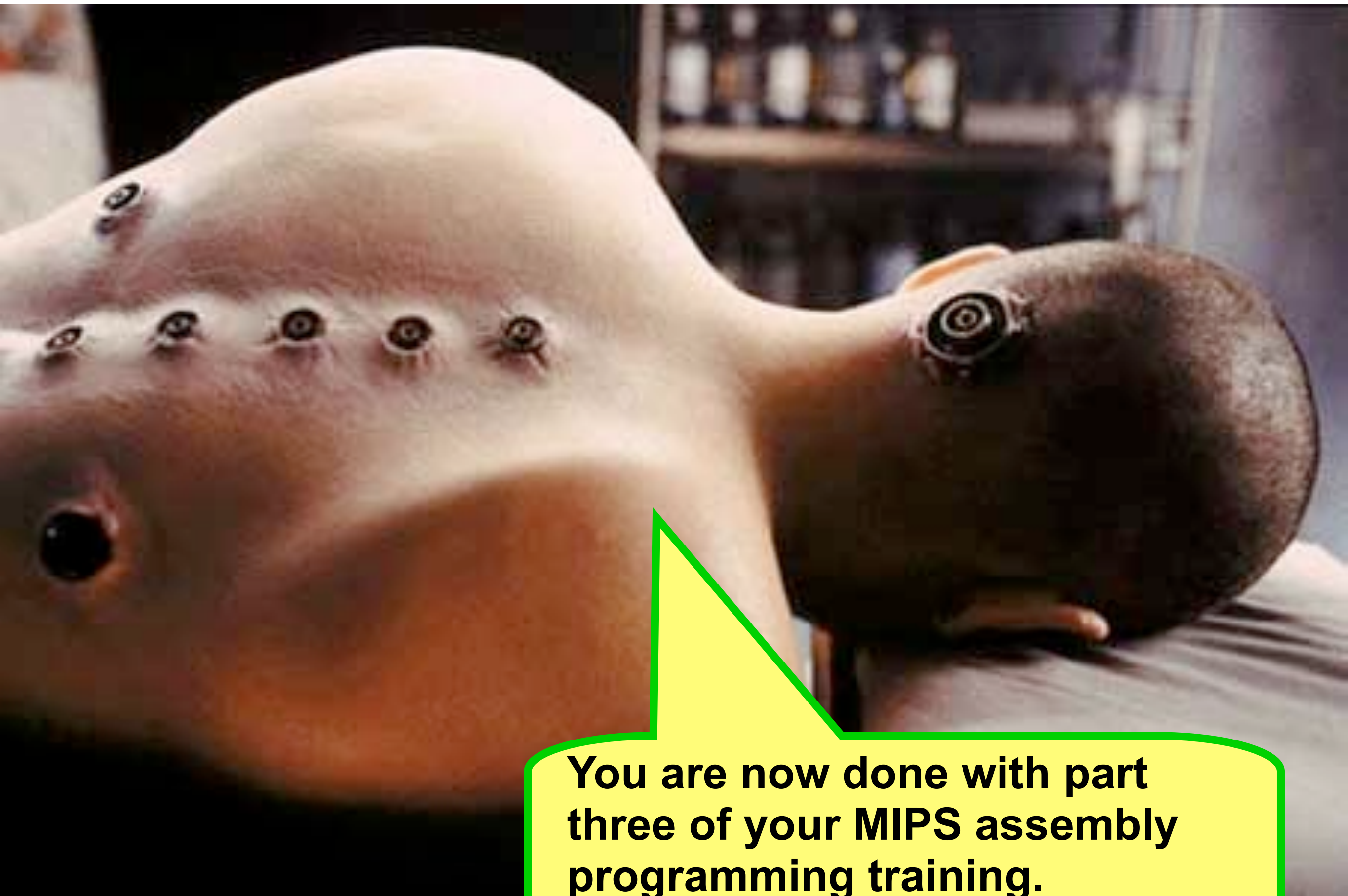
MIPS designates 8 registers ***\$s0-\$s7*** as ***saved registers***.

It's up to the callee to save the registers if they're being used.

If you use an s-register and call a subroutine, the contents of the s-registers must be unchanged after the call to the subroutine.

Temporary values in ***\$t-registers***.

Convention: A convention is a set of agreed, stipulated or generally accepted standards, norms, social norms or criteria, often taking the form of a custom.



You are now done with part three of your MIPS assembly programming training.