

Chapter 5

Type Declarations

(Version of 27 September 2004)

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5.1. Renaming existing types

Example: polynomials, revisited

Representation of a polynomial by a list of integers:

```
type poly = int list
```

- Introduction of an *abbreviation* for the type **int list**
- The two names denote the *same* type
- The object [4,2,8] is of type **poly** *and* of type **int list**

```
- type poly = int list ;  
  type poly = int list  
  
- type poly2 = int list ;  
  type poly2 = int list  
  
- val p:poly = [1,2] ;  
  val p = [1,2] : poly  
  
- val p2:poly2 = [1,2] ;  
  val p2 = [1,2] : poly2  
  
- p = p2 ;  
  val it = true : bool
```

5.2. Enumeration types

Declaration of a *new* type having a *finite* number of values

Example (`weekend.sml`)

```
datatype months = Jan | Feb | Mar | Apr | May | Jun
                | Jul | Aug | Sep | Oct | Nov | Dec
```

```
datatype days = Mon | Tue | Wed | Thu | Fri | Sat | Sun
```

```
fun weekend Sat = true
```

```
  | weekend Sun = true
```

```
  | weekend d = false
```

- **datatype** months = Jan | ... | Dec ;

```
datatype months = Jan | ... | Dec
```

- **datatype** days = Mon | ... | Sun ;

```
datatype days = Mon | ... | Sun
```

- **fun** weekend ... ;

```
val weekend = fn : days -> bool
```

- Convention (of this course):
constant identifiers (tags) start with an uppercase letter,
to avoid confusion with the function identifiers
- Possibility of using pattern matching
- Two datatype declarations cannot share the same constant,
as otherwise there would be typing problems
- Impossibility of defining sub-types in ML,
such as the integer interval `1..12` or the months `{Jul, Aug}`
- Equality is *automatically* defined on enumeration types

The bool and unit types, revisited

The types `bool` and `unit` are *not* primitive in ML:
they can be declared as enumeration types as follows:

```
datatype bool = true | false
```

```
datatype unit = ()
```

The order type

```
datatype order = LESS | EQUAL | GREATER
```

It is the result type of the comparison functions `Int.compare`,
`Real.compare`, `Char.compare`, `String.compare`, etc.

5.3. Constructed types

Example (person.sml)

```
datatype name = Name of string
datatype sex = Male | Female
datatype age = Age of int      (* expressed in years *)
datatype weight = WeightInKg of int
datatype person = Person of name * sex * age * weight
val ayşe = Person (Name "Ayşe", Female, Age 30, WeightInKg 58)
```

- **datatype** name = Name **of** string ;
datatype name = Name of string
- Name ;
val it = fn : string -> name
- **val** friend = Name "Ali" ;
val friend = Name "Ali" : name
- friend = "Ali" ;
! friend = "Ali"
! ^^^^^
! Error: operator and operand don't agree

The identifiers **Name**, **Age**, ..., **Person** are *value constructors*:

- Their declarations introduce collections of *tagged values*
- The type of a value constructor is a functional type; for instance, **Age** is of type **int** \rightarrow **age**
- The constructor **Age** may only be applied to an expression of type **int**, giving as result an object of type **age**
- If expression e reduces to the normal form n of type **int**, then **Age** e reduces to **Age** n , which is of type **age**
- The usage of the constructor **Age** is the *only* means of constructing an object of type **age**
- The tags of enumeration types are nothing else but value constructors without arguments!
- Value constructors can be used in patterns

Example

```

fun youngLady (Person ( _ , Female, Age a, _ )) = a <= 18
  | youngLady p = false
-   youngLady ayşe ;
  val it = false

```

Types with several constructors

Example 1: coordinates of a point in the plane

Handling Cartesian *and* polar coordinates (`coord.sml`)

```
datatype coord = Cart of real * real
                | Polar of real * real
```

Transformation of coordinates

```
fun toPolar (Polar (r,phi)) = ...
    | toPolar (Cart (x,y)) = ...

fun toCart (Cart (x,y)) = ...
    | toCart (Polar (r,phi)) = ...
```

Addition and multiplication of coordinates

```
fun addCoord c1 c2 =
    let val (Cart (x1,y1)) = toCart c1
        val (Cart (x2,y2)) = toCart c2
    in Cart (x1+x2, y1+y2) end

fun multCoord c1 c2 =
    let val (Polar (r1,phi1)) = toPolar c1
        val (Polar (r2,phi2)) = toPolar c2
    in Polar (r1*r2, phi1+phi2) end
```

Example 2: computation with integers and reals

Handling numbers, whether integers or reals (`number.sml`)

```
datatype number = Int of int
                | Real of real

fun toReal (Int i) = Real ( real i )
    | toReal (Real r) = Real r

fun addNumbers (Int a) (Int b) = Int (a+b)
    | addNumbers nb1 nb2 =
        let val (Real r1) = toReal nb1
            val (Real r2) = toReal nb2
        in Real (r1+r2) end
```

- The union of types can now be simulated
- The value constructors reveal the types of the objects
- Equality is defined on constructed types
if the types of the objects are not functional

Recursive types

Possibility of recursively using the declared type

Example 1: integer (linear) lists

```
datatype listInt = [ ]
                | :: of int * listInt
```

Example 2: integer binary trees

```
datatype bTreeInt = Void
                | Bt of int * bTreeInt * bTreeInt
```

Necessity of at least one non-recursive alternative!

The following objects are of type **bTreeInt**:

Void

Bt(12,Void,Void)

Bt(~5,Void,Void)

Bt(8+5, Bt(4,Void,Void), Void)

Bt(17, Bt(7,Void,Void), Bt(~1,Void,Void))

Bt(5, Bt(3, Bt(12, Bt(5,Void,Void), Bt(13,Void,Void)),

Bt(1, Bt(10,Void,Void), Void)),

Bt(8, Bt(5,Void,Void), Bt(4,Void,Void)))

Mutually recursive datatypes are declared together using **and** and **withtype**

5.4. Parameterised/polymorphic types

Example 1: (linear) lists

(Linear) lists are an example of polymorphic type:

```
datatype 'a list = [ ]
                | :: of 'a * 'a list
```

where **list** is called a *type constructor*

Example 2: binary trees

The type **bTree***int* defines binary trees with an *integer* as information associated to each node:

how to define binary trees of objects of type α ?

```
datatype 'a bTree = Void
                | Bt of 'a * 'a bTree * 'a bTree
```

where **bTree** is a *type constructor*

- **Bt**(2,Void,Void) is of type **int bTree**
- **Bt**("Ali",Void,Void) is of type **string bTree**
- **Void** is of type α **bTree**
- **Bt**(2, **Bt**("Ali",Void,Void), **Void**) is *not* a binary tree

Example 3: the predefined option type

The `option` type gives a new approach to partial functions

```
datatype 'a option = NONE
                | SOME of 'a
```

The predefined `valOf` function “removes the **SOME**” and “changes” the type of its argument:

```
function valOf expr
```

```
TYPE: 'a option  $\rightarrow$  'a
```

```
PRE: expr is of the form SOME e, otherwise exception Option is raised
```

```
POST: e
```

- `valOf (SOME (19+3)) ;`
`val it = 22 : int`
- `valOf NONE ;`
`! uncaught exception Option`

Example: factorial (`ch03/fact.sml`)

```
fun factOpt n =
```

```
  if n < 0 then NONE
```

```
  else if n = 0 then SOME 1
```

```
  else SOME (n * valOf (factOpt (n-1)))
```

- `factOpt 4 ;`
`val it = SOME 24 : int option`
- `factOpt ~3 ;`
`val it = NONE : int option`

5.5. Exceptions, revisited

- **exception** StopError ;
exception StopError
- **exception** NegArg of int ;
exception NegArg of int

The identifiers **StopError** and **NegArg** are declared as *exception constructors*

Patterns may involve value and exception constructors

An exception is *propagated* until being *caught* by an *exception handler*, at the latest by the top-level one

Example: factorial (ch03/fact.sml)

local

```
fun factAux 0 = 1
  | factAux n = n * factAux (n-1)
```

in fun factExc n =

```
(if n < 0 then raise NegArg n else Int.toString (factAux n) )
  handle NegArg n => "fact " ^ Int.toString n ^ " is undefined"
```

end

- factExc 4 ;
val it = "24" : string
- factExc ~3 ;
val it = "fact ~3 is undefined" : string

5.6. Application: expression evaluation

Evaluation of an integer arithmetic expression containing occurrences of some variable

Representation of expressions as *syntax trees*

Datatype declarations

datatype variable = Var **of** string

datatype expression = Const **of** int

| Cont **of** variable

| Plus **of** expression * expression

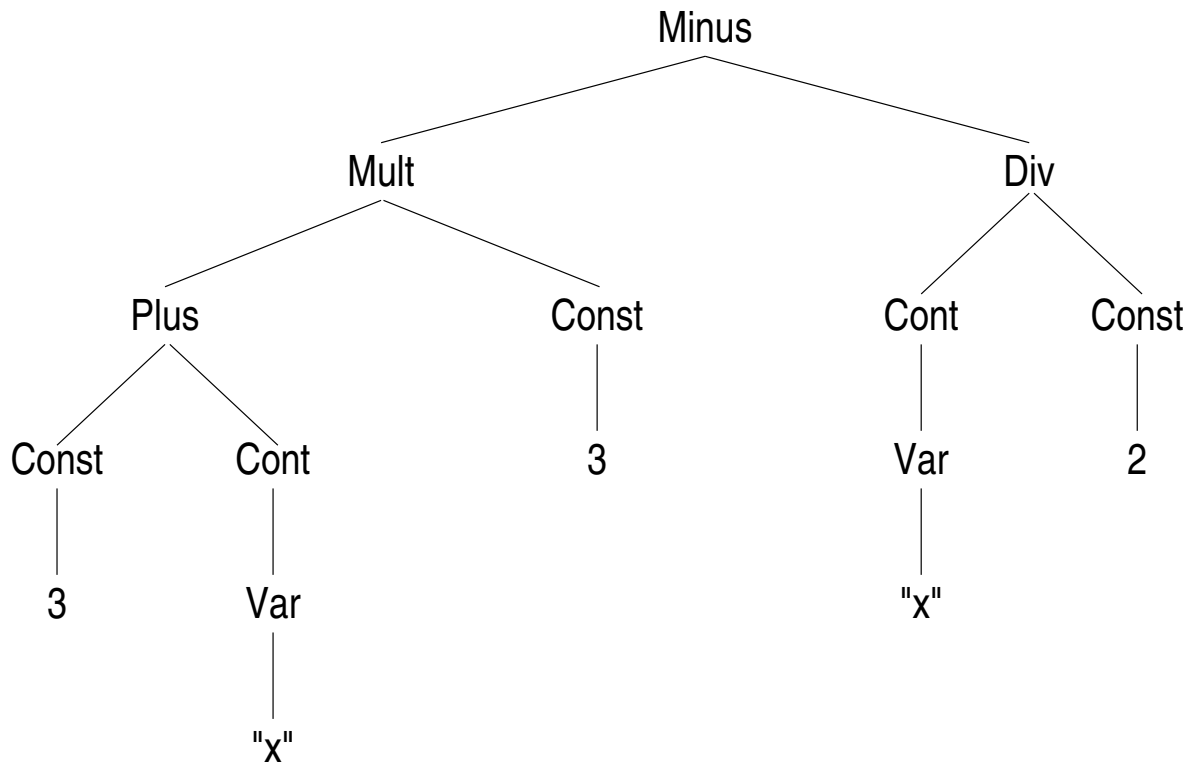
| Minus **of** expression * expression

| Mult **of** expression * expression

| Div **of** expression * expression

Example

The expression $(3 + x) * 3 - x/2$
is represented by the following syntax tree:



and is written in ML with our new datatypes as follows:

```

Minus ( Mult ( Plus ( Const 3, Cont (Var "x") ),
                Const 3 ),
        Div ( Cont (Var "x"), Const 2 ) )
  
```

Evaluation (eval.sml)

function eval exp var val

TYPE: expression \rightarrow variable \rightarrow int \rightarrow int

PRE: var is the only variable that may appear in exp

POST: the value of exp where var is replaced by val

```

fun eval (Const n) var v = n
  | eval (Cont a) var v =
      if a <> var then error "eval: pre-condition violated"
      else v
  | eval (Plus(e1,e2)) var v =
      let val n1 = eval e1 var v
          val n2 = eval e2 var v
      in n1 + n2 end
  | eval (Minus(e1,e2)) var v =
      let val n1 = eval e1 var v
          val n2 = eval e2 var v
      in n1 - n2 end
  | eval (Mult(e1,e2)) var v =
      let val n1 = eval e1 var v
          val n2 = eval e2 var v
      in n1 * n2 end
  | eval (Div(e1,e2)) var v =
      let val n1 = eval e1 var v
          val n2 = eval e2 var v
      in if n2 = 0 then error "eval: division by 0"
          else n1 div n2
      end

```

For the expression $(3 + x) * 3 - x/2$
the call `eval (Minus (...)) (Var "x") 5` returns the value `22`

Exercises

- How to integrate variables and constants of type **real**?
- Extend the `eval` function such that the expression may contain *several* variables: the second argument could then be a *list* of (variable, value) pairs