

Topic 14: Propagation¹

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Course 1DL441:
Combinatorial Optimisation and Constraint Programming,
whose part 1 is Course 1DL451:
Modelling for Combinatorial Optimisation

¹Based partly on material by Christian Schulte



Outline

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Outline

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Outline

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

Example (Agricultural experiment design, AED)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—	✓	—	✓
spelt	—	—	✓	✓	—	—	✓
wheat	—	—	✓	—	✓	✓	—

Constraints to be satisfied:

- 1 Equal growth load: Every plot grows 3 grains.
- 2 Equal sample size: Every grain is grown in 3 plots.
- 3 Balance: Every grain pair is grown in 1 common plot.

Instance: 7 plots, 7 grains, 3 grains/plot, 3 plots/grain, balance 1.

General term: **balanced incomplete block design (BIBD)**.



Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

Example (Agricultural experiment design, AED)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	1	1	1	0	0	0	0
corn	1	0	0	1	1	0	0
millet	1	0	0	0	0	1	1
oats	0	1	0	1	0	1	0
rye	0	1	0	0	1	0	1
spelt	0	0	1	1	0	0	1
wheat	0	0	1	0	1	1	0

Constraints to be satisfied:

- 1 Equal growth load: Every plot grows 3 grains.
- 2 Equal sample size: Every grain is grown in 3 plots.
- 3 Balance: Every grain pair is grown in 1 common plot.

Instance: 7 plots, 7 grains, 3 grains/plot, 3 plots/grain, balance 1.

General term: **balanced incomplete block design (BIBD)**.



In a BIBD, the plots are **blocks** and the grains are **varieties**:

Example (BIBD *integer* model: $\checkmark \rightsquigarrow 1$ and $- \rightsquigarrow 0$)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators

```
-3 enum Varieties; enum Blocks;
-2 int: blockSize; int: sampleSize; int: balance;
-1 array[Varieties,Blocks] of var 0..1: BIBD;
0 solve satisfy;
1 constraint forall(b in Blocks)
    (blockSize = count(BIBD[..,b], 1));
2 constraint forall(v in Varieties)
    (sampleSize = count(BIBD[v,..], 1));
3 constraint forall(v, w in Varieties where v < w)
    (balance = count([BIBD[v,b]*BIBD[w,b] | b in Blocks], 1));
```

Example (Instance data for our AED)

```
-3 Varieties = {barley,...,wheat}; Blocks = {plot1,...,plot7};
-2 blockSize = 3; sampleSize = 3; balance = 1;
```



Store after filling the first four rows

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	?						
spelt							
wheat							

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Store after filling the first four rows

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	?						
spelt							
wheat							

But plot1 **cannot** grow rye as that would violate the first constraint (every plot grows 3 grains).

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators



Store after filling the first four rows

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—						
spelt							
wheat							

But plot1 **cannot** grow rye as that would violate the first constraint (every plot grows 3 grains).

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators



Store after filling the first four rows

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—						
spelt							
wheat							

But plot1 **cannot** grow rye as that would violate the first constraint (every plot grows 3 grains). Actually, plot1 **cannot** grow oats, spelt, or wheat either, for the same reason, and this was **already propagated** when **trying the search guess** that plot1 grow millet!

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators



Store after filling the first four rows

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—						
spelt	—						
wheat	—						

But plot1 **cannot** grow rye as that would violate the first constraint (every plot grows 3 grains). Actually, plot1 **cannot** grow oats, spelt, or wheat either, for the same reason, and this was **already propagated** when **trying the search guess** that plot1 grow millet!

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators



Continuing ...

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	?					
spelt	—						
wheat	—						

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators



Continuing ...

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	?					
spelt	—						
wheat	—						

Guess: Let plot2 grow rye. Strategy: ✓ guesses first.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Continuing ...

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓					
spelt	—						
wheat	—						

Guess: Let plot2 grow rye. Strategy: ✓ guesses first.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓					
spelt	—						
wheat	—						

Propagation: plot2 **cannot** grow spelt and wheat as otherwise the first constraint (every plot grows 3 grains) would be violated for plot2.



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓					
spelt	—	—					
wheat	—	—					

Propagation: plot2 **cannot** grow spelt and wheat as otherwise the first constraint (every plot grows 3 grains) would be violated for plot2.



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓					
spelt	—	—					
wheat	—	—					

Propagation: plot3, plot4, and plot6 **cannot** grow rye as otherwise the third constraint (every grain pair is grown in 1 common plot) would be violated.



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—		—	
spelt	—	—					
wheat	—	—					

Propagation: plot3, plot4, and plot6 **cannot** grow rye as otherwise the third constraint (every grain pair is grown in 1 common plot) would be violated.



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—		—	
spelt	—	—					
wheat	—	—					

Propagation: plot5 and plot7 **must** grow rye as otherwise the second constraint (every grain is grown in 3 plots) would be violated for rye.



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—	✓	—	✓
spelt	—	—					
wheat	—	—					

Propagation: plot5 and plot7 **must** grow rye as otherwise the second constraint (every grain is grown in 3 plots) would be violated for rye.



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—	✓	—	✓
spelt	—	—					
wheat	—	—					

Propagation: plot3 **must** grow spelt and wheat as otherwise the first constraint (every plot grows 3 grains) would be violated for plot3.



Continuing ...

Example (BIBD *integer* model)

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
ConstraintFixpoint of Multiple
Propagators

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—	✓	—	✓
spelt	—	—	✓				
wheat	—	—	✓				

Propagation: plot3 **must** grow spelt and wheat as otherwise the first constraint (every plot grows 3 grains) would be violated for plot3.



Continuing ...

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—	✓	—	✓
spelt	—	—	✓				
wheat	—	—	✓				

Propagation: No more propagation possible.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Continuing ...

Example (BIBD *integer* model)

	plot1	plot2	plot3	plot4	plot5	plot6	plot7
barley	✓	✓	✓	—	—	—	—
corn	✓	—	—	✓	✓	—	—
millet	✓	—	—	—	—	✓	✓
oats	—	✓	—	✓	—	✓	—
rye	—	✓	—	—	✓	—	✓
spelt	—	—	✓	✓			
wheat	—	—	✓				

Guess: Let plot4 grow spelt. Strategy: ✓ guesses first.

Propagation: etc.



Outline

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

a		1	2	3	4	5	6	7	8	9
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



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Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

State $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported values **of** a :

a		1	2	3	4	5	6	7	8	9
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



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Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

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Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

State $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported values **of** b :

a			2		4		6		8	
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



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Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

a			2		4		6		8	
b			2	3	4	5				

Keep propagator for $2 \cdot a + 4 \cdot b = 24$, as **not subsumed**:
its constraint is not definitely true under the current store.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

State $a + b = 9$: **prune** unsupported values **of** a :

a			2		4		6		8	
b			2	3	4	5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

State $a + b = 9$: **prune** unsupported values **of** a :

a			2		4		6		8	
b			2	3	4	5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

State $a + b = 9$: **prune** unsupported values **of** b :

a					4		6			
b			2	3	4	5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



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Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

State $a + b = 9$: **prune** unsupported values **of** b :

a					4		6			
b			2	3	4	5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

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Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

a					4		6			
b				3		5				

Keep propagator for $a + b = 9$, as **not subsumed**:
its constraint is not definitely true under the current store.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

Run $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported values **of a** :

a					4		6			
b				3		5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

Run $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported values **of a** :

a					4		6			
b				3		5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

Run $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported values **of** b :

a							6			
b				3		5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

Run $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported values **of** b :

a							6			
b				3		5				

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

a							6			
b				3						

Dispose of propagator for $2 \cdot a + 4 \cdot b = 24$, as **subsumed**:
its constraint is definitely true under the current store.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

Run $a + b = 9$: **prune** unsupported values **of** a :

a						6			
b				3					

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

Run $a + b = 9$: **prune** unsupported values of b :

a						6			
b			3						

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

a							6			
b				3						

Dispose of propagator for $a + b = 9$, as **subsumed**:
its constraint is definitely true under the current store.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

a							6			
b				3						

No propagators are left: all solutions are found. No search!

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Domain* Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

$$a + b = 9$$

a						6			
b				3					

This **general** propagation method works for **all** systems of constraints (linear or not, equalities or inequalities, etc), no matter how many constraints and decision variables.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Outline

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Problem, Model, and Propagation

Example (Propagation to *Bounds*(*) Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

a		1	2	3	4	5	6	7	8	9
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Bounds*(*) Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

State $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported bounds **of a** :

a		1	2	3	4	5	6	7	8	9
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Bounds*(*) Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

State $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported bounds **of a** :

a		1	2	3	4	5	6	7	8	9
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Bounds*(*) Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

State $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported bounds **of** b :

a			2	3	4	5	6	7	8	
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Bounds*(*) Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

State $2 \cdot a + 4 \cdot b = 24$: **prune** unsupported bounds **of** b :

a			2	3	4	5	6	7	8	
b	0	1	2	3	4	5	6	7	8	

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Bounds*(*) Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

a			2	3	4	5	6	7	8	
b			2	3	4	5				

Keep the propagator for $2 \cdot a + 4 \cdot b = 24$, as **not subsumed**.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Problem, Model, and Propagation

Example (Propagation to *Bounds*(*) Consistency)

Find $a \in \{1, 2, \dots, 9\}$ and $b \in \{0, 1, \dots, 8\}$ such that

$$2 \cdot a + 4 \cdot b = 24$$

a			2	3	4	5	6	7	8	
b			2	3	4	5				

Keep the propagator for $2 \cdot a + 4 \cdot b = 24$, as **not subsumed**.

Some propagators are left: no solutions found yet. Search!

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Outline

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Outline

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators

Definition (Strict partial order)

A **strict partial order** is a pair $\langle X, \sqsubset \rangle$, where X is a set over which the binary relation \sqsubset is irreflexive ($\forall x \in X : x \not\sqsubset x$) and transitive ($\forall x, y, z \in X : x \sqsubset y \wedge y \sqsubset z \Rightarrow x \sqsubset z$).

Example: $(\mathbb{Z}, <)$ is a strict partial order.

Definition (Well-founded order)

A **well-founded order** is a strict partial order $\langle X, \sqsubset \rangle$ in which there is no infinite decreasing sequence $\dots \sqsubset x_3 \sqsubset x_2 \sqsubset x_1$.

Examples: $(\mathbb{N}, <)$; $(2^S, \subset)$ for a set S ; and loop variants.

Definition (Lexicographic order)

Given two well-founded orders $\langle X, \sqsubset_X \rangle$ and $\langle Y, \sqsubset_Y \rangle$, the **lexicographic order** $\langle X \times Y, \sqsubset_{\text{lex}} \rangle$ is well-founded, where $\langle x_1, y_1 \rangle \sqsubset_{\text{lex}} \langle x_2, y_2 \rangle$ iff either $x_1 \sqsubset_X x_2$ or $x_1 = x_2 \wedge y_1 \sqsubset_Y y_2$. Similarly for composing more than two orders.

Examples: `lex_less` is $(\mathbb{N}^*, <_{\text{lex}})$; loop variant of slide 28.



Functions

Intuition

- Example 1
- Example 2
- Example 3

Algorithms

- Reminders of Discrete Mathematics
- Solving: Overview
- Propagator for a Constraint
- Fixpoint of Multiple Propagators

Definition (Fixpoint)

A **fixpoint** of a function $f: X \rightarrow X$ is an element $x \in X$ that does not change under f , that is $f(x) = x$.

Example: A store of the set \mathbb{S} of all possible stores can be a fixpoint of a propagator, which is a total function in $\mathbb{S} \rightarrow \mathbb{S}$.

Idempotent functions compute fixpoints:

Definition (Idempotency)

A function f is **idempotent** iff it is equal to its composition with itself: $\forall x : f(f(x)) = f(x)$.

Example: A propagator $p: \mathbb{S} \rightarrow \mathbb{S}$ can be idempotent.



Outline

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Solving

Systematic search, for a satisfaction problem:

- 1: **propagate** all constraints; **backtrack** if empty domain
- 2: **if** only fixed variables, **then** show solution & **backtrack**
- 3: **while** there is at least one scheduled propagator **do**
- 4: **select** non-fixed variable v of current domain $\text{dom}(v)$
- 5: **partition** $\text{dom}(v)$ using **guesses** (say $v = d$ & $v \neq d$,
or $v > d$ & $v \leq d$, for a **selected** value $d \in \text{dom}(v)$)
- 6: **for each guess**: **recurse** upon adding it as constraint

For an **optimisation problem**: before backtracking at line 2
add the constraint that **any next solution must be better**.

Strategies:

- Line 4: **variable selection**: smallest domain, ...
- Line 5: **value selection**: maximum, median, ...
- Line 5: **guess selection**: equality, bisection, ...
- Tree **exploration**: depth-first search, ...



Strength of Stores

Definition (Store strength comparison, denoted $s \prec t$)

Store s is (strictly) stronger than store t

if and only if $s(v) \subseteq t(v)$ for **every** decision variable v ,
and $s(v) \subsetneq t(v)$ for **at least one** decision variable v .

(\mathbb{S}, \prec) is a well-founded (and hence partial) order.

Example (Store strength comparison)

Consider these stores for variables $\{x, y\}$ over $\{4, 5, 7\}$:

$$s_1 = \{x \mapsto \{4, 5\}, y \mapsto \{5, 7\}\}$$

$$s_2 = \{x \mapsto \{5\}, y \mapsto \{5, 7\}\}$$

$$s_3 = \{x \mapsto \{5, 7\}, y \mapsto \{4, 5, 7\}\}$$

Note: $s_2 \prec s_1$ and $s_2 \prec s_3$, but s_1 and s_3 are incomparable.

Intuition

Example 1

Example 2

Example 3

Algorithms

Reminders of
Discrete Mathematics

Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators



Outline

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview

Propagator for a
Constraint

Fixpoint of Multiple
Propagators

Definition (Propagator)

A **propagator** p_c for a constraint c modifies a store so that:

- **Contracting**: The result store is stronger than or equal to (\preceq) the input store: $p_c(s) \preceq s$ or $p_c(s) = s$, for any s .
- **Correct**: Each solution to c in the store remains there: $d \in s \Rightarrow d \in p_c(s)$, for any store s and solution d to c .
- **Checking**: For each solution to c , no domain is shrunk: $d \in c \Leftrightarrow p_c(s) = s$, for any fixed store s denoting d .
- **Monotonic** (optional): Strength-ordered stores remain ordered: $s_1 \preceq s_2 \Rightarrow p_c(s_1) \preceq p_c(s_2)$, for any s_1 and s_2 .

Example (Domain-consistency propagator for $x \leq y$)

$$p_{x \leq y}(s) = \left\{ \begin{array}{l} x \mapsto \{n \in s(x) \mid n \leq \max(s(y))\}, \\ y \mapsto \{n \in s(y) \mid n \geq \min(s(x))\} \end{array} \right\}$$

$$p_{x \leq y}(\{x \mapsto \{1, 3, 5, 9\}, y \mapsto \{0, 2, 4\}\}) = \{x \mapsto \{1, 3\}, y \mapsto \{2, 4\}\}$$



Motivation for Monotonicity

Counter-example

Consider the non-monotonic propagator for constraint c

$$p_c(s) = \text{if } s(x) = \{4, 5, 7\} \text{ then } \{x \mapsto \{4\}\} \text{ else } s$$

and the stores $s_1 = \{x \mapsto \{4, 5\}\}$ and $s_2 = \{x \mapsto \{4, 5, 7\}\}$:

$$s_1 \preceq s_2 \quad \text{but} \quad p_c(s_2) = \{x \mapsto \{4\}\} \preceq \{x \mapsto \{4, 5\}\} = p_c(s_1)$$

The result stores could also be incomparable;
note that \prec and \preceq are **partial** ordering relations.

But propagation would be propagator-order-dependent:

$$p_c(p_{x < 7}(s_2)) = \{x \mapsto \{4, 5\}\} \neq \{x \mapsto \{4\}\} = p_{x < 7}(p_c(s_2))$$

This might lead to unexpected solver behaviour.



Consequences of Propagator Definition

■ Property of propagation, if only monotonic propagators:

- **Order independence:** Propagators may be invoked in **any** order: their weakest common fixpoint is **unique**.
E.g., from $\{x, y \mapsto \{3, 4, 5\}\}$, **the** weakest fixpoint of $p_{x \geq y}$ and $p_{y > 3}$ is $\{x, y \mapsto \{4, 5\}\}$, whereas **a** strongest fixpoint is a solution store, such as $\{x, y \mapsto \{5\}\}$.

Only this property depends on monotonicity.

■ Properties of a propagator p_c for a constraint c :

- **Correctness:** Each monotonic propagator necessarily is correct, so the latter requirement does not have to be proven separately for a propagator proven monotonic.
- **Non-solution identification:** For a non-solution to c , the domain of some decision variable becomes empty.



Intuition

- Example 1
- Example 2
- Example 3

Algorithms

- Reminders of Discrete Mathematics
- Solving: Overview
- Propagator for a Constraint
- Fixpoint of Multiple Propagators

Idempotency of propagators is *not* required:

Every DC propagator is idempotent; a BC propagator may be non-idempotent: see Ex. 2.9 on p. 19 of [Course Notes](#).

Terminology:

In the literature, the deletion of domain values is also called **pruning**, **filtering**, **contraction**, or **narrowing**.

If a domain loses its last value, then we say that there was a domain **wipe-out**, and the propagator must **fail**.

Definition (Model)

A **model** of a CSP $\langle V, U, C \rangle$ is a tuple $\langle V, U, P \rangle$, where P is the set of propagators chosen for the constraints C . Similarly for a model of a COP.

For propagator algorithms, see Topic 16: Propagators.



Outline

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

1. Intuition

Example 1

Example 2

Example 3

2. Algorithms

Reminders of Discrete Mathematics

Solving: Overview

Propagator for a Constraint

Fixpoint of Multiple Propagators



Naïve Fixpoint Algorithm

Let $\langle V, U, P[f] \rangle$ be a model where there is a common domain U for all variables of V , without loss of generality.

Intuition

- Example 1
- Example 2
- Example 3

Algorithms

- Reminders of Discrete Mathematics
- Solving: Overview
- Propagator for a Constraint
- Fixpoint of Multiple Propagators

Let $s_0 = \{v \mapsto U \mid v \in V\}$ be the initial store, where every decision variable v of V is mapped to the universe U .

Call to build the root of the search tree: $\text{Propagate}(P, s_0)$.

```
function Propagate( $R, s$ )  
  while  $\exists q \in R : q(s) \not\preceq s$  do                                // variant:  $s$   
    select  $q \in R : q(s) \not\preceq s$   
     $s := q(s)$   
  return  $s$                                                        // post:  $s$  is a common fixpoint of  $R$ 
```



Toward More Realistic Propagation

Why is the previous algorithm naïve?

For the condition of its **while** loop:

- We may examine a propagator that does not depend in some sense on the propagator that was just run.
- We do not maintain the set of propagators that are known to be at fixpoint.

So we may examine a propagator that cannot prune values.

Variables of a propagator:

Let $\text{var}(p)$ denote the set of decision variables of the constraint implemented by propagator p :

- Running p has no effect on $\text{dom}(v)$, for $v \in V \setminus \text{var}(p)$.
- Running p is independent of $\text{dom}(v)$, for $v \in V \setminus \text{var}(p)$.



Variable-Directed Fixpoint Algorithm

Call to build the root of the search tree: $\text{Propagate}(P, P, s_0)$.

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

```
function Propagate( $R, Q, s$ ) //  $R$  = all prop.s; pre:  $Q \subseteq R$   
while  $Q \neq \emptyset$  do // invariant: every  $p \in R \setminus Q$  is at fixpt  
                        // variant:  $\langle s, |Q| \rangle$   
    select  $q \in Q$       // prop.s of  $Q$  are possibly not at fixpt  
     $Q := Q \setminus \{q\}$   
     $s' := q(s)$           //  $s' \preceq s$   
     $\text{ModVars} := \{v \in \text{var}(q) \mid s(v) \neq s'(v)\}$   
     $\text{DepProps} := \{p \in R \mid \exists v \in \text{var}(p) : v \in \text{ModVars}\}$   
     $Q := Q \cup \text{DepProps}$  // maybe  $q \in Q$ : optional idempot.  
     $s := s'$   
return  $s$              // post:  $s$  is a common fixpoint of  $R$ 
```



Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

Toward Further Improved Propagation

Propagators signal status to avoid some useless runs:

- Propagator p is **failed** upon a domain wipe-out.
- Propagator p is **subsumed** (or **entailed**) by store s iff all stronger stores are fixpoints: $\forall s' \preceq s : p(s') = s'$. This status is an obligation when s is a solution store. Such a propagator can safely be disposed of in the model.
- Otherwise, if so, ideally signal that p is **at fixpoint** for s .
- It is always safe to signal that a propagator p is **possibly not at fixpoint** for the result store s .

Examples (Subsumption)

$p_{x \leq y}$ is subsumed by $\{x \mapsto \{1, 3\}, y \mapsto \{3, 5\}\}$, but not by $\{x \mapsto \{1, 3, 4\}, y \mapsto \{3, 5\}\}$. A DC propagator of a unary constraint, like $x \in \{1, 3, 5\}$, is subsumed upon its first run.



Propagators with Status Message

Example (Domain-consistency propagator for $x \leq y$)

$p_{x \leq y}(s) = \text{let } s' = \left\{ \begin{array}{l} x \mapsto \{n \in s(x) \mid n \leq \max(s(y))\}, \\ y \mapsto \{n \in s(y) \mid n \geq \min(s(x))\} \end{array} \right\} \text{ in}$
 if $s'(x) = \emptyset \vee s'(y) = \emptyset$ **then** $\langle \text{Failed}, \emptyset \rangle$
 else if $\max(s'(x)) \leq \min(s'(y))$ **then** $\langle \text{Subsumed}, s' \rangle$
 else $\langle \text{AtFixpt}, s' \rangle$

Note that $\min(s(x))$ and $\max(s(y))$ do not change: hence s' is at least a fixpoint for $p_{x \leq y}$ and at best subsumes it!

Responsibility:

The burden of signalling, in reasonable runtime, a proper status message is on the programmer of a propagator.

Intuition

- Example 1
- Example 2
- Example 3

Algorithms

- Reminders of Discrete Mathematics
- Solving: Overview
- Propagator for a Constraint
- Fixpoint of Multiple Propagators



Propagator-Status-Directed Fixpoint Algo.

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint

Fixpoint of Multiple
Propagators

```
function Propagate( $R, Q, s$ ) // non-subsumed prop.s in  $R$ 
while  $Q \neq \emptyset$  do                                     // invariant: ...; variant: ...
    select  $q \in Q$ 
     $Q := Q \setminus \{q\}$ 
     $\langle m, s' \rangle := q(s)$                                      //  $s' \preceq s$ 
    if  $m = \text{Failed}$  then return  $\langle R, \emptyset \rangle$  endif
    if  $m = \text{Subsumed}$  then  $R := R \setminus \{q\}$  endif
     $\text{ModVars} := \{v \in \text{var}(q) \mid s(v) \neq s'(v)\}$ 
     $\text{DepProps} := \{p \in R \mid \exists v \in \text{var}(p) : v \in \text{ModVars}\}$ 
    if  $m = \text{AtFixpt}$  then  $\text{DepProps} := \text{DepProps} \setminus \{q\}$  endif
     $Q := Q \cup \text{DepProps}$ 
     $s := s'$ 
return  $\langle R, s \rangle$                                      // post:  $s$  is a common fixpoint of  $R$ 
```



Toward Even Further Improved Propagation

Signalling *how* domains were modified:

Mutually exclusive **modification events** for each variable v :

- 1 None(v): the domain of v was not changed.
- 2 Failed(v): the domain of v was wiped out.
- 3 Fixed(v): the domain of v was pruned to a singleton.
- 4 Min(v): the lower bound of $\text{dom}(v)$ was increased.
Max(v): the upper bound of $\text{dom}(v)$ was decreased.
- 5 Any(v): the domain of v was otherwise pruned.

Gecode: Min(v) and Max(v) are bundled into Bounded(v).

☞ It is often simple to decide whether a propagator remains at fixpoint depending on **how** another propagator prunes domains of decision variables they share: variable sharing is no longer the sole criterion for adding propagators to Q .



Propagator Conditions

Example (Domain-consistency propagator for $x \leq y$)

$$p_{x \leq y}(s) = \left\{ \begin{array}{l} x \mapsto \{n \in s(x) \mid n \leq \max(s(y))\}, \\ y \mapsto \{n \in s(y) \mid n \geq \min(s(x))\} \end{array} \right\}$$

☞ $\text{PropConds}(p_{x \leq y}) = \{\text{Min}(x), \text{Max}(y)\}$

Promise: If the propagator is at fixpoint, then it will remain at fixpoint, unless $\min(\text{dom}(x))$ or $\max(\text{dom}(y))$ changes.

Example (Domain-consistency propagator for $x \neq y$)

$$p_{x \neq y}(s) = \left\{ \begin{array}{l} x \mapsto s(x) \setminus \text{if } |s(y)| = 1 \text{ then } s(y) \text{ else } \emptyset, \\ y \mapsto s(y) \setminus \text{if } |s(x)| = 1 \text{ then } s(x) \text{ else } \emptyset \end{array} \right\}$$

☞ $\text{PropConds}(p_{x \neq y}) = \{\text{Fixed}(x), \text{Fixed}(y)\}$

Promise: If the propagator is at fixpoint, then it will remain at fixpoint, unless $\text{dom}(x)$ or $\text{dom}(y)$ becomes a singleton.

Intuition

- Example 1
- Example 2
- Example 3

Algorithms

- Reminders of Discrete Mathematics
- Solving: Overview
- Propagator for a Constraint

Fixpoint of Multiple Propagators



Assumptions

Responsibilities, under Gecode:

- The programmer of propagator p states $\text{PropConds}(p)$.
- The solver computes as follows the set $\text{Conds}(s, s')$ of propagator conditions raised by applying a propagator q to a store s , giving $s' = q(s)$:

Modification event	Conditions added to $\text{Conds}(s, s')$
Fixed(v)	Fixed(v), Bounded(v), Any(v)
Bounded(v)	Bounded(v), Any(v)
Any(v)	Any(v)
None(v)	(none)

- The solver **schedules** a propagator p (adds p to Q) if the conditions $\text{Conds}(s, s')$ raised by propagator q intersect with the propagator conditions $\text{PropConds}(p)$.

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators



Status-and-Condition-Directed Fixpt Algo.

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators

```

function Propagate( $R, Q, s$ )
  while  $Q \neq \emptyset$  do                                // invariant: ...; variant: ...
    select  $q \in Q$ 
     $Q := Q \setminus \{q\}$ 
     $\langle m, s' \rangle := q(s)$                                 //  $s' \preceq s$ 
    if  $m = \text{Failed}$  then return  $\langle R, \emptyset \rangle$  endif
    if  $m = \text{Subsumed}$  then  $R := R \setminus \{q\}$  endif
     $\text{ModVars} := \{v \in \text{var}(q) \mid s(v) \neq s'(v)\}$ 
     $\text{DepProps} :=$ 
     $\{p \in R \mid \text{Conds}(s, s') \cap \text{PropConds}(p) \neq \emptyset\}$ 
    if  $m = \text{AtFixpt}$  then  $\text{DepProps} := \text{DepProps} \setminus \{q\}$  endif
     $Q := Q \cup \text{DepProps}$ 
     $s := s'$ 
  return  $\langle R, s \rangle$                                 // post:  $s$  is a common fixpoint of  $R$ 

```



Yet Further Optimisations

Priorities: The set Q is implemented as a queue:

How to do “**select** $q \in Q$ ”?

- According to cost: cheapest first
- According to expected impact: highest impact first
- In general: first-in first-out queue

Propagator rewriting:

Example

When all domain values for x are smaller than those for y , then the propagator for $\max(x, y) = z$ can be replaced by the propagator for $y = z$.

Further reading:

For a more formal treatment of all these issues, including proofs, see the [Course Notes](#).

Intuition

Example 1
Example 2
Example 3

Algorithms

Reminders of
Discrete Mathematics
Solving: Overview
Propagator for a
Constraint
Fixpoint of Multiple
Propagators