

Algorithmic Program Verification

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**... with applications to
weak memory models ...**

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**... with applications to
weak memory models ...**

**Parosh Aziz Abdulla
Uppsala University**

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Parosh Aziz Abdulla

Algorithmic Program Verification, Spring 2022.

Postgraduate Course, 10 credits.

[News](#)

Next lecture: Wednesday, 2 February, 10:00-12:00, Zoom 589 726 9279.

[Background](#)

Current capabilities in computer technology allow enormously complicated implementations of such systems, making the task of producing error-free products more and more difficult. In particular, during the last decade, parallel systems have become a critical part of the infrastructure of our society due to the emergence of modern platforms such as multicores and cloud technology.

It is of great practical and economic importance to developing methods that make the design process less error-prone. There is a real need for techniques for testing and verifying software to guarantee a higher degree of reliability.

Reasoning about concurrent systems is often conducted under the fundamental assumption of sequential consistency (SC) where all components are strongly synchronized so that they all have a uniform view of the global state of the system. However, nowadays, most parallel software run on platforms that do not guarantee SC. More precisely, to satisfy demands on efficiency and energy-saving, such platforms implement optimizations that lead to the relaxation of the inter-component synchronization, hence offering only weak consistency guarantees. Weakly consistent platforms are found at all levels of system design, such as multiprocessors, cache protocols, programming languages, and cloud systems.

[Goal](#)

The participants will learn:

- The principles of model checking: a technique that has led to the most notable advances in algorithmic program verification during the last 20 years.
- Frameworks for verifying communication protocols, distributed systems and algorithms, timed systems, hardware circuits, concurrent programs, and multicore architectures.
- Modeling and verification through classical models such as timed automata, push-down automata, Petri nets, and lossy channels systems.
- Modeling and verifying concurrent programs running on weakly consistent platforms, such as x86-TSO, IBM Power, ARM, C11, and the cloud.

TSO, IBM Power, ARM, C11, and the cloud.

[Contents](#)

- Model checking.
- Infinite-state models.
- Reachability analysis.
- Petri nets.
- Timed automata.
- Push-down automata.
- Lossy channel systems.
- x86-TSO.
- C11.
- Well quasi-ordered systems.
- Program abstraction: monotonic abstraction, view abstraction.
- The Power and ARM architectures.
- Parameterized systems.
- Cloud platforms.

[Structure](#)

- 15 Lectures.
- Project work. The project will consist of implementing some of the algorithms discussed in the class.

[Examination](#)

- Weekly assignments.
- Project work.

[Prerequisites](#)

- Course suitable both for PhD and MSc students.
- Primary knowledge corresponding to three years of an undergraduate program in computer science. However, I do not assume any prior knowledge of formal methods, program verification, or weak memory models.

[Slides \(I will later add slides on weak consistency\)](#)

- L1. Petri Nets. [pdf](#). [keynote](#).
- L2. Well-Quasi-Orderings. [pdf](#). [keynote](#).
- L3. Lossy Channel Systems. [pdf](#). [keynote](#).
- L4. Backward Reachability Analysis. [pdf](#). [keynote](#).
- L5. Timed Petri Nets. [pdf](#). [keynote](#).
- L6. Timed Automata. [pdf](#). [keynote](#).
- L7-8. View Abstraction. [pdf](#). [keynote](#).
- L9-10. Recursive Programs. Push-Down Automata.

[Literature](#)

- [Parosh Aziz Abdulla. Well- and Better-Quasi-Ordered Transition Ssystems.](#)

[Contact Person](#)

- [Parosh Aziz Abdulla.](#)

Outline

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- **Background**

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- **Program Verification**

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- **Infinite-State Models**

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models

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models

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concurrency

common
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concurrency



motivation

concurrency

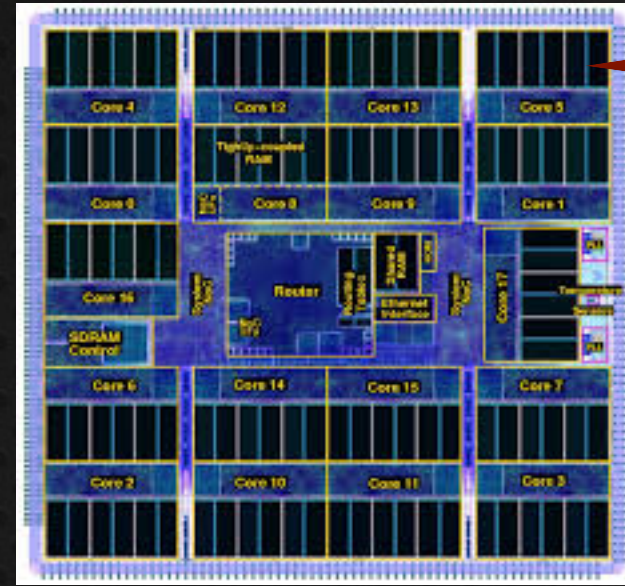
motivation

**Concurrent systems
are everywhere**

concurrency

motivation

Concurrent systems
are everywhere



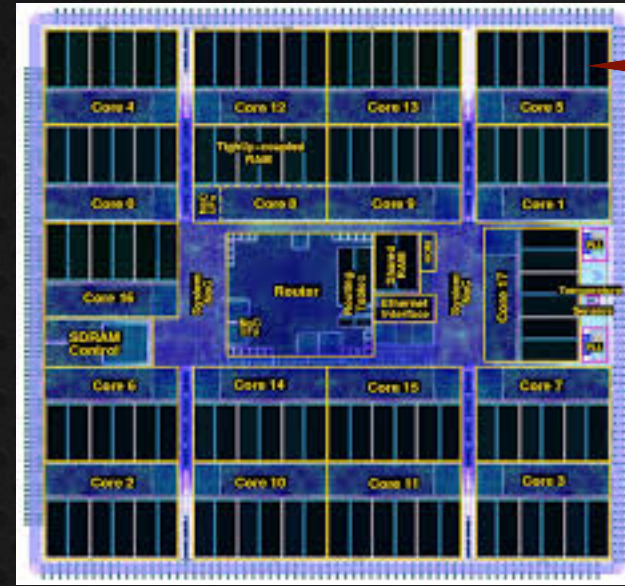
Multicore
architectures

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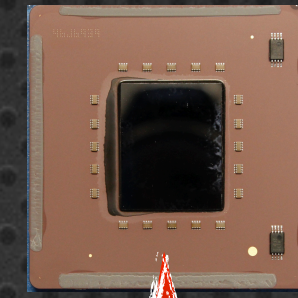
Multicore
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intel



ARM



IBM
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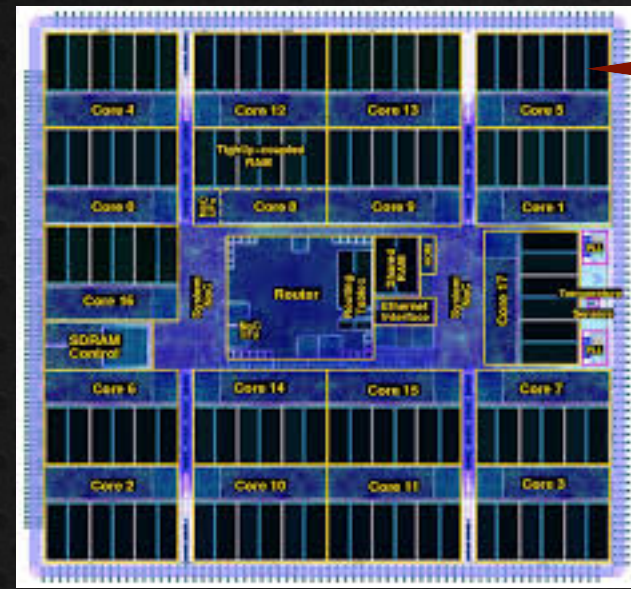


concurrency

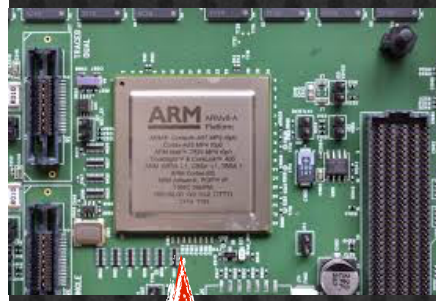
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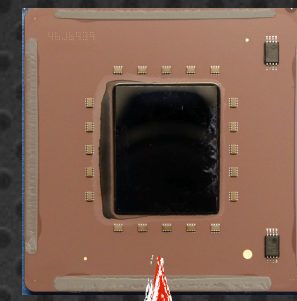
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ARM



IBM
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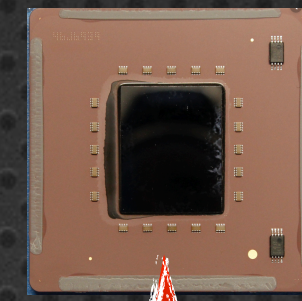
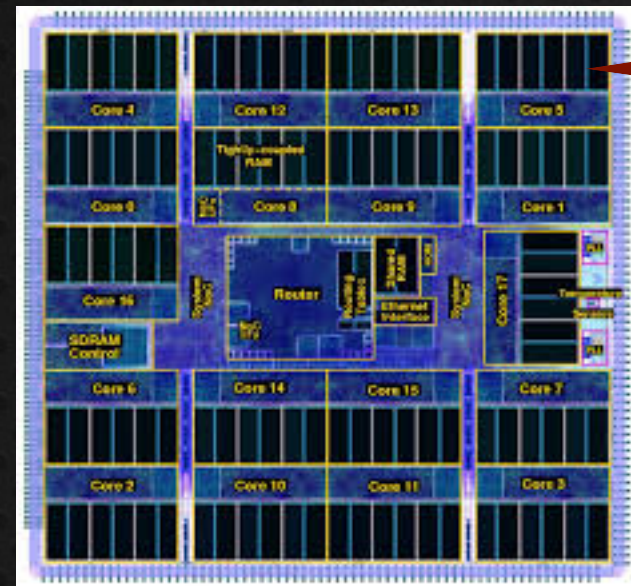
Distributed
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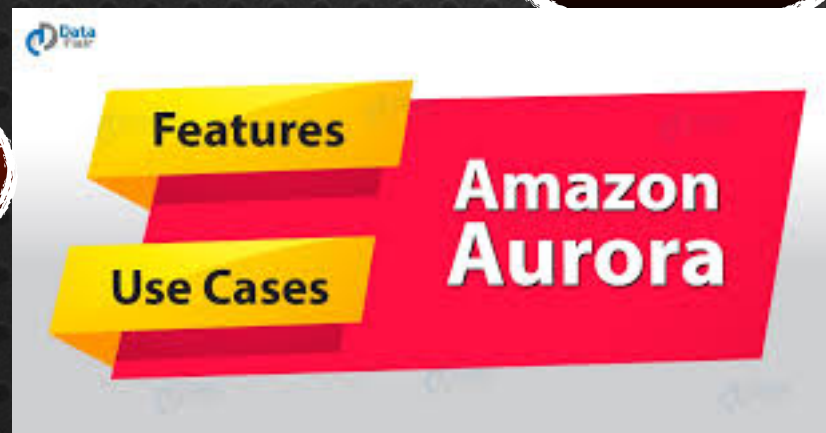
ARM

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Power



Facebook

Distributed
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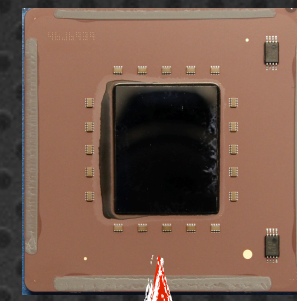
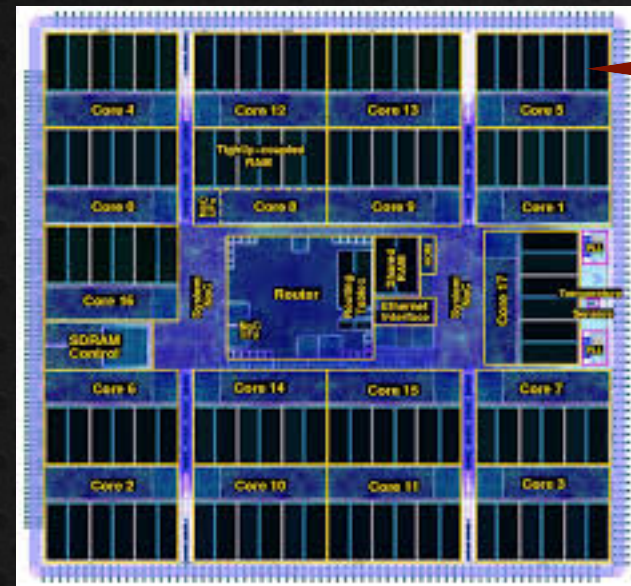


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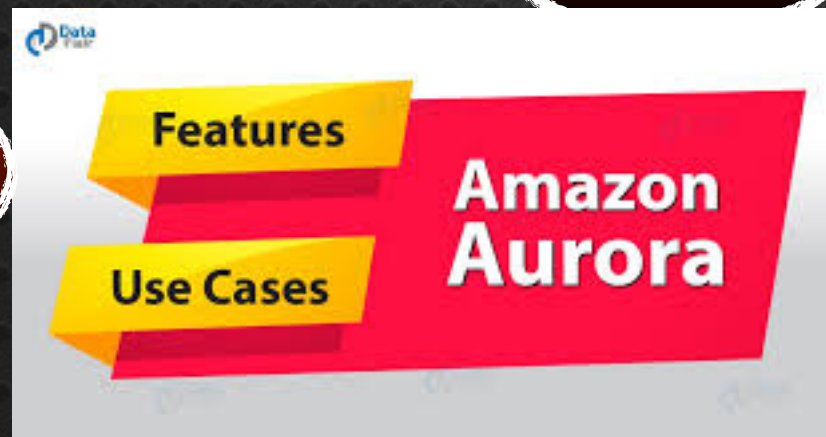


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Programming
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Distributed
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Sequential Consistency (SC)

- + simple & intuitive**
- expensive**

Weak Semantics (WS)

- + efficient, realistic**
- complicated**

**“order in which
data becomes
visible”**

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Architectures

- Intel x86: Total Store Ordering (TSO)
- ARM
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Programming Languages

- C11: Release-Acquire, Relaxed
- Java

Distributed Systems

- Eventual Consistency
- Causal Consistency

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“order in which
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Weak Persistence

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- Eventual Consistency
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- NVRAMs
- Intermittent Computing
- File Systems

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recent work



• NVRAMs

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Program Verification (SC)



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**Program Verification
(SC)**



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Model Checking

Model \models (safety) property

Classical Approach
Finite-State Systems

Model Checking
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Finite-State Systems

Model Checking

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Challenge:

Infinite-State Systems

Classical Approach
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Sources of "Infiniteness":

Classical Approach

Finite-State Systems

Model Checking

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Infinite-State Systems

Sources of "Infiniteness":

Unbounded Data Structures

- stacks (recursion)
- queues (protocols)
- counters (programs)
- clocks (time)
- lists, trees, graphs (heaps)

Classical Approach

Finite-State Systems

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Unbounded Control Structures

- parameterized systems
- multithreaded programs
- concurrent libraries
- Petri nets

Classical Approach

Finite-State Systems

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Multiple Sources:

- timed Petri nets
- recursive programs with unbounded data
- channels with time stamps
- etc