

Networks, Graphs and Matrices

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Networks

Network: a collection of items, called **vertices**, with connections between them, called **edges**.

Example: **social network**. The vertices represent persons and the edges relations, e.g. friendship relations

Networks are studied in mathematical **graph theory**

Connectedness properties of graphs:

- Are there subgraphs that are only loosely connected?
- Can a graph be divided into two separate graphs by breaking only a few edges?

For large graphs those questions are difficult to answer by visual inspection. Connectedness can be computed: eigenvalues and eigenvectors of the graph Laplacian matrix. This is **spectral graph theory**

Graphs

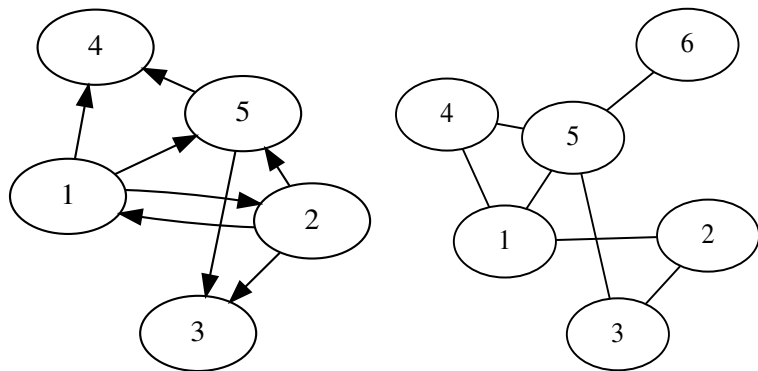
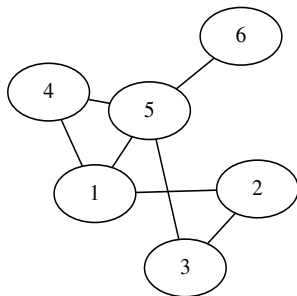


Figure: Directed and undirected graph.

Adjacency matrix

Adjacency matrix A of an undirected graph:

$$a_{ij} = \begin{cases} 0 & \text{if } i = j, \\ 1 & \text{if } i \neq j, \text{ and there is an edge between vertices } i \text{ and } j. \end{cases}$$

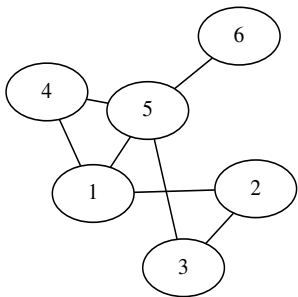


$$A = \begin{pmatrix} 0 & 1 & 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}.$$

Degrees

Degree of a vertex: the number of edges

The degrees of the vertices can be computed by the matrix-vector multiplication $d = Ae$, where e is a vector of all ones:



$$d = Ae = \begin{pmatrix} 3 \\ 2 \\ 2 \\ 2 \\ 4 \\ 1 \end{pmatrix}.$$

d is called the **degree vector**.

Connectedness and reducibility

A subgraph is a subset of the vertices and corresponding edges.

An undirected graph is called **connected** if there is no subgraph isolated from the rest of the graph.

Connectedness is equivalent to the concept of **irreducibility** of a matrix.

A symmetric matrix A is called **reducible** if there exists a permutation matrix P such that PAP^T is **block-diagonal**,

$$PAP^T = \begin{pmatrix} A_1 & 0 \\ 0 & A_2 \end{pmatrix}.$$

If there is no such permutation, then the matrix A is called **irreducible**.

Proposition *An undirected graph is connected if and only if its adjacency matrix is irreducible.*

Connectedness

Connectedness properties of graphs:

- Are there subgraphs that are only loosely connected?
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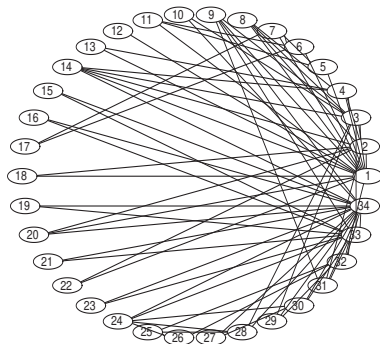


Figure: A rendition of the karate club graph.

Spectral graph partitioning

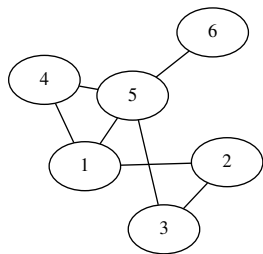
Degree matrix D : diagonal matrix with the degree vector on the diagonal:

$$D = \begin{pmatrix} d_1 & 0 & 0 & \cdots & 0 \\ 0 & d_2 & 0 & \cdots & 0 \\ \vdots & & \ddots & & \vdots \\ 0 & \cdots & 0 & d_{n-1} & 0 \\ 0 & \cdots & 0 & 0 & d_n \end{pmatrix}.$$

Laplacian matrix for an undirected graph with adjacency matrix A and degree matrix D :

$$L = D - A.$$

Laplacian



$$\begin{pmatrix} 3 & -1 & 0 & -1 & -1 & 0 \\ -1 & 2 & -1 & 0 & 0 & 0 \\ 0 & -1 & 2 & 0 & -1 & 0 \\ -1 & 0 & 0 & 2 & -1 & 0 \\ -1 & 0 & -1 & -1 & 4 & -1 \\ 0 & 0 & 0 & 0 & -1 & 1 \end{pmatrix}.$$

The sum of all elements in each row is equal to zero,

$$Le = De - Ae = 0,$$

so e is an eigenvector corresponding to the eigenvalue 0.

Second smallest eigenvalue

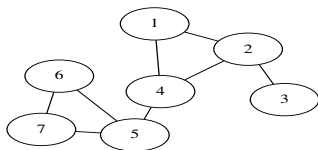
Assume for the moment that we have a disconnected graph with two subgraphs with adjacency matrix and Laplacian

$$A = \begin{pmatrix} A_1 & 0 \\ 0 & A_2 \end{pmatrix}, \quad L = \begin{pmatrix} L_1 & 0 \\ 0 & L_2 \end{pmatrix}$$

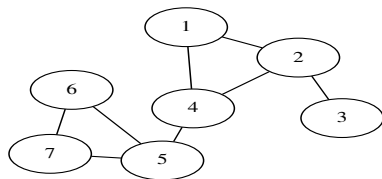
Then the Laplacian has a double eigenvalue equal to zero and eigenvectors

$$v_1 = \begin{pmatrix} e \\ 0 \end{pmatrix}, \quad v_2 = \begin{pmatrix} 0 \\ e \end{pmatrix}$$

Connect the subgraphs:



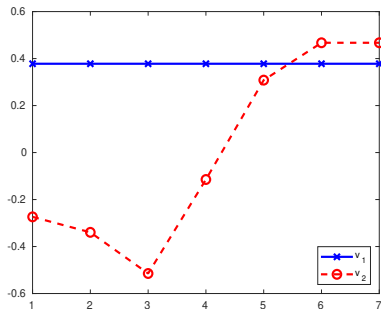
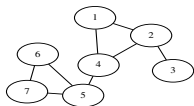
Adjacency matrix



$$A = \begin{pmatrix} 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 \end{pmatrix}$$

Eigenvectors v_1 and v_2

Eigenvalues: $\lambda_1 = 0$, $\lambda_2 = 0.34$



Eigenvectors of symmetric matrices are orthogonal: $v_1^T v_2 = 0$, so v_2 must have positive and negative entries

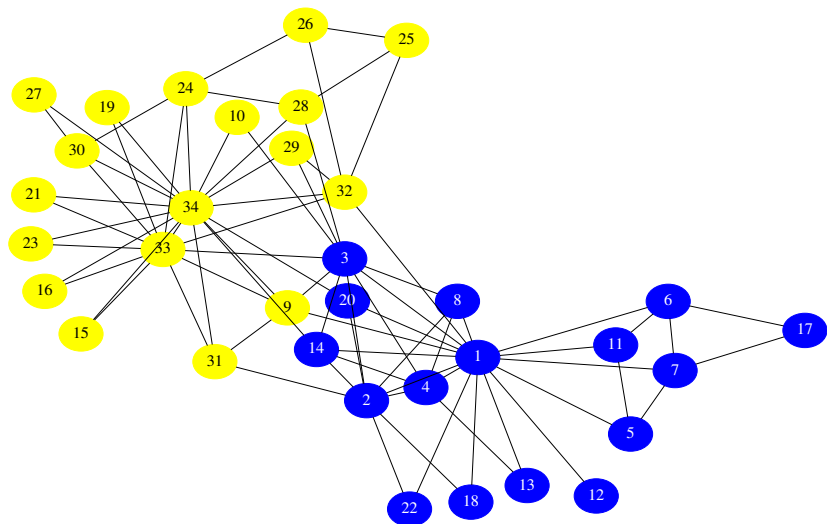
The values of the components of v_2 show how the vertices are located: 3 and 7 are furthest apart. 6 is as far away from 3 as 7. 4 is in the middle.

Spectral partitioning (simplified)

- 1 Given the Fiedler vector v_2 , reorder its elements in ascending order. This defines a permutation P of the integers $\{1, 2, \dots, n\}$, which induces a reordering of the vertices of the graph.
- 2 Apply the permutation to the vertices of the graph, and modify the set of edges accordingly. The adjacency matrix of the reordered graph is $\tilde{A} = PAP^T$.
- 3 For each partitioning in a neighborhood of the sign change, compute the corresponding cost for partitioning the graph. Choose the partitioning with the smallest conductance.

Cost: The number of edges that are broken.

Karate club graph



References

There is a huge literature on graphs and spectral partitioning. For a brief introduction, see Chapters 10 and 16 of



L. Eldén.

Matrix Methods in Data Mining and Pattern Recognition, Second Edition.

SIAM, 2019.