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An Algebraic Multigrid Tutorial

IMA Tutorial – Fast Solution Techniques November 28 - 29, 2010



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Outline

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- AMG Theory and Compatible Relaxation
- AMG for Electromagnetic Problems
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Preliminaries...

Consider solving the N×N linear system

$$Au = f$$

 Most iterative methods have the following form, where *r_k=f - Au_k* is the residual at iteration *k*

$$u_{k+1} = u_k + M^{-1}r_k$$

- Let $e_k = u u_k$ be the error, and note that $r_k = Ae_k$
- The error propagation for the iterative method is

$$e_{k+1} = (I - M^{-1}A)e_k$$



Multigrid linear solvers are optimal (O(N) operations), and hence have good scaling potential



 Weak scaling – want constant solution time as problem size grows in proportion to the number of processors

Multigrid uses a sequence of coarse grids to accelerate the fine grid solution



Simple 1D model problem

• 1D Laplace on a uniform grid with spacing *h*



• Discrete problem is a linear system Au = f with

$$A = \begin{pmatrix} 2 & -1 & & \\ -1 & 2 & -1 & \\ & & \ddots & \\ & & -1 & 2 \end{pmatrix} \quad \text{or} \quad A = \begin{bmatrix} -1 & 2 & -1 \end{bmatrix}$$

Stencil
Matrix

We will mostly use the stencil form



Multigrid components for 1D model

- Many smoother options, e.g.,
 - Weighted Jacobi, $I (\omega/2) A$, $\omega_{opt} = 2/3$
 - Gauss-Seidel (GS)
- Prolongation is linear interpolation (note brackets)

 Coarse-grid operator is coarse discretization of the problem (scaled appropriately)

$$A_c = \frac{1}{2} \left[\begin{array}{ccc} -1 & 2 & -1 \end{array} \right]$$

 In practice, a slightly different method (equivalent to cyclic reduction) solves this problem in one V-cycle



2D model problem: Laplace on a square (1)

Five-point stencil discretization on a uniform grid

- Smoothers: weighted Jacobi or GS (lexicographical or red/black)
- Full coarsening, bilinear interpolation



• Coarse discretization (scaled appropriately) for *A_c*



2D anisotropic model problem on a square

Five-point stencil discretization on a uniform grid

$$-u_{xx} - \epsilon u_{yy} = f, \quad 0 < \epsilon \ll 1 \qquad A = \begin{bmatrix} -\epsilon \\ -1 & 2 + 2\epsilon & -1 \\ -\epsilon \end{bmatrix}$$

г

Pointwise relaxation smooths only in the x direction!

Two solutions:

- Change coarse-grid correction coarsen only in the direction of smoothness (semicoarsening in *x*, linear interpolation in *x*)
- 2) Change relaxation line relaxation with points grouped along *y* lines



Parallel Multigrid

(see Yang tutorial on Monday)



Approach for parallelizing multigrid is straightforward data decomposition



- Basic communication pattern is "nearest neighbor"
 - Relaxation, interpolation, & Galerkin not hard to implement
- Different neighbor processors on coarse grids
- Many idle processors on coarse grids (100K+ on BG/L)
 - Algorithms to take advantage have had limited success

Straightforward parallelization approach is optimal for V-cycles on structured grids (5-pt Laplacian example)

Standard communication / computation models

 $T_{comm} = \alpha + m\beta$ (communicate m doubles) $T_{comp} = m\gamma$ (compute m flops)

Time to do relaxation

 $T \approx 4\alpha + 4n\beta + 5n^2\gamma$



Time to do relaxation in a V(1,0) multigrid cycle

 $T_V \approx (1 + 1 + \dots) 4\alpha + (1 + 1/2 + \dots) 4n\beta + (1 + 1/4 + \dots) 5n^2\gamma$ $\approx (\log N) 4\alpha + (2) 4n\beta + (4/3) 5n^2\gamma$

- For achieving optimality in general, the *log* term is unavoidable!
- More precise: $T_{V,better} \approx T_V + (\log P)(4\beta + 5\gamma)$



Additional comments on parallel multigrid

W-cycles scale poorly:

 $T_W \approx (2^{\log N}) 4\alpha + (\log N) 4n\beta + (2) 5n^2 \gamma$



- Lexicographical Gauss-Seidel is too sequential
 - Use red/black or multi-color GS
 - Use weighted Jacobi, hybrid Jacobi/GS, L1
 - Use C-F relaxation (Jacobi on C-pts then F-pts)
 - Use Polynomial smoothers
- Parallel smoothers are often less effective



- Recent survey on parallel multigrid:
 - "A Survey of Parallelization Techniques for Multigrid Solvers," Chow, Falgout, Hu, Tuminaro, and Yang, Parallel Processing For Scientific Computing, Heroux, Raghavan, and Simon, editors, SIAM, series on Software, Environments, and Tools (2006)
- Recent paper on parallel smoothers:
 - "Multigrid Smoothers for Ultra-Parallel Computing," Baker, Falgout, Kolev, and Yang, SIAM J. Sci. Comput., submitted. LLNL-JRNL-435315

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Example weak scaling results on Dawn (an IBM BG/P system at LLNL) in 2010



- Laplacian on a cube; 40³ = 64K grid points per processor; largest problem had 3 billion unknowns!
- PFMG is a semicoarsening multigrid solver in *hypre*
- Still room to improve setup implementation (these results already employ the assumed partition algorithm described later)

Basic multigrid research challenge

- Optimal O(N) multigrid methods don't exist for some applications, even in serial
- Need to invent methods for these applications
- However …
- Some of the classical and most proven techniques used in multigrid methods don't parallelize
 - Gauss-Seidel smoothers are inherently sequential
 - W-cycles have poor parallel scaling
- Parallel computing imposes additional restrictions on multigrid algorithmic development

Algebraic Multigrid (AMG)

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Algebraic Multigrid (AMG) is based on MG principles, but uses matrix coefficients

For best results, geometry alone is not enough



 AMG ignores geometric information altogether, but captures both linear & operator-dep interpolation

$$(A\mathbf{u})_i = a_{i,i-1}u_{i-1} + a_{i,i}u_i + a_{i,i+1}u_{i+1}$$

$$u_{i} = \left(-\frac{a_{i,i-1}}{a_{i,i}}\right)u_{i-1} + \left(-\frac{a_{i,i+1}}{a_{i,i}}\right)u_{i+1}$$



AMG is an ideal method for unstructured grid problems

- Many algorithms (AMG alphabet soup)
- Automatically coarsens "grids"





Algebraically smooth error may not be smooth in a geometric sense

Error left by relaxation can be geometrically oscillatory

7 GS sweeps on

$$-au_{xx} - bu_{yy} = f$$

$$a = b$$
 $a \gg b$



- This example...
 - targets geometric smoothness
 - uses pointwise smoothers
- Not sufficient for some problems!



AMG coarsens grids in the direction of geometric smoothness



Preliminaries... the Galerkin coarse-grid operator

As before, consider solving the N×N linear system

$$Au = f$$

- Let P be prolongation (interpolation) and P^T restriction
- The coarse-grid operator is defined by the Galerkin procedure, $A_c = P^T A P$
- This gives the "best" coarse-grid correction in the sense that the solution e_c of the coarse system

$$A_c e_c = P^T r$$

satisfies

$$e_c = \arg \min \|e - Pe_c\|_A$$



Preliminaries... AMG "grids"

- Matrix adjacency graphs play an important role in AMG:
 - grid = set of graph vertices
 - grid point i = vertex i
- As a visual aid, it is highly instructive to relate the matrix equations to an underlying PDE and discretization
- We will often draw the grid points in their geometric locations
- Remember that AMG doesn't actually use this geometric information!



Classical AMG (C-AMG)

(Brandt, McCormick, Ruge, Stüben)

C-AMG targets geometric smoothness

 From theory (later): smooth error is characterized by small eigenmodes, hence satisfies (A scaled to have norm 1)

$$\mathbf{e}^T A \mathbf{e} = \lambda \ll \mathbf{1}$$

Constant is geometrically smooth, so assume zero row sum

$$e^{T}Ae = \sum_{i} e_{i} \left(a_{ii}e_{i} + \sum_{j \neq i} a_{ij}e_{j} \right)$$

$$= \sum_{i} e_{i} \left(\sum_{j \neq i} (-a_{ij})(e_{i} - e_{j}) \right)$$

$$a_{ii} = -\sum_{j \neq i} a_{ij}$$

$$= \sum_{i < j} (-a_{ij})e_{i}(e_{i} - e_{j}) + \sum_{i > j} (-a_{ij})e_{i}(e_{i} - e_{j})$$

$$= \sum_{i < j} (-a_{ij})e_{i}(e_{i} - e_{j}) - \sum_{i < j} (-a_{ji})e_{j}(e_{i} - e_{j})$$

$$= \sum_{i < j} (-a_{ij})(e_{i} - e_{j})^{2} \ll 1$$

C-AMG targets geometric smoothness through strength-of-connection

Assuming geometric smoothness, can show

$$\mathbf{e}^T A \mathbf{e} = \sum_{i < j} (-a_{ij}) (e_i - e_j)^2 \ll 1$$

- C-AMG Smoothness Heuristic: Smooth error varies slowly in the direction of "large" matrix coefficients
- Strength of connection: Given a threshold $0 < \theta \le 1$, we say that variable u_i strongly depends on variable u_i if

$$-a_{ij} \ge \theta \max_{k \neq i} \{-a_{ik}\}$$

- In practice, positive off-diagonals are weak
- Note that this definition of strength is not symmetric

Choosing the coarse grid

- In C-AMG, the coarse grid is a subset of the fine grid
- The basic coarsening procedure is as follows:
 - Define a strength matrix A_s by deleting weak connections in A
 - First pass: Choose an independent set of fine-grid points based on the graph of A_s
 - Second pass: Choose additional points if needed to satisfy interpolation requirements
- Coarsening partitions the grid into C- and F-points



- select C-pt with maximal measure
- select neighbors as F-pts
- Jupdate measures of F-pt neighbors



- → select C-pt with maximal measure
- select neighbors as F-pts
- update measures of F-pt neighbors





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C-AMG coarsening is inherently sequential



- select C-pt with maximal measure
- select neighbors as F-pts
- Jupdate measures of F-pt neighbors



C-AMG coarsening – second pass

- Recall: Second pass chooses additional points if needed to satisfy interpolation requirements
- C-AMG interpolation (discussed next) requires that each pair of strongly connected *F*-points be strongly connected to a common *C*-point
- C-AMG second pass: search for *F*-point pairs that don't satisfy this requirement and changes one to a *C*-point
- Second pass can lead to high complexity
- Idea: eliminate second pass, modify interpolation
AMG grid hierarchies for several 2D problems





C-AMG Interpolation – collapsing the stencil

Smooth error means "small" residuals

 $r \ll \|A\|$

To derive interpolation, assume that

$$r_i = (Ae)_i = 0$$

Hence,

$$\begin{aligned} a_{ii}e_i &= -\sum_{j \in C_i} a_{ij}e_j - \sum_{j \in F_i^s} a_{ij}e_j - \sum_{j \in N_i^w} a_{ij}e_j \\ C_i &: C\text{-points strongly connected to } i \end{aligned}$$

- F_i^s : F-points strongly connected to i
- N_i^w : all points weakly connected to *i*.
- The trick is to rewrite the e_j in F^s_i and N^w_i in terms of either the interpolatory points in C_i or the F-point i

C-AMG Interpolation – collapsing the stencil... Isotropic 9-pt Laplacian



C-AMG Interpolation – collapsing the stencil... Anisotropic 9-pt Laplacian, $\theta > 0.25$



C-AMG Interpolation – algebraic derivation

$$a_{ii}e_i = -\sum_{j \in C_i} a_{ij}e_j - \sum_{j \in F_i^s} a_{ij}e_j - \sum_{j \in N_i^w} a_{ij}e_j$$

- C_i : C-points strongly connected to i
- F_i^s : F-points strongly connected to i N_i^w : all points weakly connected to i.

Write

$$e_{j} = \begin{cases} \sum_{k \in C_{i}} \left(\frac{a_{jk}}{\sum_{l \in C_{i}} a_{jl}} \right) e_{k}, & j \in F_{i}^{s} \\ e_{i}, & j \in N_{i}^{w} \end{cases}$$

Then

$$w_{ij} = -\left(a_{ij} + \sum_{k \in F_i^s} \left(\frac{a_{kj}}{\sum_{l \in C_i} a_{kl}}\right) a_{ik}\right) / \left(a_{ii} + \sum_{k \in N_i^w} a_{ik}\right)$$

Example C-AMG results

$$-au_{xx} - bu_{yy} = f \quad \begin{vmatrix} a = b \\ a & b \end{vmatrix}$$



C-AMG coarse grids

		Conv	Coarse	Grid	Oper	Setup	Solve
N	Iters	factor	grids	comp	comp	time	time
61×61	10	0.23	6	1.6	1.6	0.01	0.02
121 × 121	9	0.23	8	1.6	1.7	0.05	0.07
241 × 241	9	0.23	9	1.6	1.7	0.25	0.32
481 × 481	9	0.23	12	1.7	1.7	1.02	1.27
961×961	11	0.29	13	1.7	1.7	4.42	6.28

- Grid complexity total # of grid points divided by the # of fine grid points
- Operator complexity total # of nonzeroes in the system operators divided by the # of nonzeroes in the fine grid operator

Parallel AMG

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Parallel Coarsening Algorithms

- C-AMG coarsening algorithm is inherently sequential
- Several parallel algorithms (in *hypre*):
 - CLJP (Cleary-Luby-Jones-Plassmann) one-pass approach with random numbers to get concurrency (illustrated next)
 - Falgout C-AMG on processor interior, then CLJP to finish
 - PMIS CLJP without the 'C'; parallel version of C-AMG first pass
 - HMIS C-AMG on processor interior, then PMIS to finish
 - CGC (Griebel, Metsch, Schweitzer) compute several coarse grids on each processor, then solve a global graph problem to select the grids with the best "fit"

• ...

Other parallel AMG codes use similar approaches





- select C-pts with maximal measure locally
- remove neighbor edges
- update neighbor measures



→ select C-pts with maximal measure locally

- remove neighbor edges
- update neighbor measures













10 C-points selected

Standard AMG selects 9 C-points





Parallel coarse-grid selection in AMG can produce unwanted side effects

- Non-uniform grids can lead to increased operator complexity and poor convergence
- Operator "stencil growth" reduces parallel efficiency



- Currently no guaranteed ways to control complexity
- Can ameliorate with more aggressive coarsening
- Requires long-range interpolation approaches

C-AMG interpolation is not suitable for more aggressive coarsening

 PMIS is parallel and eliminates the second pass, which can lead to the following scenarios:



- Want above *i*-points to interpolate from both C-points
- Long-range (distance two) interpolation!

One possibility for long-range interpolation is extended interpolation

- **C-AMG**: $C_i = \{j,k\}$
- Long-range: $C_i = \{j,k,m,n\}$
- Extended interpolation apply C-AMG interpolation to an extended stencil
- Extended+i interpolation is the same as extended, but also collapses to point i
- Improves overall quality



New parallel coarsening and long-range interpolation methods are improving scalability

- Unstructured 3D problem with material discontinuities
- About 90K unknowns per processor on MCR (Linux cluster)
- AMG GMRES(10)



Parallel AMG in hypre now scales to 130K processors on BG/L ... and beyond





- Largest problem above: 2B unknowns
- Largest problem to date: 26B unknowns on 98K processors of BG/L
- Most processors to date: 16B unknowns on 196K cores of Jaguar (Cray XT5 at ORNL)

Smoothed Aggregation (SA)

(Vaněk, Mandel, Brezina)



SA views the prolongation operator columnwise, as a set of local basis functions

ID Laplacian example:



- Range(P) contains the (smooth) constant vector 1
- SA approach for building prolongation decompose near null space into a basis with local support

SA builds interpolation by first chopping up a global basis, then smoothing it

 Tentative interpolation is constructed from "aggregates" (local QR factorization is used to orthonormalize)



 Smoothing adds basis overlap and improves approximation property

 $P = S\hat{P}$



Phase 1:

a) Pick root pt not adjacent to agg

b) Aggregate root and neighbors

Phase 2:





Phase 1:
a) Pick root pt not adjacent to agg
b) Aggregate root and neighbors

Phase 2:





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Phase 1:

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Phase 2:



SA coarsening is traditionally more aggressive than C-AMG coarsening (5-pt Laplacian example)



Operator complexities are usually smaller, too



Additional comments on SA...

- Usual prolongator smoother is damped Jacobi
- Strength of connection is usually defined differently $|a_{ij}| > \theta \sqrt{|a_{ii}a_{jj}|}$
- Special care must be taken for anisotropic problems to keep complexity low
 - Thresholded prolongator smoothing
 - Basis shifting approach
- Parallel SA coarsening has issues similar to C-AMG

AMG Theory & Compatible Relaxation


GAMG preliminaries...

- Consider solving Au = f, A SPD
- Consider smoothers with error propagation

$$e_{k+1} = (I - M^{-1}A)e_k$$

where we assume that $(M+M^T-A)$ is SPD (necessary and sufficient condition for convergence)

- Note: *M* may be symmetric or nonsymmetric
- Denote the symmetrized smoother operator by

$$\widetilde{M} = M(M^T + M - A)^{-1}M^T$$

that is, $(I - \widetilde{M}^{-1}A) = (I - M^{-T}A)(I - M^{-1}A)$



GAMG preliminaries continued...

- Let $P: \Re^{n_c} \to \Re^n$ be interpolation (prolongation)
- Let $R: \mathfrak{R}^n \to \mathfrak{R}^{n_c}$ be some "restriction" operator
 - Note that *R* is not the MG restriction operator
- Define s.t. RP=I and PR is a projection onto range(P)
- For any SPD matrix *X* and any full-rank matrix *B*, denote the *X*-orthogonal projection onto range(*B*) by $\pi_X(B) = B(B^T X B)^{-1} B^T X$
- Define the two-grid multigrid error propagator by $E_{TG} = (I - M^{-1}A)(I - \pi_A(P))$



GAMG two-grid theory splits construction of coarsegrid correction into two parts

• Theorem:

$$||E_{TG}||_A^2 \le 1 - \frac{1}{K}; \quad K = \sup_e \frac{||(I - PR)e||_{\widetilde{M}}^2}{||e||_A^2}$$

- Now, fix R so that it does not depend on P
 - Defines the coarse-grid variables, $u_c = Ru$
 - Example: R = [0, I] $(P^T = [W^T, I]^T)$, i.e., subset of the fine grid
- Theorem:

$$K \leq \eta K_{\star}; \quad \eta = \|PR\|_A; \quad K_{\star} = \inf_P \sup_e \frac{\|(I - PR)e\|_{\widetilde{M}}^2}{\|e\|_A^2}$$

- Small K_{\star} insures coarse grid quality use CR
- Small η insures interpolation quality necessary condition that does not depend on relaxation!

CR is an efficient method for measuring the quality of the set of coarse variables

- CR (Brandt, 2000) is a modified relaxation scheme that keeps the coarse-level variables, Ru, invariant
- Theorem: (fast convergence) good coarse grid)

$$K_{\star} \leq \left(\frac{\Delta^2}{2-\omega}\right) \frac{1}{1-\rho_{cr}}$$

 $\Delta \ge 1$ measures the deviation of *M* from its symmetric part M_{σ} and $0 < \omega < 2$ is a kind of smoothing parameter

$$\Delta^2 = \left\| M_{\sigma}^{-1/2} \ M \ M_{\sigma}^{-1/2} \right\|^2; \quad \omega = \lambda_{\max}(M_{\sigma}^{-1}A)$$

- Must insure "good" constants
 - in particular, $\omega \ll 2$



Several general CR methods

- Define *S* such that $\Re^n = \operatorname{range}(S) \oplus \operatorname{range}(R^T)$ and RS = 0
 - Example: $R = [0, I]; S = [I, 0]^T; P^T = [W^T, I]^T$
- Primary CR method feasible for relaxation based on matrix splittings, where M is explicitly available

$$I - (S^T M S)^{-1} (S^T A S)$$

Habituated CR – not as sharp, but always computable

$$S^T(I - M^{-1}A)S; \quad S^TS = I$$



Sharp Theory insightful for improving CR prediction

GAMG theory

$$||E_{TG}||_A^2 \le 1 - \frac{1}{K}; \quad K = \sup_e \frac{||(I - PR)e||_{\widetilde{M}}^2}{||e||_A^2}$$

Sharp theory

$$||E_{TG}||_A^2 = 1 - \frac{1}{K_{\sharp}}; \quad K_{\sharp} = \sup_e \frac{\left\| (I - \pi_{\widetilde{M}}(P))e \right\|_{\widetilde{M}}^2}{\|e\|_A^2}$$

- Differ only in form of the projection
- Careful comparison \rightarrow optimal *R* is given by $R = P^T \widetilde{M}$
- But, we don't have P yet (we're trying to build it)



AMG and ILU

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Ideal interpolation

• Recall 2-level theory:

$$||E_{TG}||_A^2 \le 1 - \frac{1}{K}; \qquad K = \sup_e \frac{||(I - PR)e||_{\widetilde{M}}^2}{||e||_A^2}$$

- Consider $R = [0, I_c], P^T = [W^T, I_c]^T$ case
- "Ideal" P is given by

$$P_{\star} = \arg\min_{P} \sup_{e} \frac{\|(I - PR)e\|_{\widetilde{M}}^{2}}{\|e\|_{A}^{2}} = \begin{bmatrix} -A_{ff}^{-1}A_{fc} \\ I_{c} \end{bmatrix}$$

Not a practical choice in general

AMG and ILU

• Can factor *A* as follows

$$A = \begin{bmatrix} A_{ff} & A_{fc} \\ A_{cf} & A_{cc} \end{bmatrix} = \begin{bmatrix} I & 0 \\ (A_{ff}^{-1}A_{fc})^T & I \end{bmatrix} \begin{bmatrix} A_{ff} & 0 \\ 0 & S_{cc} \end{bmatrix} \begin{bmatrix} I & (A_{ff}^{-1}A_{fc}) \\ 0 & I \end{bmatrix}$$
$$A^{-1} = \begin{bmatrix} I & -(A_{ff}^{-1}A_{fc}) \\ 0 & I \end{bmatrix} \begin{bmatrix} A_{ff}^{-1} & 0 \\ 0 & S_{cc}^{-1} \end{bmatrix} \begin{bmatrix} I & 0 \\ -(A_{ff}^{-1}A_{fc})^T & I \end{bmatrix}$$
$$= S(S^T A S)^{-1} S^T + P_{\star}(P_{\star}^T A P_{\star})^{-1} P_{\star}^T$$

Thinking of ILU, write the error propagator

$$(I - (A^{-1})A) = (I - S(S^T A S)^{-1} S^T A) (I - P_{\star} (P_{\star}^T A P_{\star})^{-1} P_{\star}^T A)$$

F-relaxation

Coarse-grid correction

AMG for Electromagnetic Problems

(see Kolev Poster on Monday)



Electromagnetic (EM) problems have huge oscillatory near null spaces

- Definite Maxwell, Indefinite Maxwell, Helmholtz
- Require specialized smoothers and coarse grids



Local: specialized relaxation (Definite / Indefinite Maxwell)



Global: specialized coarse grids (Helmholtz, Indefinite Maxwell)

Definite Maxwell, Nédélec edge FEM discretization

 $\nabla \times \alpha \nabla \times \boldsymbol{E} + \beta \boldsymbol{E} = f \quad \alpha, \beta > 0$

Near null-space characterized by gradients

 $\nabla \times (\nabla p_h) = 0$

Geometric multigrid for definite Maxwell

Helmholtz decomposition

 $E_{h} = v_{h} + \nabla p_{h}$ divergence-free curl-free • Smooth both components (Hiptmair, SINUM 1998) $R_{h} = R_{e,h} + G_{h}R_{v,h}G_{h}^{T}$ Discrete Gradient Point smoother for Point smoother for A_{h} $G_{h}^{T}A_{h}G_{h}$

- Block smoother (Arnold, Falk, Winther, Num. Math. 2000)
- Natural FE interpolation
- Difficulties extending to
 - unstructured meshes
 - variable coefficients



Auxiliary-space Maxwell solver (AMS) utilizes a new decomposition

Based on Hiptmair, Xu (2006)

$$\boldsymbol{E}_h = \boldsymbol{v}_h + \nabla p_h + \boldsymbol{\Pi}_h \boldsymbol{z}_h$$

Define preconditioner based on nodal solvers

$$B_{h} = R_{h} + G_{h}B_{v,h}G_{h}^{T} + \Pi_{h}B_{v,h}\Pi_{h}^{T}$$
Point smoother for AMG solver for AMG solver for A_{h} $G_{h}^{T}A_{h}G_{h}$ $\Pi_{h}^{T}A_{h}\Pi_{h}$



- User provides A, G_h and vertex coordinates
- Fast computation of Π_h (~ 3 mat-vec multiplies)
- AMS is a variational form of Hiptmair-Xu

Auxiliary-space Maxwell Solver (AMS) is improving solve times by up to 25x for some EM problems

- Hiptmair-Xu / AMS are the first provably scalable solvers for EM on unstructured grids
- Employs BoomerAMG
- Highly robust
 - Materials with widely varying electromagnetic properties
 - Unstructured grids
- Example: 1.2B unknowns on 1.9K processors took 355s (23 iterations)



Adaptive AMG

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Adaptive AMG is well-suited for QCD

- Quantum Chromodynamics (QCD) is the theory of strong forces in the Standard Model of particle physics
- Scalable solvers for the Dirac equations have been elusive until recently
- Challenges:
 - The system is complex and indefinite
 - The system can be extremely ill-conditioned
 - Near null space is unknown and oscillatory!









Adaptive AMG idea: use the method to improve the method

- Requires no a-priori knowledge of the near null space
- Idea: uncover *representatives* of slowly-converging error by applying the "current method" to Ax = 0, then use these to adapt (improve) the method
- Achi Brandt's *Bootstrap AMG* is an adaptive method
- PCG can be viewed as an adaptive method
 - Not optimal because it uses a global view
 - The key is to view representatives locally
- We developed 2 methods: αAMG and αSA (SISC pubs)



To build effective interpolation, it is important to interpret the near null space in a local way

- (2-level) Coarse-grid correction is a projection $(I - P(P^TAP)^{-1}P^TA)e$
- Better to break up near null space into a local basis



Deflation – not optimal

Multigrid – optimal

 Get full approximation property (low-frequency Fourier modes in this example)

SA builds interpolation by first chopping up a global basis, then smoothing it

 Tentative interpolation is constructed from "aggregates" (local QR factorization is used to orthonormalize)



 Smoothing adds basis overlap and improves approximation property

 $P = S\hat{P}$

Adaptive smoothed aggregation (α SA) automatically builds the global basis for SA

- Generate the basis one vector at a time
 - Start with relaxation on $Au=0 \rightarrow u_1 \rightarrow \alpha SA(u_1)$
 - Use α SA(u_1) on $Au=0 \rightarrow u_2 \rightarrow \alpha$ SA(u_1, u_2)
 - Etc., until we have a good method
- Setup is expensive, but is amortized over many RHS's
- Published in 2004, highlighted in SIAM Review in 2005
 - Brezina, Falgout, MacLachlan, Manteuffel, McCormick, and Ruge, "Adaptive smoothed aggregation (αSA)," SIAM J. Sci. Comput. (2004)
- Successfully applied to 2D QED
 - Brannick, Brezina, Keyes, Livne, Livshits, MacLachlan, Manteuffel, McCormick, Ruge, and Zikatanov, "Adaptive smoothed aggregation in lattice QCD," Springer (2006)

4D Wilson-Dirac Results: D-MG shows no critical slowing down (Time)



• Parameters: N=16³x32, β =6.0, m_{crit} = -0.8049

D-MG Parameters: 4⁴x3x2 blocking, 3 levels, W(2,2,4) cycle, N_v = 20, setup run at m_{crit}

Summary

- Multigrid methods are optimal and have good scaling potential
- AMG is based primarily on matrix entries
- In practice, some additional properties of the underlying system are assumed (near null space)
- Adaptive AMG uncovers near null space information
- AMG can solve a large class of problems and can scale to BG/Lclass machines
- Parallel computing imposes additional restrictions on AMG algorithmic development
- Still many outstanding research questions

The Scalable Linear Solvers Team



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See http://www.llnl.gov/casc/linear_solvers for publications, presentations, and software (hypre)

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