Abstract Syntax Trees
&
Top-Down Parsing

Review of Parsing

- Given a language L(G), a parser consumes a sequence of tokens s and produces a parse tree
- Issues:
  - How do we recognize that $s \in L(G)$?
  - A parse tree of s describes how $s \in L(G)$
  - Ambiguity: more than one parse tree (possible interpretation) for some string s
  - Error: no parse tree for some string s
  - How do we construct the parse tree?

Abstract Syntax Trees

- So far, a parser traces the derivation of a sequence of tokens
- The rest of the compiler needs a structural representation of the program
- Abstract syntax trees
  - Like parse trees but ignore some details
  - Abbreviated as AST

Abstract Syntax Trees (Cont.)

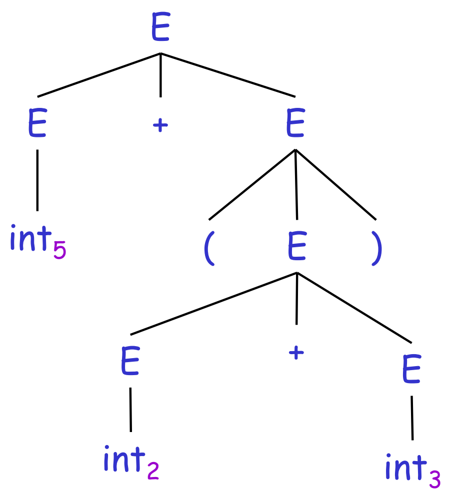
- Consider the grammar

  $E \rightarrow \text{int} \mid (E) \mid E + E$
- And the string

  5 + (2 + 3)
- After lexical analysis (a list of tokens)

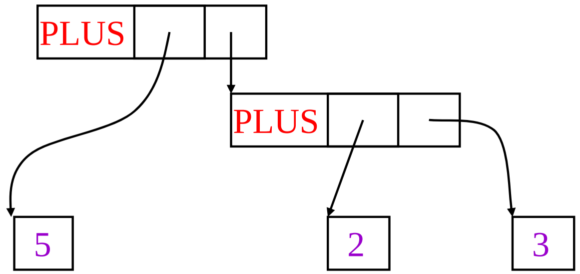
  int_5 '+' '(' int_2 '+' int_3 ')'
- During parsing we build a parse tree …

Example of Parse Tree



- Traces the operation of the parser
- Captures the nesting structure
- But too much info
  - Parentheses
  - Single-successor nodes

Example of Abstract Syntax Tree



- Also captures the nesting structure
- But <u>abstracts</u> from the concrete syntax
  $\mapsto$ more compact and easier to use
- An important data structure in a compiler

Semantic Actions

- This is what we'll use to construct ASTs
- Each grammar symbol may have <u>attributes</u>
  - An attribute is a property of a programming language construct
  - For terminal symbols (lexical tokens) attributes can be calculated by the lexer
- Each production may have an <u>action</u>
  - Written as: $X \to Y_1 \ldots Y_n$ { action }
  - That can refer to or compute symbol attributes

Semantic Actions: An Example

- Consider the grammar

$$E \to \text{int} \mid E + E \mid (E)$$

- For each symbol X define an attribute X.val
  - For terminals, val is the associated lexeme
  - For non-terminals, val is the expression's value (which is computed from values of subexpressions)
- We annotate the grammar with actions:

$$\begin{array}{lll} E \to & \text{int} & \{ E.\text{val} = \text{int.val} \} \\ \mid & E_1 + E_2 & \{ E.\text{val} = E_1.\text{val} + E_2.\text{val} \} \\ \mid & (E_1) & \{ E.\text{val} = E_1.\text{val} \} \end{array}$$

Semantic Actions: An Example (Cont.)

- String: 5 + (2 + 3)
- Tokens: int_5 '+' '(' int_2 '+' int_3 ')'

Productions	Equations
$E \rightarrow E_1 + E_2$	$E.val = E_1.val + E_2.val$
$E_1 \rightarrow int_5$	$E_1.val = int_5.val = 5$
$E_2 \rightarrow (E_3)$	$E_2.val = E_3.val$
$E_3 \rightarrow E_4 + E_5$	$E_3.val = E_4.val + E_5.val$
$E_4 \rightarrow int_2$	$E_4.val = int_2.val = 2$
$E_5 \rightarrow int_3$	$E_5.val = int_3.val = 3$

Semantic Actions: Dependencies

Semantic actions specify a system of equations

- Order of executing the actions is not specified

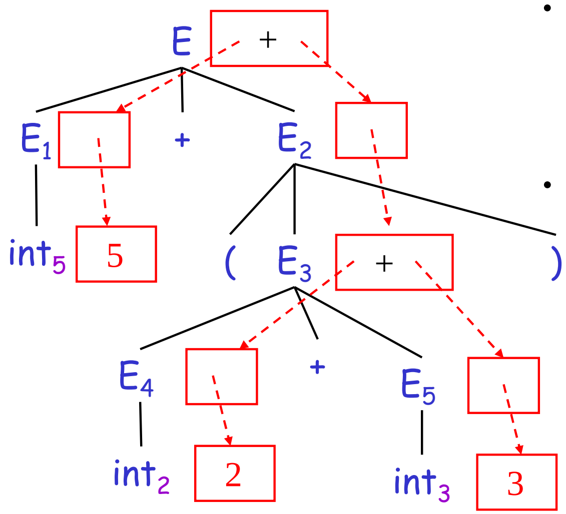
- Example:

$$E_3.val = E_4.val + E_5.val$$

  - Must compute $E_4.val$ and $E_5.val$ before $E_3.val$
  - We say that $E_3.val$ *depends on* $E_4.val$ and $E_5.val$

- The parser must find the order of evaluation

Dependency Graph



- Each node labeled with a non-terminal E has one slot for its val attribute
- Note the dependencies

Evaluating Attributes

- An attribute must be computed after all its successors in the dependency graph have been computed
  - In the previous example attributes can be computed bottom-up

- Such an order exists when there are no cycles
  - Cyclically defined attributes are not legal

Semantic Actions: Notes (Cont.)

- Synthesized attributes
  - Calculated from attributes of descendents in the parse tree
  - E.val is a synthesized attribute
  - Can always be calculated in a bottom-up order
- Grammars with only synthesized attributes are called S-attributed grammars
  - Most frequent kinds of grammars

Inherited Attributes

- Another kind of attributes
- Calculated from attributes of the parent node(s) and/or siblings in the parse tree
- Example: a line calculator

A Line Calculator

- Each line contains an expression

  $E \to int \;|\; E + E$
- Each line is terminated with the = sign

  $L \to E = \;|\; + E =$
- In the second form, the value of evaluation of the previous line is used as starting value
- A program is a sequence of lines

  $P \to \varepsilon \;|\; P\, L$

Attributes for the Line Calculator

- Each E has a synthesized attribute val
  - Calculated as before
- Each L has a synthesized attribute val

  L → E = { L.val = E.val }

  | + E = { L.val = E.val + L.prev }
- We need the value of the previous line
- We use an inherited attribute L.prev

Attributes for the Line Calculator (Cont.)

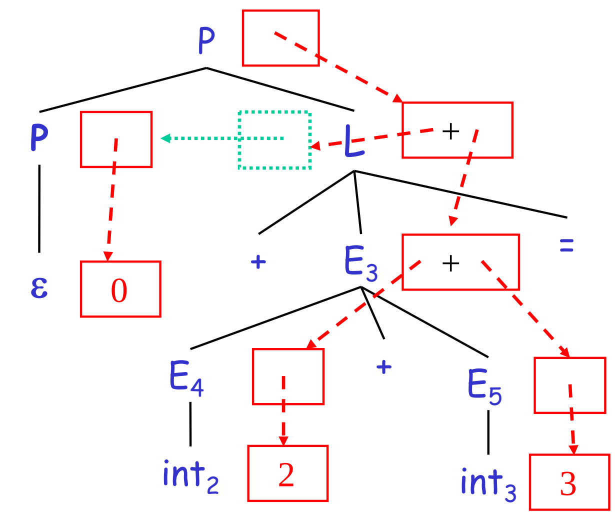
- Each P has a synthesized attribute val
  - The value of its last line

$P \to \varepsilon$ { P.val = 0 }
 | P_1 L { P.val = L.val;
 L.prev = P_1.val }

- Each L has an inherited attribute prev
  - L.prev is inherited from sibling P_1.val

- Example …

Example of Inherited Attributes



- val synthesized
- prev inherited
- All can be computed in depth-first order

Semantic Actions: Notes (Cont.)

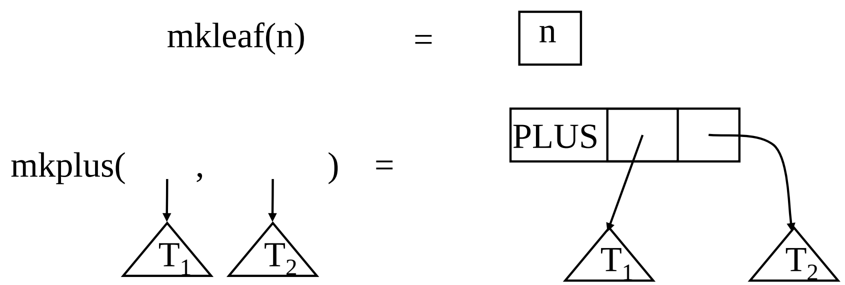
- Semantic actions can be used to build ASTs

- And many other things as well
  - Also used for type checking, code generation, …

- Process is called <u>syntax-directed translation</u>
  - Substantial generalization over CFGs

Constructing an AST

- We first define the AST data type
- Consider an abstract tree type with two constructors:



Constructing a Parse Tree

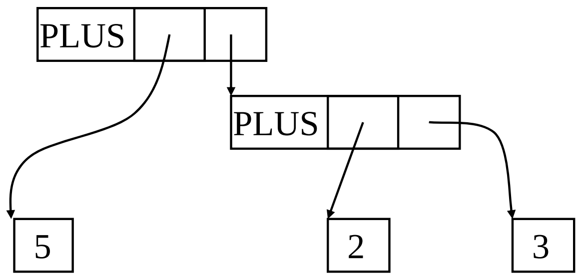
- We define a synthesized attribute ast
  - Values of ast values are ASTs
  - We assume that int.lexval is the value of the integer lexeme
  - Computed using semantic actions

$E \rightarrow$ int	{ E.ast = mkleaf(int.lexval) }
$\mid E_1 + E_2$	{ E.ast = mkplus(E_1.ast, E_2.ast) }
$\mid (E_1)$	{ E.ast = E_1.ast }

Parse Tree Example

- Consider the string int_5 '+' '(' int_2 '+' int_3 ')'
- A bottom-up evaluation of the ast attribute:

  E.ast = mkplus(mkleaf(5),
  mkplus(mkleaf(2), mkleaf(3))



Review of Abstract Syntax Trees

- We can specify language syntax using CFG
- A parser will answer whether $s \in L(G)$
- … and will build a parse tree
- … which we convert to an AST
- … and pass on to the rest of the compiler

- Next two & a half lectures:
  - How do we answer $s \in L(G)$ and build a parse tree?
- After that: from AST to assembly language

Second-Half of Lecture 5: Outline

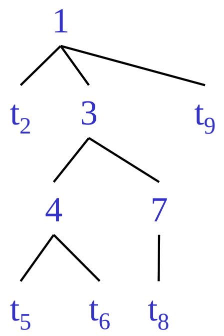
- Implementation of parsers
- Two approaches
  - Top-down
  - Bottom-up
- Today: Top-Down
  - Easier to understand and program manually
- Then: Bottom-Up
  - More powerful and used by most parser generators

Introduction to Top-Down Parsing

- Terminals are seen in order of appearance in the token stream:

  $t_2 \ t_5 \ t_6 \ t_8 \ t_9$

- The parse tree is constructed
  - From the top
  - From left to right



Recursive Descent Parsing

- Consider the grammar

  $E \to T + E \mid T$

  $T \to int \mid int * T \mid (E)$
- Token stream is: $int_5 * int_2$
- Start with top-level non-terminal E

- Try the rules for E in order

Recursive Descent Parsing. Example (Cont.)

Token stream: int5 * int2

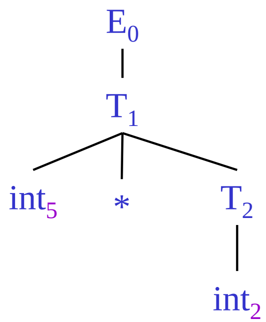
- Try $E_0 \to T_1 + E_2$
- Then try a rule for $T_1 \to (E_3)$
  - But (does not match input token int_5
- Try $T_1 \to int$. Token matches.
  - But + after T_1 does not match input token *
- Try $T_1 \to int * T_2$
  - This will match but + after T_1 will be unmatched
- Has exhausted the choices for T_1
  - Backtrack to choice for E_0

$E \to T + E \mid T$
$T \to (E) \mid int \mid int * T$

Recursive Descent Parsing. Example (Cont.)

Token stream: int5 * int2

- Try $E_0 \to T_1$
- Follow same steps as before for T_1
  - And succeed with $T_1 \to int_5 * T_2$ and $T_2 \to int_2$
  - With the following parse tree



$E \to T + E \mid T$
$T \to (E) \mid int \mid int * T$

Recursive Descent Parsing. Notes.

- Easy to implement by hand
- Somewhat inefficient (due to backtracking)
- But does not always work …

When Recursive Descent Does Not Work

- Consider a production $S \to S\ a$

  bool S_1() { return S() && term(a); }

  bool S() { return S_1(); }

- S() will get into an infinite loop
- A <u>left-recursive grammar</u> has a non-terminal S

  $S \to^{+} S\alpha$ for some α

- Recursive descent does not work in such cases

Elimination of Left Recursion

- Consider the left-recursive grammar

  $S \to S\ \alpha \mid \beta$

- S generates all strings starting with a β and followed by any number of α's
- The grammar can be rewritten using right-recursion

  $S \to \beta\ S'$

  $S' \to \alpha\ S' \mid \varepsilon$

More Elimination of Left-Recursion

- In general

  $S \to S\ \alpha_1 \mid \dots \mid S\ \alpha_n \mid \beta_1 \mid \dots \mid \beta_m$

- All strings derived from S start with one of $\beta_1,\dots,\beta_m$ and continue with several instances of $\alpha_1,\dots,\alpha_n$
- Rewrite as

  $S \to \beta_1\ S' \mid \dots \mid \beta_m\ S'$

  $S' \to \alpha_1\ S' \mid \dots \mid \alpha_n\ S' \mid \varepsilon$

General Left Recursion

- The grammar

$$S \to A\ \alpha \mid \delta$$
$$A \to S\ \beta$$

is also left-recursive because

$$S \to^{+} S\ \beta\ \alpha$$

- This left-recursion can also be eliminated
- See a Compilers book for a general algorithm

Summary of Recursive Descent

- Simple and general parsing strategy
  - Left-recursion must be eliminated first
  - … but that can be done automatically
- Unpopular because of backtracking
  - Thought to be too inefficient

- In practice, backtracking is eliminated by restricting the grammar

Predictive Parsers

- Like recursive-descent but parser can "predict" which production to use
  - By looking at the next few tokens
  - No backtracking
- Predictive parsers accept LL(k) grammars
  - L means "left-to-right" scan of input
  - L means "leftmost derivation"
  - k means "predict based on k tokens of lookahead"
- In practice, LL(1) is used

LL(1) Languages

- In recursive-descent, for each non-terminal and input token there may be a choice of production
- LL(1) means that for each non-terminal and token there is only one production
- Can be specified via 2D tables
  - One dimension for current non-terminal to expand
  - One dimension for next token
  - A table entry contains one production

Predictive Parsing and Left Factoring

- Recall the grammar for arithmetic expressions

  E → T + E | T

  T → (E) | int | int * T

- Hard to predict because
  - For T two productions start with int
  - For E it is not clear how to predict

- A grammar must be <u>left-factored</u> before it is used for predictive parsing

Left-Factoring Example

- Recall the grammar

  E → T + E | T

  T → (E) | int | int * T

- Factor out common prefixes of productions

  E → T X

  X → + E | ε

  T → (E) | int Y

  Y → * T | ε

LL(1) Parsing Table Example

- Left-factored grammar

  E → T X　　　X → + E | ε

  T → (E) | int Y　　　Y → * T | ε

- The LL(1) parsing table:

	int	*	+	(	)	$
E	T X			T X		
X			+ E		ε	ε
T	int Y			(E)		
Y		* T	ε		ε	ε

LL(1) Parsing Table Example (Cont.)

- Consider the [E, int] entry
  - "When current non-terminal is E and next input is int, use production E → T X
  - This production can generate an int in the first place
- Consider the [Y,+] entry
  - "When current non-terminal is Y and current token is +, get rid of Y"
  - Y can be followed by + only in a derivation in which Y → ε

LL(1) Parsing Tables: Errors

- Blank entries indicate error situations
  - Consider the [E,*] entry
  - "There is no way to derive a string starting with * from non-terminal E"

Using Parsing Tables

- Method similar to recursive descent, except
  - For each non-terminal S
  - We look at the next token a
  - And chose the production shown at [S,a]
- We use a stack to keep track of pending non-terminals
- We reject when we encounter an error state
- We accept when we encounter end-of-input

LL(1) Parsing Algorithm

```
initialize stack = <S $> and next
repeat
  case stack of
    <X, rest>  : if T[X,*next] = Y1…Yn
                     then stack ← <Y1…Yn rest>;
                     else  error();
    <t, rest>  : if t == *next++
                     then stack ← <rest>;
                     else  error();
until stack == <>
```


LL(1) Parsing Example

Stack	Input	Action
E $	int * int $	T X
T X $	int * int $	int Y
int Y X $	int * int $	terminal
Y X $	* int $	* T
* T X $	* int $	terminal
T X $	int $	int Y
int Y X $	int $	terminal
Y X $	$	ε
X $	$	ε
$	$	ACCEPT

	int	*	+	(	)	$
E	T X			T X		
X			+ E		ε	ε
T	int Y			(E)		
Y		* T	ε		ε	ε

Constructing Parsing Tables

- LL(1) languages are those defined by a parsing table for the LL(1) algorithm
- No table entry can be multiply defined

- We want to generate parsing tables from CFG

Constructing Parsing Tables (Cont.)

- If $A \to \alpha$, where in the line of A we place α ?
- In the column of t where t can start a string derived from α
  - $\alpha \to^* t\ \beta$
  - We say that $t \in First(\alpha)$
- In the column of t if α is ε and t can follow an A
  - $S \to^* \beta\ A\ t\ \delta$
  - We say $t \in Follow(A)$

Computing First Sets

Definition

$$First(X) = \{ t \mid X \to^* t\alpha \} \cup \{ \varepsilon \mid X \to^* \varepsilon \}$$

Algorithm sketch

1. $First(t) = \{ t \}$
2. $\varepsilon \in First(X)$ if $X \to \varepsilon$ is a production
3. $\varepsilon \in First(X)$ if $X \to A_1 \dots A_n$ and $\varepsilon \in First(A_i)$ for each $1 \le i \le n$
4. $First(\alpha) \subseteq First(X)$ if $X \to A_1 \dots A_n\ \alpha$ and $\varepsilon \in First(A_i)$ for each $1 \le i \le n$

First Sets: Example

- Recall the grammar

$E \to T\ X$ $\qquad$ $X \to +\ E \mid \varepsilon$

$T \to (\ E\ ) \mid int\ Y$ $\qquad$ $Y \to *\ T \mid \varepsilon$

- First sets

First(( ) = { (} $\qquad$ First( )) = {) }

First(+) = { + } $\qquad$ First(*) = { * }

First(int) = { int }

First(T) = { int, (}

First(E) = { int, (}

First(X) = { +, ε }

First(Y) = { *, ε }

Computing Follow Sets

- **Definition**

$$\text{Follow}(X) = \{ t \mid S \to^* \beta\, X\, t\, \delta \}$$

- **Intuition**
  - If $X \to A\ B$ then $\text{First}(B) \subseteq \text{Follow}(A)$ and $\text{Follow}(X) \subseteq \text{Follow}(B)$
  - Also if $B \to^* \varepsilon$ then $\text{Follow}(X) \subseteq \text{Follow}(A)$
  - If S is the start symbol then $\$ \in \text{Follow}(S)$

Computing Follow Sets (Cont.)

Algorithm sketch

1. $\$ \in \text{Follow}(S)$
2. $\text{First}(\beta) - \{\varepsilon\} \subseteq \text{Follow}(X)$

   For each production $A \to \alpha\ X\ \beta$
3. $\text{Follow}(A) \subseteq \text{Follow}(X)$

   For each production $A \to \alpha\ X\ \beta$ where $\varepsilon \in \text{First}(\beta)$

Follow Sets: Example

- Recall the grammar

$E \to T\ X$ $\qquad X \to +\ E \mid \varepsilon$

$T \to (\ E\ ) \mid \text{int}\ Y$ $\qquad Y \to *\ T \mid \varepsilon$

- Follow sets

Follow(+) = { int, (} Follow(*) = { int, (}

Follow(( ) = { int, (} Follow(E) = {), $ }

Follow(X) = { $,) } Follow(T) = { +,) , $ }

Follow( )) = { +,) , $ } Follow(Y) = { +,) , $ }

Follow(int) = { *, +,) , $ }

Constructing LL(1) Parsing Tables

- Construct a parsing table T for CFG G

- For each production $A \to \alpha$ in G do:
  - For each terminal $t \in \text{First}(\alpha)$ do
    - $T[A, t] = \alpha$
  - If $\varepsilon \in \text{First}(\alpha)$, for each $t \in \text{Follow}(A)$ do
    - $T[A, t] = \alpha$
  - If $\varepsilon \in \text{First}(\alpha)$ and $\$ \in \text{Follow}(A)$ do
    - $T[A, \$] = \alpha$

Notes on LL(1) Parsing Tables

- If any entry is multiply defined then *G* is not LL(1)
  - If *G* is ambiguous
  - If *G* is left recursive
  - If *G* is not left-factored
  - And in other cases as well
- Most programming language grammars are not LL(1)
- There are tools that build LL(1) tables

Review

- For some grammars there is a simple parsing strategy

  Predictive parsing

- Next time: a more powerful parsing strategy