

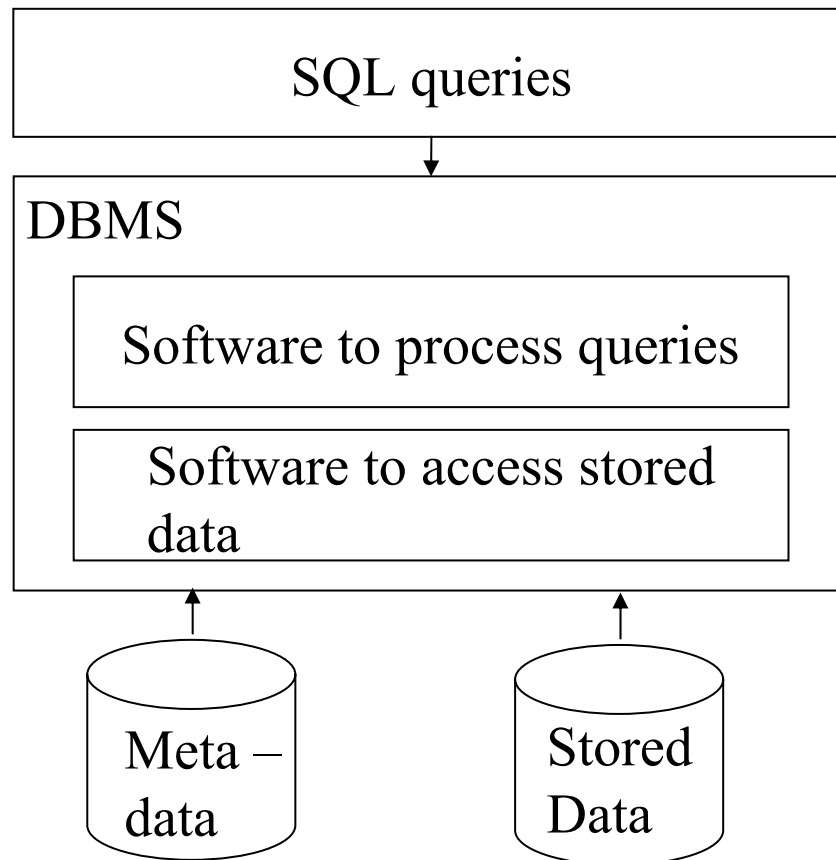
# Data Stream Management Systems

Principles of Modern Database Systems  
2007

Tore Risch  
Dept. of information technology  
Uppsala University  
Sweden

# What is a Data Base Management System?

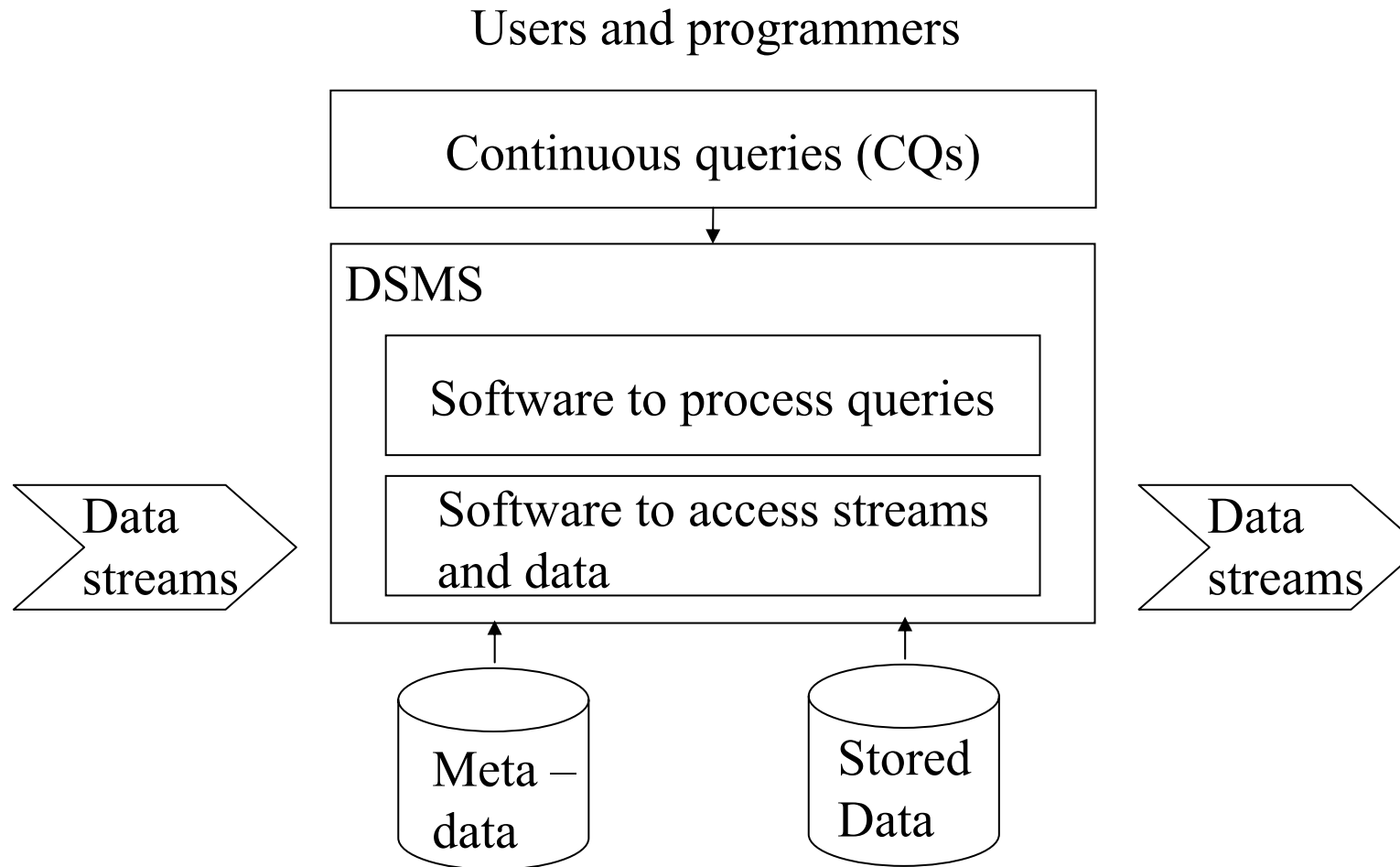
Users and programmers



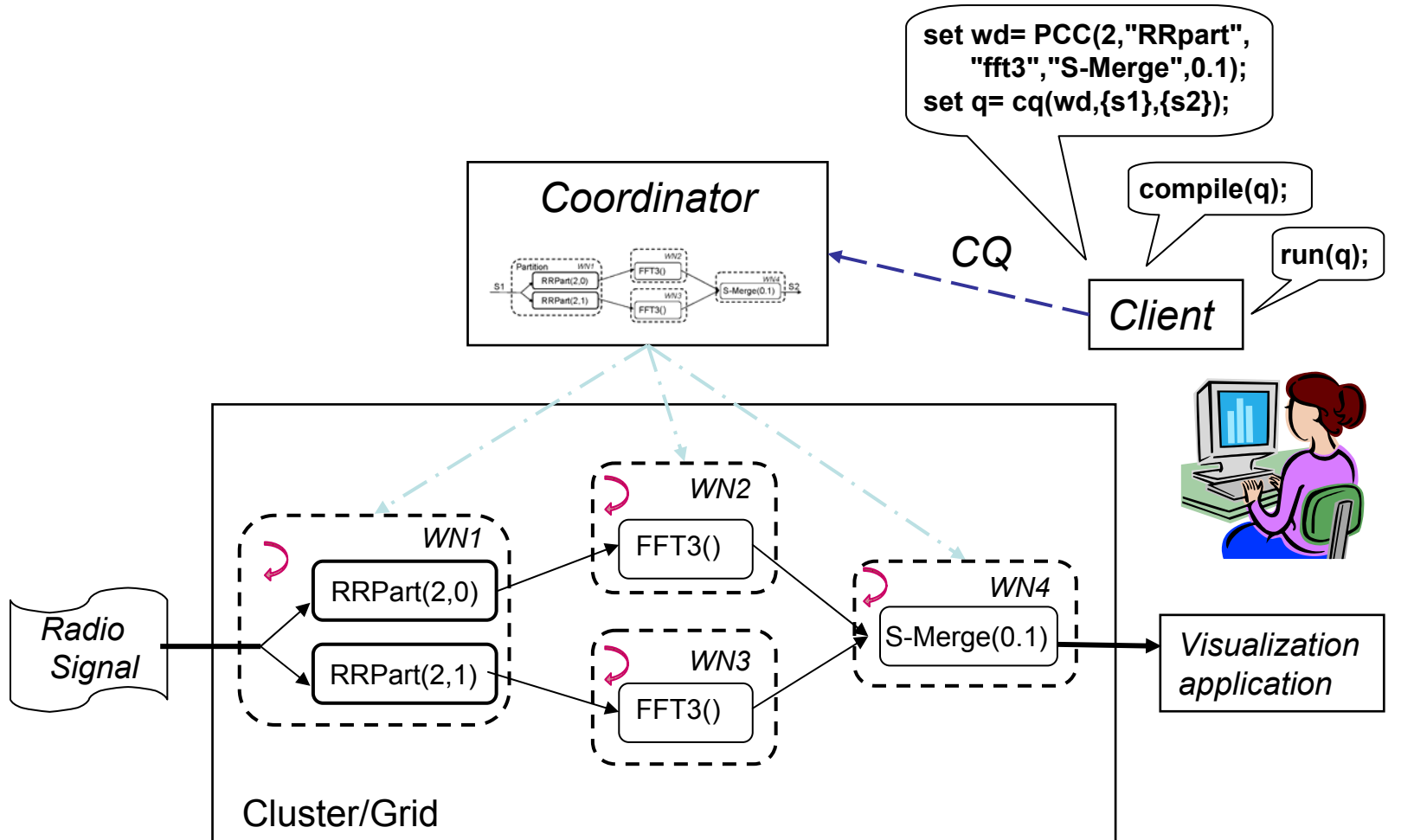
# New applications

- Data comes as large data streams, e.g.
    - Satellite data
    - Scientific instruments
    - Colliders
    - Patient monitoring
    - Stock data
    - Process industry
    - Traffic control
- ⇒ Would like to query data in *streams*

# What is a Data Stream Management System?



# DSMS Scenario

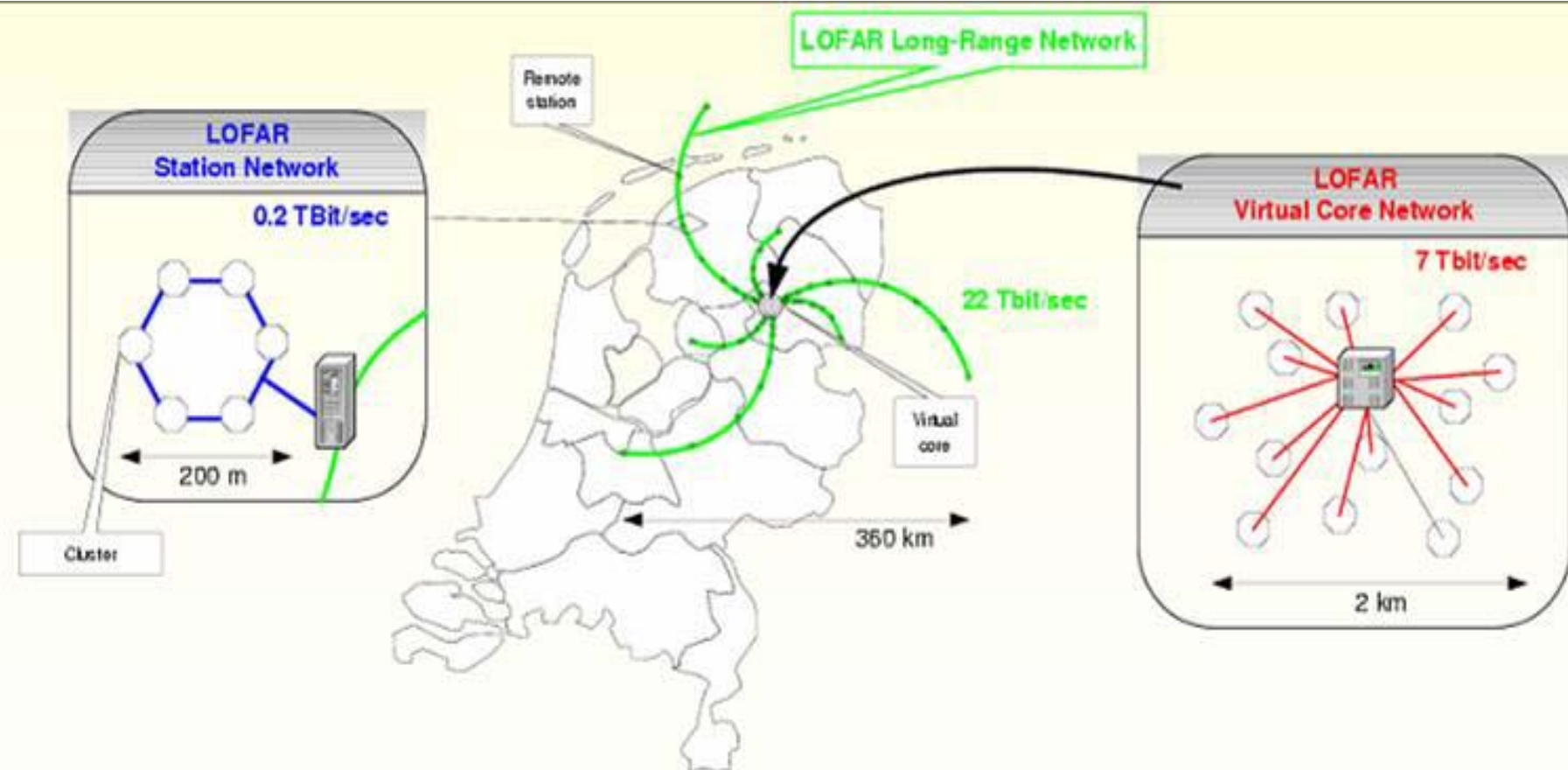


# Overview paper

⇒ L. Golab and T. Özsu: Issues in Stream Data Management, *SIGMOD Records*, 32(2), June 2003,  
<http://www.acm.org/sigmod/record/issues/0306/1.golab-ozsu1.p>

# The LOFAR Instrument

- 13000 antennas
- Distributed over 100 stations
- Producing ~20Tbps raw data



**U: Developing a scalable DSMS to process LOFAR stream queries**

# Streams vs tables

- Streams potentially *infinite* in size
  - Regular DBs based on queries to finite tables
- Streams ordered, i.e. *sequence* data
  - Regular DBs are based on sets and bags
- *Stop condition* indicates when/if streams end
- Often very high stream data volume and rate
  - Regular DBs usually less demanding
- Real-time delivery, Quality of Service
  - Regular DBs weak here
- Active query model, *continuous* queries
  - Regular DB queries *passive*



# Continuous queries

- CQs are turned on and run until stop condition true
  - Regular queries executed until finished by demand
- CQs return *unbounded* data (streams) as result
  - Regular queries *bounded* by size of tables
- CQs operators usually *montone*, i.e. cannot re-read stream
  - Regular queries can access same table many times
- CQs specified over *stream windows* (i.e. bounded stream segments)
  - Regular queries specified over entire tables
- CQs often based on time stamps (logs) of stream elements (*temporal*)
  - Regular queries not temporal
- CQ join operators *approximate*
  - Regular join operators usually exactly match data

# Stream windows

- Need monotone *window operator* to chop stream into segments
- Window *size* (*sz*) based on:
  - Number of elements  
E.g. last 10 elements
  - Time  
E.g. elements last second
- *Landmark* window:
  - Window from start of stream
  - Continuously growing
  - Not bounded
  - Materialization
- Windows also have *stride* (*str*)
  - Rule for how they move forward

# Window stride

- How fast the window moves forward
- *Jumping* window  
 $sz = str$ 
  - => Output data rate  $o$  = input data rate  $i$
  - => No overlap between windows
  - => All data processed once
  - => C.f. "window rate"  $wr=i/sz$
- *Sliding* windows  
 $str < sz$ 
  - =>  $o > i$  ( $o = i*sz/str$ )
  - => Overlaps between windows
  - => Data processed more than once
- *Sampling* window  
 $str > sz$ 
  - =>  $o < i$
  - => No overlaps
  - => Some data not processed
  - => a form of *schredding*

# Joining streams

- Streams infinite
  - => Monotone join operators needed
  - => regular join impossible (not monotone)
- Instead streams are *merged*:
  1. Split stream into segments by *window operator*
  2. Join windows from each stream
  3. Merge the result
- Stream merge is *approximate* join method
  - Window size determines quality of result
- Stream joins need to deal with rate differences, blocking
  - => *Time-out* when data blocks
  - => *Load shredding* skips stream elements
  - => Can also do *approximations* (e.g. aggregation)
  - => Need to deal with nulls (c.f. outer joins)

# Stream joining methods

- Special join methods different from table joins
- Xjoin:  
T. Urhan and M. Franklin. Dynamic pipeline scheduling for improving interactive performance of online queries. *Proceedings of the VLDB Conference*, 2001.
- Mjoin:  
S. Viglas, J. Naughton, and J. Burger. Maximizing the output rate of multi-join queries over streaming information sources. In *Proc. of the VLDB Conference 2003*
- Hybride:  
Babu, Munagala, Widom, Motwani: Adaptive Caching for Continuous Queries, *Proc. 21st International Conference on Data Engineering (ICDE 2005)*

# Punctuations

- Can be seen as corresponding to transactions
- Condition for a unit of work  
E.g. deal is done => new data about it ignored
- Add *punctuation* token in stream
- May improve performance
- Synchronization
- Punctuated joins:

Ding, Mehta, Rundensteiner, Heineman: Joining Punctuated Streams, *EDBT 2004*

# DSMS Systems

- Aurora* (Brown, MIT, Brandeis): Carney et al: Monitoring Streams – A New Class of Data Management Applications, *VLDB 2003*
- TelegraphCQ* (Berkeley): Chandrasekaran et al: TelegraphCQ: Continuous Dataflow Processing for an Uncertain World, *CIDR 2003*
- Gigascop*e (AT & T): Cranor et al: Gigascop: High Performance Network Monitoring with an SQL Interface, *SIGMOD 2002*
- STREAM* (Stanford): StreaMon: Baby & Widom: An Adaptive Engine for Stream Query Processing, *SIGMOD 2004*
- Borealis* (Brown & Brandeis): Ahmad et al: StreaMon: An Adaptive Engine for Stream Query Processing, *SIGMOD 2005* (distributed streams)
- Wavescope* (MIT): Girod et al: The Case for a Signal-Oriented Data Stream Management System, *CIDR 2007*

# Own related efforts

SCSQ (Zeitler & Risch): Processing high-volume stream queries on a supercomputer, ICDE Ph.D. Workshop 2006 (distributed, numerical)

*GSDM* (Ivanova & Risch): Customizable Parallel Execution of Scientific Stream Queries, VLDB 2005 (distributed, numerical)

L.Lin, T. Risch: Querying Continuous Time Sequences , VLDB 1998 (numerical time series)



# Aggregation over stream windows

E.g. SCSQ:

```
select avg(winagg(s,100,30))  
  from Stream s  
 where id(source(s))=2;
```

- Lots of work on similarity search over time sequences
- Indexing time series

Bulut and Singh: A Unified Framework for Monitoring Data Streams in Real Time, ICDE 2005

Zhu and Shasha: Warping Indexes with Envelope Transforms for Query by Humming, SIGMOD 2003

# Scientific Databases

- Optimization of queries with numerical functions

Wolniewicz and Graefe: Algebraic Optimization of Computations over Scientific Databases, VLDB 1999

- Function approximation and caching

Panda, Riedewald, Pope, Gehrke, Chew: Indexing for Function Approximation, VLDB 2006

Denny & Franklin: Adaptive Execution of Variable-Accuracy Functions, VLDB 2006

# Scientific Databases

- Scientific workflows

Berkley et al: Incorporating Semantics in Scientific Workflow Authoring, SSDBM 2005

- Tracking changes and sources

Buneman et al: Provenance Management in Curated Databases, SIGMOD 2006

- Spatial indexing (c.f. multimedia databases)

Csabail et al: Spatial Indexing of Large Multidimensional Databases, CIDR 2007