

UDBL

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Introduction to Object-Oriented and Object-Relational Database Systems



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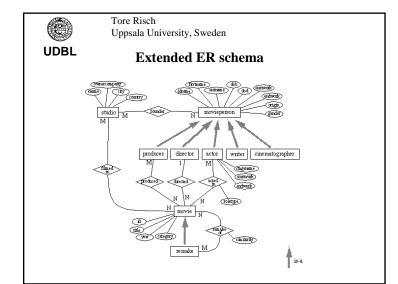
Database Design

•Logical Database Design:

How to translate a schema in the conceptual data model (e.g. extended ER-schemas) to a schema in the DBMS data model (e.g relational tables)

PROBLEM:

Semantics may disappear or be blurred when data is translated from extended ER-model to less expressive relational data model





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Database Design

•Physical Database Design:

E.g by indexes:

- permit fast matching of records in table satisfying certain search conditions.

PROBLEM:

New applications may require data and index structures that are not supported by the DBMS.

E.g. calendars, numerical arrays, geographical data, images, text, voice, etc.



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Database Manipulation

- •Typical query language *operations* are:
 - Searching for records fulfilling certain selection conditions
 - Iterating over entire tables applying *update operations*

PROBLEM: Would like to be able to customize and extend query language for different application areas, maps, time series, images, etc.



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Classical DBMSs

Applications:

Administrative applications

e.g. banking (ATMs)

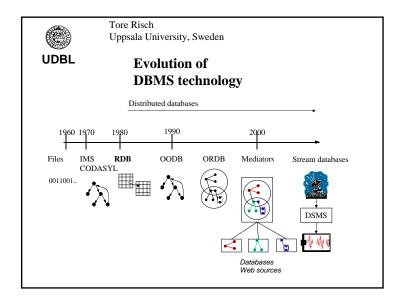
•Properties:

Very large structured data volumes

Very many small Transactions On-line (High transaction rates)

Occasional batch programs

High Security/Consistency





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New DBMS applications areas

CAD Computer Aided Design

Multi-media databases (images, maps, voice, time series,...)

Scientific Applications (measurements, logs)

Hypertext databases/documents (WWW/HTML/XML)



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New DBMS applications (for OODB)

•New needs for e.g. CAD and scientific databases:

Extensibility (on all levels)

Better performance

Expressability

E.g. Object-Orientation needed

Tight programming language interfaces

E.g. C++, Java

Long transactions

E.g. Engineering requires checkin/checkout model

Very large objects



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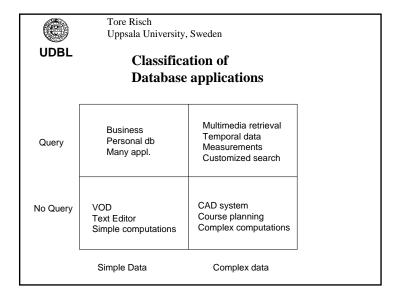
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Kinds of **DBMS** support

Query	Relational DBMSs	Object-Relational DBMSs
No Query	File systems (scalable) Storage managers	Object Stores

Simple Data

Complex data





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Object Stores (OO databases)

•First generation ODBs (around 1990)

Extend OO programming language with DBMS primitives

E.g. Smalltalk, C++, Java

Allow *persistent* data structures in C++ programs

Navigate through database using C++ primitives (as CODASYL)

An object store for C++

•Many products, e.g.:

Objectivity, Versant, ObjectStore

•Special embedded (C++/Java) OO Query language proposal: OQL



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Object Stores (OO databases)

- •Pros and cons:
 - + Long transactions with checkin/checkout model
 - + Always same *host* language (C++/Java)
 - + High efficiency only for checked-out data
 - Primitive 'query languages'
 - No methods in database (all code executes in client)
 - Rudimentary data independence (no views)
 - Limited concurrency
 - -Unsafe, database may crash
 - -Slow for many small transactions (e.g. ATM applications)



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Orthogonal Persistence in Object Stores

Pointer swizzling:

Automatic conversion from disk addresses to pointers References to data structures on disk (OIDs) look like regular C++/Java pointers/references!

Navigational access style.

Fast when database cached in main-memory of client! Preprocessed by OODBMS for convenient extention of C++ (JDK support in Pjama research project)



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Orthogonal Persistence in Object Stores

Integrated with programming language

E.g. C++/Java with persistent objects (e.g. ObjectStore/Pjama) class PERSON { ... };

....

{PERSON P; // Local within block... }

static PERSON p; // Local for execution

persistent PERSON p; // Exists between program executions



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Object-Relational Databases

- •Second generation ODBs (around 1997)
- •Idea:

Extend on RDBMS functionality

Customized (abstract) data types

Customized index structures

Customized query optimizers

Use declarative query language, SQL:99

•Extensible DBMS technology:

 ${\it Object-orientation}\ {\rm for\ abstract\ data\ types}$

Data blades provide:

User definable index structures

Cost hints and for the query optimizer



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Object-Relational Databases

- •Pros and cons:
 - + Support for high-level SQL queries, compatibility
 - + Views, logical data independence possible with queries
 - + Programming language independence
 - + Stored procedures, triggers, constraints
 - + High transaction performance by avoiding data shipping
 - Overkill for application needing just a C++ object store
 Performance may suffer compared to OODBs for applications needing just an object store
 - May be very difficult to extend index structures and query optimizers



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OO/OR Comparison

Object identity

E.g. for structure sharing:

Unique OIDs maintained by DBMS

E.g. OR:

create Person instances :tore, :kalle, :ulla;

In OO: use OO programming (C++, Java) constructs.

•Complex objects

Not only tables, numbers, strings

But sets, bags, lists, and arrays, i.e. non-1NF relations.

E.g. OR: set courses(:tore) = {:c1,:c2,:c3};

OO: use OO programming constructs in e.g. C++.



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Object-to-Relational Bridges

•Idea:

Object-Store with relational database as back-end Persistent objects in Java stored in relational databases Interface stubs generated for easy Java programming Some query language support

Pros and cons:

- + No need to develop new storage manager
- + Scalable search in back-end possible using JDBC/SQL
- Slower than JDBC
- No control over database schema!

Products: Hibernate, ObjectRelationalBridge (Apache), Castor, TopLink(Oracle)



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OO/OR Comparison

Extensibility

User defined data types, OR: create type Picture;

create type Timepoint;

User Defined Functions (UDFs) on new datatypes, OR: similar (Picture, Picture) ->Number

Extensible query operators through UDFs, OR:

select t1.image, t2.image

from albums1 t1, albums2 t2

where similar(t1.image,t2.image)>0.9;

OO has abstract datatypes through OO host language

OR databases: Also extensible indexing and query processing



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OO/OR Comparison

 $\hbox{\bf \bullet Class Hierarchies as modelling tool (both OO/OR)}$

Classification

E.g. OR:

create type Student under Person;

Students are subsets of persons.

Specialization

Student subtype of Person with extra attributes University, Classes, ...



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OO/OR Comparison

Concurrency

OO databases: Long transactions with checkin/checkout

OR database: Normally short transactions

•Ad Hoc Query Facility

OO Databases: Weak

OR Databases: Very strong and extensible

•Data independence

OO Databases: Very weak

OR Databases: Strong, e.g. using views



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OO/OR Comparison

•Computational completeness

OR databases: SQL:99: Turing complete stored procedures

language executed in database server

OO Databases: C++/Java code executed in client

Persistence

OR databases: Embedded queries to access persistent objects

OO databases: Transparent access to persistent objects

by swizzling

•Secondary storage management

OR databases: Indexes can be implemented by user

(difficult!)



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Object-Oriented DBMS Standard

•The ODMG standard proposal:

R. Cattell, Ed.: The ODMG-93 Standard

for Object Databases

Morgan-Kaufmann Publishers, San Mateo,

California, 1993.

http://www.odbms.org/odmg.html

•The SQL:99 standard proposal:

ISO standard